

## Markov Chains

**Stochastic process:** family  $\{X_i; i \in I\}$  of rvs  $I$  the **index set**. Often  $I \subset \mathbb{R}$ , e.g.  $[0, \infty)$ ,  $[0, 1]$   $\mathbb{Z}$  or  $\mathbb{N}$ .

**Continuous time:**  $I$  is an interval

**Discrete time:**  $I \subset \mathbb{Z}$ .

Generally all  $X_n$  take values in **state space**  $S$ . In following  $S$  is a finite or countable set; each  $X_n$  is discrete.

Usually  $S$  is  $\mathbb{Z}$ ,  $\mathbb{N}$  or  $\{0, \dots, m\}$  for some finite  $m$ .

**Markov Chain:** stochastic process  $X_n; n \in \mathbb{N}$ . taking values in a finite or countable set  $S$  such that for every  $n$  and every event of the form

$$A = \{(X_0, \dots, X_{n-1}) \in B \subset S^n\}$$

we have

$$P(X_{n+1} = j | X_n = i, A) = P(X_1 = j | X_0 = i) \quad (1)$$

Notation:  $\mathbf{P}$  is the (possibly infinite) array with elements

$$P_{ij} = P(X_1 = j | X_0 = i)$$

indexed by  $i, j \in S$ .

$\mathbf{P}$  is the (one step) **transition matrix** of the Markov Chain.

WARNING: in (1) we require the condition to hold **only** when

$$P(X_n = i, A) > 0$$

Evidently the entries in  $\mathbf{P}$  are non-negative and

$$\sum_j P_{ij} = 1$$

for all  $i \in S$ . Any such matrix is called **stochastic**.

We define powers of  $\mathbf{P}$  by

$$(\mathbf{P}^n)_{ij} = \sum_k (\mathbf{P}^{n-1})_{ik} P_{kj}$$

Notice that even if  $S$  is infinite these sums converge absolutely.

## Chapman-Kolmogorov Equations

Condition on  $X_{l+n-1}$  to compute

$$P(X_{l+n} = j | X_l = i)$$

$$\begin{aligned} P(X_{l+n} = j | X_l = i) &= \sum_k P(X_{l+n} = j, X_{l+n-1} = k | X_l = i) \\ &= \sum_k P(X_{l+n} = j | X_{l+n-1} = k, X_l = i) \\ &\quad \times P(X_{l+n-1} = k | X_l = i) \\ &= \sum_k P(X_1 = j | X_0 = k) \\ &\quad \times P(X_{l+n-1} = k | X_l = i) \\ &= \sum_k P(X_{l+n-1} = k | X_l = i) \mathbf{P}_{kj} \end{aligned}$$

Now condition on  $X_{l+n-2}$  to get

$$\begin{aligned} P(X_{l+n} = j | X_l = i) &= \\ &\sum_{k_1 k_2} \mathbf{P}_{k_1 k_2} \mathbf{P}_{k_2 j} P(X_{l+n-2} = k_1 | X_l = i) \end{aligned}$$

Notice: sum over  $k_2$  computes  $k_1, j$  entry in matrix  $\mathbf{P}\mathbf{P} = \mathbf{P}^2$ .

$$P(X_{l+n} = j | X_l = i) = \sum_{k_1} (\mathbf{P}^2)_{k_1, j} P(X_{l+n-2} = k_1 | X_l = i)$$

We may now prove by induction on  $n$  that

$$P(X_{l+n} = j | X_l = i) = (\mathbf{P}^n)_{ij}.$$

This proves Chapman-Kolmogorov equations:

$$\begin{aligned} P(X_{l+m+n} = j | X_l = i) &= \\ \sum_k P(X_{l+m} = k | X_l = i) & \\ \times P(X_{l+m+n} = j | X_{l+m} = k) & \end{aligned}$$

These are simply a restatement of the identity

$$\mathbf{P}^{n+m} = \mathbf{P}^n \mathbf{P}^m.$$

**Remark:** It is important to notice that these probabilities depend on  $m$  and  $n$  but **not** on  $l$ . We say the chain has **stationary** transition probabilities. A more general definition of Markov chain than (1) is

$$\begin{aligned} P(X_{n+1} = j | X_n = i, A) \\ = P(X_{n+1} = j | X_n = i). \end{aligned}$$

Notice RHS now permitted to depend on  $n$ .

Define  $\mathbf{P}^{n,m}$ : matrix with  $i, j$ th entry

$$P(X_m = j | X_n = i)$$

for  $m > n$ . Then

$$\mathbf{P}^{r,s} \mathbf{P}^{s,t} = \mathbf{P}^{r,t}$$

Also called Chapman-Kolmogorov equations. This chain does not have stationary transitions.

**Remark:** The calculations above involve sums in which all terms are positive. They therefore apply even if the state space  $S$  is countably infinite.

## Extensions of the Markov Property

Function  $f(x_0, x_1, \dots)$  defined on  $S^\infty =$  all infinite sequences of points in  $S$ .

Let  $B_n$  be the event

$$f(X_n, X_{n+1}, \dots) \in C$$

for suitable  $C$  in range space of  $f$ . Then

$$P(B_n | X_n = x, A) = P(B_0 | X_0 = x) \quad (2)$$

for any event  $A$  of the form

$$\{(X_0, \dots, X_{n-1}) \in D\}$$

Also

$$P(AB_n | X_n = x) = P(A | X_n = x)P(B_n | X_n = x) \quad (3)$$

“Given the present the past and future are conditionally independent.”

Proof of (2):

Special case:

$$B_n = \{(X_n = x, X_{n+1} = x_1, \dots, X_{n+m} = x_m)\}$$

LHS of (2) evaluated by repeated conditioning  
(cf. Chapman-Kolmogorov):

$$\mathbf{P}_{x,x_1} \mathbf{P}_{x_1,x_2} \cdots \mathbf{P}_{x_{m-1},x_m}$$

Same for RHS.

Events defined from  $X_n, \dots, X_{n+m}$ : sum over  
appropriate vectors  $x, x_1, \dots, x_m$ .

General case: monotone class techniques.



To prove (3) write

$$\begin{aligned} P(AB_n|X_n = x) &= P(B_n|X_n = x, A)P(A|X_n = x) \\ &= P(B_n|X_n = x)P(A|X_n = x) \end{aligned}$$

using (2).

## Classification of States

If an entry  $\mathbf{P}_{ij}$  is 0 it is not possible to go from state  $i$  to state  $j$  in one step. It may be possible to make the transition in some larger number of steps, however. We say  $i$  **leads to**  $j$  (or  $j$  is accessible from  $i$ ) if there is an integer  $n \geq 0$  such that

$$P(X_n = j | X_0 = i) > 0.$$

We use the notation  $i \rightsquigarrow j$ . Define  $\mathbf{P}^0$  to be identity matrix  $\mathbf{I}$ . Then  $i \rightsquigarrow j$  if there is an  $n \geq 0$  for which  $(\mathbf{P}^n)_{ij} > 0$ .

States  $i$  and  $j$  **communicate** if  $i \rightsquigarrow j$  and  $j \rightsquigarrow i$ .

Write  $i \leftrightarrow j$  if  $i$  and  $j$  communicate.

Communication is an equivalence relation: reflexive, symmetric, transitive relation on states of  $S$ .

More precisely:

**Reflexive:** for all  $i$  we have  $i \leftrightarrow i$ .

**Symmetric:** if  $i \leftrightarrow j$  then  $j \leftrightarrow i$ .

**Transitive:** if  $i \leftrightarrow j$  and  $j \leftrightarrow k$  then  $i \leftrightarrow k$ .

Proof:

Reflexive: follows from inclusion of  $n = 0$  in definition of leads to.

Symmetry is obvious.

Transitivity: suffices to check that  $i \rightsquigarrow j$  and  $j \rightsquigarrow k$  imply that  $i \rightsquigarrow k$ . But if  $(\mathbf{P}^m)_{ij} > 0$  and  $(\mathbf{P}^n)_{jk} > 0$  then

$$\begin{aligned}(\mathbf{P}^{m+n})_{ik} &= \sum_l (\mathbf{P}^m)_{il} (\mathbf{P}^n)_{lk} \\ &\geq (\mathbf{P}^m)_{ij} (\mathbf{P}^n)_{jk} \\ &> 0\end{aligned}$$

Any equivalence relation on a set partitions the set into **equivalence classes**; two elements are in the same equivalence class if and only if they are equivalent.

Communication partitions  $S$  into equivalence classes called **communicating classes**.

Example:

$$\mathbf{P} = \begin{bmatrix} \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & 0 & 0 & 0 & 0 \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & 0 & 0 & 0 & 0 \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & 0 & 0 & 0 & 0 \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & 0 & 0 & 0 & 0 \\ \frac{1}{4} & \frac{1}{4} & 0 & 0 & \frac{1}{2} & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ 0 & 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix}$$

Find communicating classes: start with say state 1, see where it leads.

- $1 \rightsquigarrow 2$ ,  $1 \rightsquigarrow 3$  and  $1 \rightsquigarrow 4$  in row 1.
- Row 4:  $4 \rightsquigarrow 1$ . So: (transitivity) 1, 2, 3 and 4 all in the same communicating class.
- Claim: none of these leads to 5, 6, 7 or 8.

Suppose  $i \in \{1, 2, 3, 4\}$  and  $j \in \{5, 6, 7, 8\}$ . Then  $(\mathbf{P}^n)_{ij}$  is sum of products of  $\mathbf{P}_{kl}$ . Cannot be positive unless there is a sequence  $i_0 = i, i_1, \dots, i_n = j$  with  $\mathbf{P}_{i_{k-1}, i_k} > 0$  for  $k = 1, \dots, n$ .

Consider first  $k$  for which  $i_k \in \{5, 6, 7, 8\}$ . Then  $i_{k-1} \in \{1, 2, 3, 4\}$  and so  $\mathbf{P}_{i_{k-1}, i_k} = 0$ .

So:  $\{1, 2, 3, 4\}$  is a communicating class.

- $5 \rightsquigarrow 1, 5 \rightsquigarrow 2, 5 \rightsquigarrow 3$  and  $5 \rightsquigarrow 4$ .
- None of these lead to any of  $\{5, 6, 7, 8\}$  so  $\{5\}$  must be communicating class.
- Similarly  $\{6\}$  and  $\{7, 8\}$  are communicating classes.

Note: states 5 and 6 have special property. Each time you are in either state you run a risk of going to one of the states 1, 2, 3 or 4. Eventually you will make such a transition and then never return to state 5 or 6.

States 5 and 6 are **transient**.

To make this precise define hitting times:

$$T_k = \min\{n > 0 : X_n = k\}$$

We define

$$f_k = P(T_k < \infty | X_0 = k)$$

State  $k$  is **transient** if  $f_k < 1$  and **recurrent** if  $f_k = 1$ .



Let  $N_k$  be number of times chain is ever in state  $k$ .

Claims:

1. If  $f_k < 1$  then  $N_k$  has a Geometric distribution:

$$P(N_k = r | X_0 = k) = f_k^{r-1} (1 - f_k)$$

for  $r = 1, 2, \dots$

2. If  $f_k = 1$  then

$$P(N_k = \infty | X_0 = k) = 1$$

Proof using **Strong Markov Property**:

**Stopping time** for the Markov chain is a random variable  $T$  taking values in  $\{0, 1, \dots\} \cup \{\infty\}$  such that for each finite  $k$  there is a function  $f_k$  such that

$$1(T = k) = f_k(X_0, \dots, X_k)$$

Notice that  $T_k$  in theorem is a stopping time.

Standard shorthand notation: by

$$P^x(A)$$

we mean

$$P(A|X_0 = x).$$

Similarly we define

$$E^x(Y) = E(Y|X_0 = x).$$

Goal: explain and prove

$$E(f(X_T, \dots) | X_T, \dots, X_0) = E^{X_T}(f(X_0, \dots))$$

Simpler claim:

$$P(X_{T+1} = j | X_T = i) = \mathbf{P}_{ij} = P^i(X_1 = j).$$

Notation:  $A_k = \{X_k = i, T = k\}$

Notice:  $A_k = \{X_T = i, T = k\}$ :

$$\begin{aligned} P(X_{T+1} = j | X_T = i) &= \frac{P(X_{T+1} = j, X_T = i)}{P(X_T = i)} \\ &= \frac{\sum_k P(X_{T+1} = j, X_T = i, T = k)}{\sum_k P(X_T = i, T = k)} \\ &= \frac{\sum_k P(X_{k+1} = j, A_k)}{\sum_k P(A_k)} \\ &= \frac{\sum_k P(X_{k+1} = j | A_k) P(A_k)}{\sum_k P(A_k)} \\ &= \frac{\sum_k P(X_1 = j | X_0 = i) P(A_k)}{\sum_k P(A_k)} \\ &= \mathbf{P}_{i,j} \end{aligned}$$

Notice use of fact that  $T = k$  is event defined in terms of  $X_0, \dots, X_k$ .

Technical problems with proof:

- It might be that  $P(T = \infty) > 0$ . What are  $X_T$  and  $X_{T+1}$  on the event  $T = \infty$ .

Answer: condition also on  $T < \infty$ .

- Prove formula only for stopping times where  $\{T < \infty\} \cap \{X_T = i\}$  has positive probability.

We will now fix up these technical details.

Suppose  $f(x_0, x_1, \dots)$  is a (measurable) function on  $S^{\mathbb{N}}$ . Put

$$Y_n = f(X_n, X_{n+1}, \dots).$$

Assume  $E(|Y_0| | X_0 = x) < \infty$  for all  $x$ . Claim:

$$E(Y_n | X_n, A) = E^{X_n}(Y_0) \quad (4)$$

whenever  $A$  is any event defined in terms of  $X_0, \dots, X_n$ .

### **Proof:**

**1** Family of  $f$  for which claim holds includes all indicators; see extension of Markov Property in previous lecture.

**2** family of  $f$  for which claim is true is vector space (so if  $f, g$  in family then so is  $af + bg$  for any constants  $a$  and  $b$ ).

- So family of  $f$  for which claim is true includes all simple functions.
- family of  $f$  for which claim true is closed under monotone increasing limits (of non-negative  $f_n$ ) by the Monotone Convergence theorem.
- So claim true for every non-negative integrable  $f$ .
- Claim follows for integrable  $f$  by linearity.

Aside on “measurable”: what sorts of events can be defined in terms of a family  $\{Y_i : i \in I\}$ ?

Natural: any event of form  $(Y_{i_1}, \dots, Y_{i_k}) \in C$  is “defined in terms of the family” for any finite set  $i_1, \dots, i_k$  and any (Borel) set  $C$  in  $S^k$ .

For countable  $S$ : each singleton  $(s_1, \dots, s_k) \in S^k$  Borel. So every subset of  $S^k$  Borel.

Natural: if you can define each of a sequence of events  $A_n$  in terms of the  $Y$ s then the definition “there exists an  $n$  such that (definition of  $A_n$ )...” defines  $\cup A_n$ .

Natural: if  $A$  is definable in terms of the  $Y$ s then  $A^c$  can be defined from the  $Y$ s by just inserting the phrase “It is not true that” in front of the definition of  $A$ .

So family of events definable in terms of the family  $\{Y_i : i \in I\}$  is a  $\sigma$ -field which includes every event of the form  $(Y_{i_1}, \dots, Y_{i_k}) \in C$ . We call the smallest such  $\sigma$ -field,  $\mathcal{F}(\{Y_i : i \in I\})$ , the  $\sigma$ -field generated by the family  $\{Y_i : i \in I\}$ .

Using the Markov property:

Toss coin till I get a head. What is the expected number of tosses?

Define state to be 0 if toss is tail and 1 if toss is heads.

Define  $X_0 = 0$ .

Let  $N = \min\{n > 0 : X_n = 1\}$ . Want

$$E(N) = E^0(N)$$

Note: if  $X_1 = 1$  then  $N = 1$ . If  $X_1 = 0$  then  $N = 1 + \min\{n > 0 : X_{n+1} = 1\}$ .

In symbols:

$$N = \min\{n > 0 : X_n = 1\} = f(X_1, X_2, \dots)$$

and

$$N = 1 + 1(X_1 = 0)f(X_2, X_3, \dots)$$



Take expected values starting from 0:

$$E^0(N) = 1 + E^0\{1(X_1 = 0)f(X_2, X_3, \dots)\}$$

Condition on  $X_1$  and get

$$E^0(N) = 1 + E^0[E\{1(X_1 = 0)f(X_2, \dots)|X_1\}]$$

But

$$\begin{aligned} E\{1(X_1 = 0)f(X_2, X_3, \dots)|X_1\} \\ &= 1(X_1 = 0)E^{X_1}\{f(X_1, X_2, \dots)\} \\ &= 1(X_1 = 0)E^0\{f(X_1, X_2, \dots)\} \\ &= 1(X_1 = 0)E^0(N) \end{aligned}$$

so that

$$E^0(N) = 1 + pE^0\{N\}$$

where  $p$  is the probability of tails. Solve for  $E(N)$  to get

$$E(N) = \frac{1}{1-p}$$

This is the formula for expected value of the sort of geometric which starts at 1 and has  $p$  being the probability of failure.

## Initial Distributions

Meaning of unconditional expected values?

Markov property specifies only cond'l probs; no way to deduce marginal distributions.

For every dstbn  $\pi$  on  $S$  and transition matrix  $\mathbf{P}$  there is a a stochastic process  $X_0, X_1, \dots$  with

$$P(X_0 = k) = \pi_k$$

and which is a Markov Chain with transition matrix  $\mathbf{P}$ .

Note Strong Markov Property proof used only conditional expectations.

Notation:  $\pi$  a probability on  $S$ .  $E^\pi$  and  $P^\pi$  are expected values and probabilities for chain with initial distribution  $\pi$ .

Summary of easily verified facts:

- For any sequence of states  $i_0, \dots, i_k$

$$P(X_0 = i_0, \dots, X_k = i_k) = \pi_{i_0} \mathbf{P}_{i_0 i_1} \cdots \mathbf{P}_{i_{k-1} i_k}$$

- For any event  $A$ :

$$\mathbf{P}^\pi(A) = \sum_k \pi_k \mathbf{P}^k(A)$$

- For any bounded rv  $Y = f(X_0, \dots)$

$$\mathbf{E}^\pi(Y) = \sum_k \pi_k \mathbf{E}^k(Y)$$

## Recurrence and Transience

Now consider a transient state  $k$ , that is, a state for which

$$f_k = P^k(T_k < \infty) < 1$$

Note that  $T_k = \min\{n > 0 : X_n = k\}$  is a stopping time. Let  $N_k$  be the number of visits to state  $k$ . That is

$$N_k = \sum_{n=0}^{\infty} 1(X_n = k)$$

Notice that if we define the function

$$f(x_0, x_1, \dots) = \sum_{n=0}^{\infty} 1(x_n = k)$$

then

$$N_k = f(X_0, X_1, \dots)$$

Notice, also, that on the event  $T_k < \infty$

$$N_k = 1 + f(X_{T_k}, X_{T_k+1}, \dots)$$

and on the event  $T_k = \infty$  we have

$$N_k = 1$$

In short:

$$N_k = 1 + f(X_{T_k}, X_{T_k+1}, \dots) \mathbf{1}(T_k < \infty)$$

Hence

$$\begin{aligned} \mathbf{P}^k(N_k = r) &= \mathbf{E}^k \{P(N_k = r | \mathcal{F}_T)\} \\ &= \mathbf{E}^k \left[ P \left\{ \mathbf{1} + f(X_{T_k}, X_{T_k+1}, \dots) \right. \right. \\ &\quad \left. \left. \times \mathbf{1}(T_k < \infty) = r | \mathcal{F}_T \right\} \right] \\ &= \mathbf{E}^k [\mathbf{1}(T_k < \infty) \\ &\quad \times P^{X_{T_k}} \{f(X_0, X_1, \dots) = r - 1\}] \\ &= \mathbf{E}^k \left\{ \mathbf{1}(T_k < \infty) P^k(N_k = r - 1) \right\} \\ &= \mathbf{E}^k \{ \mathbf{1}(T_k < \infty) \} P^k(N_k = r - 1) \\ &= f_k P^k(N_k = r - 1) \end{aligned}$$

It is easily verified by induction, then, that

$$\mathbf{P}^k(N_k = r) = f_k^{r-1} P^k(N_k = 1)$$

But  $N_k = 1$  if and only if  $T_k = \infty$  so

$$\mathbf{P}^k(N_k = r) = f_k^{r-1}(1 - f_k)$$

so  $N_k$  has (chain starts from  $k$ ) Geometric dist'n, mean  $1/(1 - f_k)$ . Argument also shows that if  $f_k = 1$  then

$$P^k(N_k = 1) = P^k(N_k = 2) = \dots$$

which can only happen if all these probabilities are 0. Thus if  $f_k = 1$

$$P(N_k = \infty) = 1$$

Since

$$N_k = \sum_{n=0}^{\infty} \mathbf{1}(X_n = k)$$

$$\mathbb{E}^k(N_k) = \sum_{n=0}^{\infty} (\mathbf{P}^n)_{kk}$$

So: State  $k$  is transient if and only if

$$\sum_{n=0}^{\infty} (\mathbf{P}^n)_{kk} < \infty$$

and this sum is  $1/(1 - f_k)$ .

**Proposition 1** *Recurrence (or transience) is a class property. That is, if  $i$  and  $j$  are in the same communicating class then  $i$  is recurrent (respectively transient) if and only if  $j$  is recurrent (respectively transient).*

**Proof:** Suppose  $i$  is recurrent and  $i \leftrightarrow j$ . There are integers  $m$  and  $n$  such that

$$(\mathbf{P}^m)_{ji} > 0 \quad \text{and} \quad (\mathbf{P}^n)_{ij} > 0$$

Then

$$\begin{aligned} \sum_k (\mathbf{P}^k)_{jj} &\geq \sum_{k \geq 0} (\mathbf{P}^{m+k+n})_{jj} \\ &\geq \sum_{k \geq 0} (\mathbf{P}^m)_{ji} (\mathbf{P}^k)_{ii} (\mathbf{P}^n)_{ij} \\ &= (\mathbf{P}^m)_{ji} \left\{ \sum_{k \geq 0} (\mathbf{P}^k)_{ii} \right\} (\mathbf{P}^n)_{ij} \end{aligned}$$

The middle term is infinite and the two outside terms positive so

$$\sum_k (\mathbf{P}^k)_{jj} = \infty$$

which shows  $j$  is recurrent.

A finite state space chain has at least one recurrent state:

If all states were transient we would have for each  $k$   $P(N_k < \infty) = 1$ . This would mean  $P(\forall k N_k < \infty) = 1$ . But for any  $\omega$  there must be at least one  $k$  for which  $N_k = \infty$  (the total of a finite list of finite numbers is finite).

Infinite state space chain may have all states transient:

The chain  $X_n$  satisfying  $X_{n+1} = X_n + 1$  on the integers has all states transient.



More interesting example:

- Toss a coin repeatedly.
- Let  $X_n$  be  $X_0$  plus the number of heads minus the number of tails in the first  $n$  tosses.
- Let  $p$  denote the probability of heads on an individual trial.

$X_n - X_0$  is a sum of  $n$  iid random variables  $Y_i$  where  $P(Y_i = 1) = p$  and  $P(Y_i = -1) = 1 - p$ .

SLLN shows  $X_n/n$  converges almost surely to  $2p - 1$ . If  $p \neq 1/2$  this is not 0.

In order for  $X_n/n$  to have a positive limit we must have  $X_n \rightarrow \infty$  almost surely so all states are visited only finitely many times. That is, all states are transient. Similarly for  $p < 1/2$   $X_n \rightarrow -\infty$  almost surely and all states are transient.

Now look at  $p = 1/2$ . The law of large numbers argument no longer shows anything. I will show that all states are recurrent.

Proof: We evaluate  $\sum_n (\mathbf{P}^n)_{00}$  and show the sum is infinite. If  $n$  is odd then  $(\mathbf{P}^n)_{00} = 0$  so we evaluate

$$\sum_m (\mathbf{P}^{2m})_{00}$$

Now

$$(\mathbf{P}^{2m})_{00} = \binom{2m}{m} 2^{-2m}$$

According to Stirling's approximation

$$\lim_{m \rightarrow \infty} \frac{m!}{m^{m+1/2} e^{-m} \sqrt{2\pi}} = 1$$

Hence

$$\lim_{m \rightarrow \infty} \sqrt{m} (\mathbf{P}^{2m})_{00} = \frac{1}{\sqrt{\pi}}$$

Since

$$\sum \frac{1}{\sqrt{m}} = \infty$$

we are done.

## Mean return times

Compute expected times to return. For  $x \in S$  let  $T_x$  denote the hitting time for  $x$ .

Suppose  $x$  recurrent in **irreducible** chain (only one communicating class).

Derive equations for expected values of different  $T_x$ .

Each  $T_x$  is a certain function  $f_x$  applied to  $X_1, \dots$ . Setting  $\mu_{ij} = \mathbb{E}^i(T_j)$  we find

$$\mu_{ij} = \sum_k \mathbb{E}^i(T_j 1(X_1 = k))$$

Note that if  $X_1 = x$  then  $T_x = 1$  so

$$\mathbb{E}^i(T_j 1(X_1 = j)) = P_{ij}$$

For  $k \neq j$ , if  $X_1 = k$  then

$$T_j = 1 + f_j(X_2, X_3, \dots)$$

and, by conditioning on  $X_1 = k$  we find

$$\mathbb{E}^i(T_j \mathbf{1}(X_1 = k)) = \mathbf{P}_{ik} \{1 + \mathbb{E}^k(T_j)\}$$

This gives

$$\mu_{ij} = 1 + \sum_{k \neq j} \mathbf{P}_{ik} \mu_{kj} \quad (5)$$

Technically, I should check that the expectations in (5) are finite. All the random variables involved are non-negative, however, and the equation actually makes sense even if some terms are infinite. (To prove this you actually study

$$T_{x,n} = \min(T_x, n)$$

deriving an identity for a fixed  $n$ , letting  $n \rightarrow \infty$  and applying the monotone convergence theorem.)

Here is a simple example:

$$\mathbf{P} = \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & 0 & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix}$$

The identity (5) becomes

$$\mu_{1,1} = 1 + \frac{1}{2}\mu_{2,1} + \frac{1}{2}\mu_{3,1}$$

$$\mu_{1,2} = 1 + \frac{1}{2}\mu_{3,2}$$

$$\mu_{1,3} = 1 + \frac{1}{2}\mu_{2,3}$$

$$\mu_{2,1} = 1 + \frac{1}{2}\mu_{3,1}$$

$$\mu_{2,2} = 1 + \frac{1}{2}\mu_{1,2} + \frac{1}{2}\mu_{3,2}$$

$$\mu_{2,3} = 1 + \frac{1}{2}\mu_{1,3}$$

$$\mu_{3,1} = 1 + \frac{1}{2}\mu_{2,1}$$

$$\mu_{3,2} = 1 + \frac{1}{2}\mu_{1,2}$$

$$\mu_{3,3} = 1 + \frac{1}{2}\mu_{1,3} + \frac{1}{2}\mu_{2,3}$$

Seventh and fourth show  $\mu_{2,1} = \mu_{3,1}$ . Similar calculations give  $\mu_{ii} = 3$  and for  $i \neq j$   $\mu_{i,j} = 2$ .

**Example:** Coin tossing Markov Chain with  $p = 1/2$  shows situation can be different when  $S$  is infinite. Equations above become:

$$m_{0,0} = 1 + \frac{1}{2}m_{1,0} + \frac{1}{2}m_{-1,0}$$
$$m_{1,0} = 1 + \frac{1}{2}m_{2,0}$$

and many more.

Some observations:

Have to go through 1 to get to 0 from 2 so

$$m_{2,0} = m_{2,1} + m_{1,0}$$

Symmetry (switching H and T):

$$m_{1,0} = m_{-1,0}$$

Transition probabilities are **homogeneous**:

$$m_{2,1} = m_{1,0}$$

Conclusion:

$$\begin{aligned} m_{0,0} &= 1 + m_{1,0} \\ &= 1 + 1 + \frac{1}{2}m_{2,0} \\ &= 2 + m_{1,0} \end{aligned}$$

Notice that there are **no** finite solutions!

Summary of the situation:

Every state is recurrent.

All the expected hitting times  $m_{ij}$  are infinite.

All entries  $\mathbf{P}_{ij}^n$  converge to 0.

Jargon: The states in this chain are null recurrent.

Model: 2 state MC for weather: 'Dry' or 'Wet'.

```
> p:= matrix(2,2,[[3/5,2/5],[1/5,4/5]]);
```

$$p := \begin{bmatrix} 3/5 & 2/5 \\ 1/5 & 4/5 \end{bmatrix}$$

```
> p2:=evalm(p*p):
```

```
> p4:=evalm(p2*p2):
```

```
> p8:=evalm(p4*p4):
```

```
> p16:=evalm(p8*p8):
```

This computes the powers (evalm understands matrix algebra).

Fact:

$$\lim_{n \rightarrow \infty} \mathbf{P}^n = \begin{bmatrix} \frac{1}{3} & \frac{2}{3} \\ \frac{1}{3} & \frac{2}{3} \end{bmatrix}$$



```

> evalf(evalm(p));
      [.60000000000    .40000000000]
      [                ]
      [.20000000000    .80000000000]
> evalf(evalm(p2));
      [.44000000000    .56000000000]
      [                ]
      [.28000000000    .72000000000]
> evalf(evalm(p4));
      [.35040000000    .64960000000]
      [                ]
      [.32480000000    .67520000000]
> evalf(evalm(p8));
      [.3337702400    .6662297600]
      [                ]
      [.3331148800    .6668851200]
> evalf(evalm(p16));
      [.3333336197    .6666663803]
      [                ]
      [.3333331902    .6666668098]

```

Where did  $1/3$  and  $2/3$  come from?

Suppose we toss a coin  $P(H) = \alpha_D$  and start the chain with Dry if we get heads and Wet if we get tails.

Then

$$P(X_0 = x) = \begin{cases} \alpha_D & x = \text{Dry} \\ \alpha_W = 1 - \alpha_D & x = \text{Wet} \end{cases}$$

and

$$\begin{aligned} P(X_1 = x) &= \sum_y P(X_1 = x | X_0 = y) P(X_0 = y) \\ &= \sum_y \alpha_y P_{y,x} \end{aligned}$$

Notice last line is a matrix multiplication of row vector  $\alpha$  by matrix  $\mathbf{P}$ . A special  $\alpha$ : if we put  $\alpha_D = 1/3$  and  $\alpha_W = 2/3$  then

$$\begin{bmatrix} \frac{1}{3} & \frac{2}{3} \end{bmatrix} \begin{bmatrix} \frac{3}{5} & \frac{2}{5} \\ \frac{1}{5} & \frac{4}{5} \end{bmatrix} = \begin{bmatrix} \frac{1}{3} & \frac{2}{3} \end{bmatrix}$$

So: if  $P(X_0 = D) = 1/3$  then  $P(X_1 = D) = 1/3$  and analogously for  $W$ . This means that  $X_0$  and  $X_1$  have the same distribution.

A probability vector  $\alpha$  is called the initial distribution for the chain if

$$P(X_0 = i) = \alpha_i$$

A Markov Chain is **stationary** if

$$P(X_1 = i) = P(X_0 = i)$$

for all  $i$

Finding stationary initial distributions. Consider  $\mathbf{P}$  above. The equation

$$\alpha\mathbf{P} = \alpha$$

is really

$$\alpha_D = 3\alpha_D/5 + \alpha_W/5$$

$$\alpha_W = 2\alpha_D/5 + 4\alpha_W/5$$

The first can be rearranged to

$$\alpha_W = 2\alpha_D.$$

So can the second. If  $\alpha$  is probability vector then

$$\alpha_W + \alpha_D = 1$$

so we get

$$1 - \alpha_D = 2\alpha_D$$

leading to

$$\alpha_D = 1/3$$

Some more examples:

$$\mathbf{P} = \begin{bmatrix} 0 & 1/3 & 0 & 2/3 \\ 1/3 & 0 & 2/3 & 0 \\ 0 & 2/3 & 0 & 1/3 \\ 2/3 & 0 & 1/3 & 0 \end{bmatrix}$$

Set  $\alpha\mathbf{P} = \alpha$  and get

$$\alpha_1 = \alpha_2/3 + 2\alpha_4/3$$

$$\alpha_2 = \alpha_1/3 + 2\alpha_3/3$$

$$\alpha_3 = 2\alpha_2/3 + \alpha_4/3$$

$$\alpha_4 = 2\alpha_1/3 + \alpha_3/3$$

$$1 = \alpha_1 + \alpha_2 + \alpha_3 + \alpha_4$$

First plus third gives

$$\alpha_1 + \alpha_3 = \alpha_2 + \alpha_4$$

so both sums  $1/2$ . Continue algebra to get

$$(1/4, 1/4, 1/4, 1/4).$$

```
p:=matrix([[0,1/3,0,2/3],[1/3,0,2/3,0],
           [0,2/3,0,1/3],[2/3,0,1/3,0]]);
```

```

           [ 0      1/3      0      2/3]
           [
           [1/3      0      2/3      0 ]
p := [
           [
           [ 0      2/3      0      1/3]
           [
           [2/3      0      1/3      0 ]
```

```
> p2:=evalm(p*p);
```

```

           [5/9      0      4/9      0 ]
           [
           [ 0      5/9      0      4/9]
p2:= [
           [4/9      0      5/9      0 ]
           [
           [ 0      4/9      0      5/9]
```

```
> p4:=evalm(p2*p2):
```

```
> p8:=evalm(p4*p4):
```

```
> p16:=evalm(p8*p8):
```

```
> p17:=evalm(p8*p8*p):
```

```
> evalf(evalm(p16));
[.5000000116 , 0 , .4999999884 , 0]
[
]
[0 , .5000000116 , 0 , .4999999884]
[
]
[.4999999884 , 0 , .5000000116 , 0]
[
]
[0 , .4999999884 , 0 , .5000000116]
> evalf(evalm(p17));
[0 , .4999999961 , 0 , .5000000039]
[
]
[.4999999961 , 0 , .5000000039 , 0]
[
]
[0 , .5000000039 , 0 , .4999999961]
[
]
[.5000000039 , 0 , .4999999961 , 0]
```

```

> evalf(evalm((p16+p17)/2));
[.2500, .2500, .2500, .2500]
[
[.2500, .2500, .2500, .2500]
[
[.2500, .2500, .2500, .2500]
[
[.2500, .2500, .2500, .2500]

```

$\mathbf{P}^n$  doesn't converges but  $(\mathbf{P}^n + \mathbf{P}^{n+1})/2$  does.

Next example:

$$\mathbf{P} = \begin{bmatrix} \frac{2}{5} & \frac{3}{5} & 0 & 0 \\ \frac{1}{5} & \frac{4}{5} & 0 & 0 \\ 0 & 0 & \frac{2}{5} & \frac{3}{5} \\ 0 & 0 & \frac{1}{5} & \frac{4}{5} \end{bmatrix}$$



Solve  $\alpha\mathbf{P} = \alpha$ :

$$\alpha_1 = \frac{2}{5}\alpha_1 + \frac{1}{5}\alpha_2$$

$$\alpha_2 = \frac{3}{5}\alpha_1 + \frac{4}{5}\alpha_2$$

$$\alpha_3 = \frac{2}{5}\alpha_3 + \frac{1}{5}\alpha_4$$

$$\alpha_4 = \frac{3}{5}\alpha_3 + \frac{4}{5}\alpha_4$$

$$1 = \alpha_1 + \alpha_2 + \alpha_3 + \alpha_4$$

Second and fourth equations redundant. Get

$$\alpha_2 = 3\alpha_1$$

$$3\alpha_3 = \alpha_4$$

$$1 = 4\alpha_1 + 4\alpha_3$$

Pick  $\alpha_1$  in  $[0, 1/4]$ ; put  $\alpha_3 = 1/4 - \alpha_1$ .

$$\alpha = (\alpha_1, 3\alpha_1, 1/4 - \alpha_1, 3(1/4 - \alpha_1))$$

solves  $\alpha\mathbf{P} = \alpha$ . So solution is not unique.

```

> p:=matrix([[2/5,3/5,0,0],[1/5,4/5,0,0],
             [0,0,2/5,3/5],[0,0,1/5,4/5]]);
             [2/5    3/5    0    0 ]
             [
             [1/5    4/5    0    0 ]
p := [
       [ 0    0    2/5    3/5]
       [
       [ 0    0    1/5    4/5]

> p2:=evalm(p*p):
> p4:=evalm(p2*p2):
> p8:=evalm(p4*p4):
> evalf(evalm(p8*p8));
      [.2500000000 , .7500000000 , 0 , 0]
      [
      [.2500000000 , .7500000000 , 0 , 0]
      [
      [0 , 0 , .2500000000 , .7500000000]
      [
      [0 , 0 , .2500000000 , .7500000000]

```

Notice that rows converge but to two different vectors:

$$\alpha^{(1)} = (1/4, 3/4, 0, 0)$$

and

$$\alpha^{(2)} = (0, 0, 1/4, 3/4)$$

Solutions of  $\alpha\mathbf{P} = \alpha$  revisited? Check that

$$\alpha^{(1)}\mathbf{P} = \alpha^{(1)}$$

and

$$\alpha^{(2)}\mathbf{P} = \alpha^{(2)}$$

If  $\alpha = \lambda\alpha^{(1)} + (1 - \lambda)\alpha^{(2)}$  ( $0 \leq \lambda \leq 1$ ) then

$$\alpha\mathbf{P} = \alpha$$

so again solution is not unique.

Last example:

```
> p:=matrix([[2/5,3/5,0],[1/5,4/5,0],
             [1/2,0,1/2]]);
```

```
           [2/5    3/5    0 ]
           [
p := [1/5    4/5    0 ]
           [
           [1/2    0    1/2]
```

```
> p2:=evalm(p*p):
```

```
> p4:=evalm(p2*p2):
```

```
> p8:=evalm(p4*p4):
```

```
> evalf(evalm(p8*p8));
```

```
[.2500000000 .7500000000      0      ]
[
[.2500000000 .7500000000      0      ]
[
[.2500152588 .7499694824 .00001525878906]
```

## Interpretation of examples

- For some  $\mathbf{P}$  all rows converge to some  $\alpha$ . In this case this  $\alpha$  is a stationary initial distribution.
- For some  $\mathbf{P}$  the locations of zeros flip flop.  $\mathbf{P}^n$  does not converge. Observation: average

$$\frac{\mathbf{P} + \mathbf{P}^2 + \dots + \mathbf{P}^n}{n}$$

*does* converge.

- For some  $\mathbf{P}$  some rows converge to one  $\alpha$  and some to another. In this case the solution of  $\alpha\mathbf{P} = \alpha$  is not unique.

Basic distinguishing features: pattern of 0s in matrix  $\mathbf{P}$ .

## The ergodic theorem

Consider a finite state space chain. If  $x$  is a vector then the  $i$ th entry in  $\mathbf{P}x$  is

$$\sum_j \mathbf{P}_{ij} x_j$$

Rows of  $\mathbf{P}$  probability vectors, so a weighted average of the entries in  $x$ .

If weights strictly between 0, 1 and largest and smallest entries in  $x$  not same then  $\sum_j \mathbf{P}_{ij} x_j$  strictly between largest and smallest entries in  $x$ . In fact

$$\begin{aligned} \sum_j \mathbf{P}_{ij} x_j - \min(x_k) &= \sum_j \mathbf{P}_{ij} \{x_j - \min(x_k)\} \\ &\geq \min_j \{p_{ij}\} (\max\{x_k\} - \min\{x_k\}) \end{aligned}$$

and

$$\begin{aligned} \max\{x_j\} - \sum_j \mathbf{P}_{ij} x_j \\ &\geq \min_j \{p_{ij}\} (\max\{x_k\} - \min\{x_k\}) \end{aligned}$$

Now multiply  $\mathbf{P}^r$  by  $\mathbf{P}^m$ .

$ij$ th entry in  $\mathbf{P}^{r+m}$  is a weighted average of the  $j$ th column of  $\mathbf{P}^m$ .

So, if all the entries in row  $i$  of  $\mathbf{P}^r$  are positive and the  $j$ th column of  $\mathbf{P}^m$  is not constant, the  $i$ th entry in the  $j$ th column of  $\mathbf{P}^{r+m}$  must be strictly between the minimum and maximum entries of the  $j$ th column of  $\mathbf{P}^m$ .

In fact, fix a  $j$ .

$\bar{x}_m =$  maximum entry in column  $j$  of  $\mathbf{P}^m$

$\underline{x}_m$  the minimum entry.

Suppose all entries of  $\mathbf{P}^r$  are positive.

Let  $\delta > 0$  be the smallest entry in  $\mathbf{P}^r$ . Our argument above shows that

$$\bar{x}_{m+r} \leq \bar{x}_m - \delta(\bar{x}_m - \underline{x}_m)$$

and

$$\underline{x}_{m+r} \geq \underline{x}_m + \delta(\bar{x}_m - \underline{x}_m)$$

Putting these together gives

$$(\bar{x}_{m+r} - \underline{x}_{m+r}) \leq (1 - 2\delta)(\bar{x}_m - \underline{x}_m)$$

In summary the column maximum decreases, the column minimum increases and the gap between the two decreases exponentially along the sequence  $m, m + r, m + 2r, \dots$



This idea can be used to prove

**Proposition 2** *Suppose  $X_n$  finite state space Markov Chain with stationary transition matrix  $\mathbf{P}$ . Assume that there is a power  $r$  such that all entries in  $\mathbf{P}^r$  are positive. Then for  $\mathbf{P}^k$  has all entries positive for all  $k \geq r$  and  $\mathbf{P}^n$  converges, as  $n \rightarrow \infty$  to a matrix  $\mathbf{P}^\infty$ . Moreover,*

$$(\mathbf{P}^\infty)_{ij} = \pi_j$$

where  $\pi$  is the unique row vector satisfying

$$\pi = \pi \mathbf{P}$$

whose entries sum to 1.

**Proof:** First for  $k > r$

$$(\mathbf{P}^k)_{ij} = \sum_{\ell} (\mathbf{P}^{k-r})_{i\ell} (\mathbf{P}^r)_{\ell j}$$

For each  $i$  there is an  $\ell$  for which  $(\mathbf{P}^{k-r})_{i\ell} > 0$  and since  $(\mathbf{P}^r)_{\ell j} > 0$  we see  $(\mathbf{P}^k)_{ij} > 0$ .

The argument before the proposition shows that

$$\lim_{j \rightarrow \infty} \mathbf{P}^{m+jk}$$

exists for each  $m$  and  $k \geq r$ . This proves  $\mathbf{P}^n$  has a limit which we call  $\mathbf{P}^\infty$ . Since  $\mathbf{P}^{n-1}$  also converges to  $\mathbf{P}^\infty$  we find

$$\mathbf{P}^\infty = \mathbf{P}^\infty \mathbf{P}$$

Hence each row of  $\mathbf{P}^\infty$  is a solution of  $x\mathbf{P} = x$ . The argument before the statement of the proposition shows all rows of  $\mathbf{P}^\infty$  are equal. Let  $\pi$  be this common row.

Now if  $\alpha$  is any vector whose entries sum to 1 then  $\alpha\mathbf{P}^n$  converges to

$$\alpha\mathbf{P}^\infty = \pi$$

If  $\alpha$  is any solution of  $x = x\mathbf{P}$  we have by induction  $\alpha\mathbf{P}^n = \alpha$  so  $\alpha\mathbf{P}^\infty = \alpha$  so  $\alpha = \pi$ . That is exactly one vector whose entries sum to 1 satisfies  $x = x\mathbf{P}$ . •

Note conditions:

There is an  $r$  for which all entries in  $\mathbf{P}^r$  are positive.

The chain has a finite state space.

Consider finite state space case:  $\mathbf{P}^n$  need not have limit. Example:

$$\mathbf{P} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

Note  $\mathbf{P}^{2n}$  is the identity while  $\mathbf{P}^{2n+1} = \mathbf{P}$ .  
Note, too, that

$$\frac{\mathbf{P}^0 + \dots + \mathbf{P}^n}{n+1} \rightarrow \begin{bmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{bmatrix}$$

Consider the equations  $\pi = \pi\mathbf{P}$  with  $\pi_1 + \pi_2 = 1$ . We get

$$\pi_1 = \frac{1}{2}\pi_1 + \frac{1}{2}(1 - \pi_1) = \frac{1}{2}$$

so that the solution to  $\pi = \pi\mathbf{P}$  is again unique.

**Def'n:** The period  $d$  of a state  $i$  is the greatest common divisor of

$$\{n : (\mathbf{P}^n)_{ii} > 0\}$$

**Lemma 1** *If  $i \leftrightarrow j$  then  $i$  and  $j$  have the same period.*

**Def'n:** A state is **aperiodic** if its period is 1.

**Proof:** I do the case  $d = 1$ . Fix  $i$ . Let

$$G = \{k : (\mathbf{P}^k)_{ii} > 0\}$$

If  $k_1, k_2 \in G$  then  $k_1 + k_2 \in G$ .

This (and aperiodic) implies (number theory argument) that there is an  $r$  such that  $k \geq r$  implies  $k \in G$ .

Now find  $m$  and  $n$  so that

$$(\mathbf{P}^m)_{ij} > 0 \text{ and } (\mathbf{P}^n)_{ji} > 0$$

For  $k > r + m + n$  we see  $(\mathbf{P}^k)_{jj} > 0$  so the gcd of the set of  $k$  such that  $(\mathbf{P}^k)_{jj} > 0$  is 1. •

The case of period  $d > 1$  can be dealt with by considering  $\mathbf{P}^d$ .

$$\mathbf{P} = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

For this example  $\{1, 2, 3\}$  is a class of period 3 states and  $\{4, 5\}$  a class of period 2 states.

$$\mathbf{P} = \begin{bmatrix} 0 & 1/2 & 1/2 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

has a single communicating class of period 2.

A chain is **aperiodic** if all its states are aperiodic.

## Hitting Times

Start irreducible recurrent chain  $X_n$  in state  $i$ .  
Let  $T_j$  be first  $n > 0$  such that  $X_n = j$ . Define

$$m_{ij} = \mathbb{E}(T_j | X_0 = i)$$

First step analysis:

$$\begin{aligned} m_{ij} &= 1 \cdot P(X_1 = j | X_0 = i) \\ &\quad + \sum_{k \neq j} (1 + \mathbb{E}(T_j | X_0 = k)) P_{ik} \\ &= \sum_j P_{ij} + \sum_{k \neq j} P_{ik} m_{kj} \\ &= 1 + \sum_{k \neq j} P_{ik} m_{kj} \end{aligned}$$

Example

$$\mathbf{P} = \begin{bmatrix} \frac{3}{5} & \frac{2}{5} \\ \frac{1}{5} & \frac{4}{5} \end{bmatrix}$$

The equations are

$$\begin{aligned}m_{11} &= 1 + \frac{2}{5}m_{21} \\m_{12} &= 1 + \frac{3}{5}m_{12} \\m_{21} &= 1 + \frac{4}{5}m_{21} \\m_{22} &= 1 + \frac{1}{5}m_{12}\end{aligned}$$

The second and third equations give immediately

$$\begin{aligned}m_{12} &= \frac{5}{2} \\m_{21} &= 5\end{aligned}$$

Then plug in to the others to get

$$\begin{aligned}m_{11} &= 3 \\m_{22} &= \frac{3}{2}\end{aligned}$$

Notice stationary initial distribution is

$$\left( \frac{1}{m_{11}}, \frac{1}{m_{22}} \right)$$



Consider fraction of time spent in state  $j$ :

$$\frac{1(X_0 = j) + \cdots + 1(X_n = j)}{n + 1}$$

Imagine chain starts in chain  $i$ ; take expected value.

$$\frac{\sum_{r=1}^n \mathbf{P}_{ij}^r + 1(i = j)}{n + 1}$$

If rows of  $\mathbf{P}^r$  converge to  $\pi$  then fraction converges to  $\pi_j$ ; i.e. limiting fraction of time in state  $j$  is  $\pi_j$ .

Heuristic: start chain in  $i$ . Expect to return to  $i$  every  $m_{ii}$  time units. So are in state  $i$  about once every  $m_{ii}$  time units; i.e. limiting fraction of time in state  $i$  is  $1/m_{ii}$ .

Conclusion: for an irreducible recurrent finite state space Markov chain

$$\pi_i = \frac{1}{m_{ii}}.$$

Real proof: Renewal theorem or variant.

Idea:  $S_1 < S_2 < \dots$  are times of visits to  $i$ .  
Segment  $i$ :

$$X_{S_{i-1}+1}, \dots, X_{S_i}.$$

Segments are iid by Strong Markov.

Number of visits to  $i$  by time  $S_k$  is exactly  $k$ .

Total elapsed time is  $S_k = T_1 + \dots + T_k$  where  $T_i$  are iid.

Fraction of time in state  $i$  by time  $S_k$  is

$$\frac{k}{S_k} \rightarrow \frac{1}{m_{ii}}$$

by SLLN. So if fraction converges to  $\pi_i$  must have

$$\pi_i = \frac{1}{m_{ii}}.$$

## Summary of Theoretical Results:

For an irreducible aperiodic positive recurrent Markov Chain:

1.  $\mathbf{P}^n$  converges to a stochastic matrix  $\mathbf{P}^\infty$ .
2. Each row of  $\mathbf{P}^\infty$  is  $\pi$  the unique stationary initial distribution.
3. The stationary initial distribution is given by

$$\pi_i = 1/m_i$$

where  $m_i$  is the mean return time to state  $i$  from state  $i$ .

If the state space is finite an irreducible chain is positive recurrent.

## Ergodic Theorem

Notice slight of hand: I showed

$$\frac{\mathbb{E} \left\{ \sum_{i=0}^n \mathbf{1}(X_i = k) \right\}}{n} \rightarrow \pi_k$$

but claimed

$$\frac{\sum_{i=0}^n \mathbf{1}(X_i = k)}{n} \rightarrow \pi_k$$

almost surely which is also true. This is a step in the proof of the ergodic theorem. For an irreducible positive recurrent Markov chain and any  $f$  on  $S$  such that  $\mathbb{E}^\pi(f(X_0)) < \infty$ :

$$\frac{\sum_0^n f(X_i)}{n} \rightarrow \sum \pi_j f(j)$$

almost surely. The limit works in other senses, too. You also get

$$\frac{\sum_0^n f(X_i, \dots, X_{i+k})}{n} \rightarrow \mathbb{E}^\pi \{f(X_0, \dots, X_k)\}$$

E.g. fraction of transitions from  $i$  to  $j$  goes to

$$\pi_i \mathbf{P}_{ij}$$

For an irreducible positive recurrent chain of period  $d$ :

1.  $\mathbf{P}^d$  has  $d$  communicating classes each of which forms an irreducible aperiodic positive recurrent chain.
2.  $(\mathbf{P}^{n+1} + \dots + \mathbf{P}^{n+d})/d$  has a limit  $\mathbf{P}^\infty$ .
3. Each row of  $\mathbf{P}^\infty$  is  $\pi$  the unique stationary initial distribution.
4. Stationary initial distribution places probability  $1/d$  on each of the communicating classes in 1.

For an irreducible null recurrent chain:

1.  $\mathbf{P}^n$  converges to 0 (pointwise).
2. there is no stationary initial distribution.

For an irreducible transient chain:

1.  $\mathbf{P}^n$  converges to 0 (pointwise).
2. there is no stationary initial distribution.

For a chain with more than 1 communicating class:

1. If  $\mathcal{C}$  is a recurrent class the submatrix  $\mathbf{P}_{\mathcal{C}}$  of  $\mathbf{P}$  made by picking out rows  $i$  and columns  $j$  for which  $i, j \in \mathcal{C}$  is a stochastic matrix. The corresponding entries in  $\mathbf{P}^n$  are just  $(\mathbf{P}_{\mathcal{C}})^n$  so you can apply the conclusions above.
2. For any transient or null recurrent class the corresponding columns in  $\mathbf{P}^n$  converge to 0.
3. If there are multiple positive recurrent communicating classes then the stationary initial distribution is not unique.