

# Modeling Packet Scheduling Algorithms in IP Routers

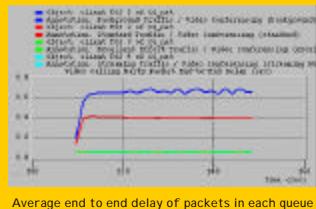
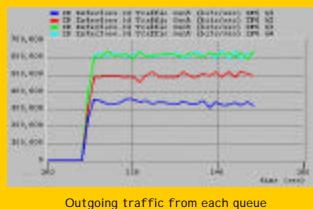
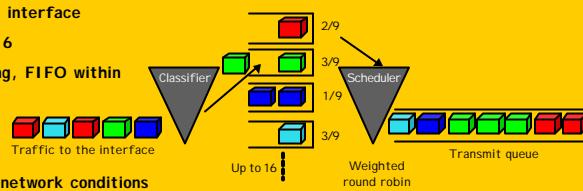
Nazy Alborz, Bruce Chen, and Ljiljana Trajkovic

{nalborz, bchenb, ljiljja}@cs.sfu.ca, <http://www.ensc.sfu.ca/research/cnl>

School of Engineering Science, Simon Fraser University

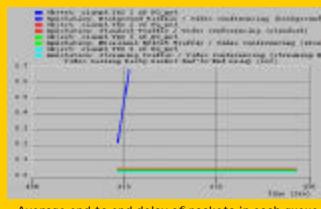
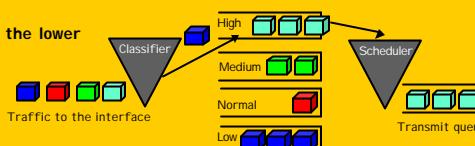
## Custom Queuing

- Packet classification: based on protocol or interface
- Number of queues: equal to or less than 16
- Scheduling: weighted round robin scheduling, FIFO within priority queues
- Usage: not used for voice
- Advantage: guarantees BW per queue
- Disadvantage: does not adapt to changing network conditions



## Priority Queuing

- Packet classification: based on protocol or interface
- Number of queues (4): high, medium, normal, and low
- Scheduling: serves packets from each queue (high to low priority) in FIFO order until each queue becomes empty
- Disadvantages:
  - delay is transferred from the high priority queue to the lower priority queues
  - lower priority queues may starve
- Usage:
  - works well for a small number of traffic types
  - not used for voice

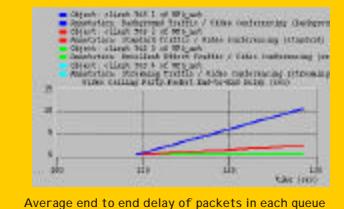
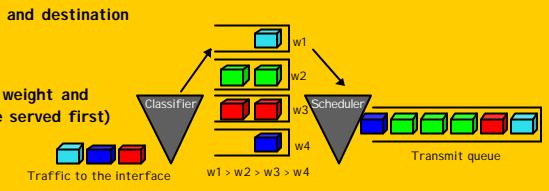


### References:

- [1] L. Zhang, "VirtualClock: A new traffic control algorithm for packet switching networks," ACM SIGCOMM, Sep. 1990.
- [2] Cisco Systems Inc. documentation on QoS: <http://www.cisco.com/wrap/public/732/tech/quality.shtml>
- [3] OPNET Technologies Inc., Washington DC, OPNET documentation V.7.0.L.

## Weighted Fair Queuing

- Packet classification: based on protocol or source and destination addresses
- Number of queues: equal to the number of flows
- Scheduling: "fair" algorithm based on both queue weight and the arrival time of each packet (lower weights are served first)
- Usage:
  - works well for data traffic
  - not recommended for voice, but can be used when no alternative is available
- Advantage: dynamically adapts to changing network conditions



## VirtualClock

- Each flow has its own clock
- Assigns a separate queue to each flow according to:
  - IP address of the source
  - IP address of the destination
- Upon receiving a packet from flow<sub>i</sub>,
  - $VClock_i = \max(VClock_i, realtime)$
  - clock ticks with  $(1/AR_i)$ ,  $VClock_i \leq VClock_i + VTick_i (= 1/AR_i)$
- Stamps packets with  $VClock_i$  values
- Services packets in increasing stamp order

