

2 Polyhedra

Preliminaries

By default, we treat elements of \mathbb{R}^m as column vectors and if $x, y \in \mathbb{R}^m$ we let $x \cdot y = x^\top y$.

Polyhedron: A *polyhedron* $P \subseteq \mathbb{R}^n$ is any set of the form $\{x \in \mathbb{R}^n : Ax \leq b\}$ where A is an $m \times n$ matrix and $b \in \mathbb{R}^m$. If $c \in \mathbb{R}^n$ and $x \cdot c \geq t$ for every $x \in P$ then $\{x \in P : x \cdot c = t\}$ is a *face* of P . A *vertex* of P is a face which consists of a single point.

Observation 2.1 *If x is a vertex of the polyhedron P , then there exists $w \in \mathbb{Z}^n$ and $\lambda \in \mathbb{R}$ so that $\{x\} = \{y \in P : y^\top w \geq \lambda\}$.*

Proof: Choose a vector u and a number μ so that $\{x\} = \{y \in P : y^\top u \geq \mu\}$. It follows from the density of the rationals that we may assume u is rational. Let N be the least common multiple of the denominators which appear in the entries of the vector u . Now setting $w = Nu$ and $\lambda = N\mu$ we have that $\{x\} = \{y \in P : y^\top w \geq \lambda\}$ as required. \square

Pointed: We say that P is *pointed* if every minimal face is a vertex. Note that if $P \subseteq \mathbb{R}_+^m$ then P must be pointed.

Integral: A point $x \in \mathbb{R}^m$ is *integral* if every coordinate is an integer. A polyhedron $P \subseteq \mathbb{R}^m$ is *integral* if every minimal face contains an integral point. Note that a pointed polyhedron is integral if and only if every vertex is integral.

Convex Hull: The *convex hull* of a set $X \subseteq \mathbb{R}^m$ is the unique minimal convex set which includes X . Note that the *convex hull* of the rows of A is precisely

$$\{y^\top A : y \in \mathbb{R}_+^m \text{ and } y \cdot 1 = 1\}.$$

Up & Down-Monotone: We say that a set $X \subseteq \mathbb{R}^m$ is *up-monotone* if whenever $x \in X$ and $y \geq x$ we have $y \in X$. Similarly, we say that X is *down monotone in \mathbb{R}_+^m* if whenever $x \in X$ and $0 \leq y \leq x$ we have $y \in X$.

Up & Down-Hull: If $X \subseteq \mathbb{R}_+^m$ the *up-hull* of X is the unique minimal up-monotone set which includes X . Similarly, the *down-hull of X in \mathbb{R}_+^m* is the unique minimal down-monotone set in \mathbb{R}_+^m which includes X .

Essential Vertex: If $P \subseteq \mathbb{R}_+^m$ is down monotone in \mathbb{R}_+^m then a vertex x of P is *essential* if x is not contained in the down-hull of $P \setminus x$.

Observation 2.2 Let $P \subseteq \mathbb{R}_+^m$ be a polyhedron. Then

- (i) If P is bounded, then P is the convex hull of its vertices.
- (ii) If P is up-monotone, then P is the up-hull of its vertices.
- (iii) If P is down-monotone in \mathbb{R}_+^m , then P is the down-hull of its essential vertices.

Up & Down: Let \mathcal{C} be a clutter and let $M = M(\mathcal{C})$. We define the polyhedra $Up(\mathcal{C})$ to be the up-hull of the incidence vectors of the edges of \mathcal{C} and $Down(\mathcal{C})$ to be the down-hull of the incidence vectors of the edges of \mathcal{C} in \mathbb{R}_+^V . Thus, we have:

$$Up(\mathcal{C}) = \{x \in \mathbb{R}_+^m : x^\top \geq y^\top A \text{ for some } y \in \mathbb{R}_+^m \text{ with } y \cdot 1 = 1\}$$

$$Down(\mathcal{C}) = \{x \in \mathbb{R}_+^m : x^\top \leq y^\top A \text{ for some } y \in \mathbb{R}_+^m \text{ with } y \cdot 1 = 1\}$$

Proposition 2.3 Let $\mathcal{C} = (V, E)$ be a clutter and let $X \subseteq \mathbb{R}_+^V$ be the set of incidence vectors of its edges. Then X is the set of vertices of $Up(\mathcal{C})$ and the set of essential vertices of $Down(\mathcal{C})$.

Proof: Homework.

The Packing Polyhedron and Ideal Clutters

Blocking Polyhedra: If $P \subseteq \mathbb{R}_+^S$ is an up-monotone polyhedron, its *blocker* is

$$b(P) = \{y \in \mathbb{R}_+^S : x^\top y \geq 1 \text{ for every } x \in P\}$$

Lemma 2.4 Let $P \subseteq \mathbb{R}_+^S$ be an up-monotone polyhedron.

- (i) if $z \notin P$ there exist $y \in \mathbb{R}_+^S$ and $\lambda \in \mathbb{R}_+$ s.t. $y^\top z < \lambda$ and $y^\top x \geq \lambda$ for every $x \in P$.
- (ii) if $x \in P$ is a vertex, there exist $y \in \mathbb{R}_+^S$ and $\lambda \in \mathbb{R}_+$ s.t. $\{x\} = \{z \in P : y^\top z \leq \lambda\}$.

Proof: (i): Since P is closed and convex and $z \notin P$, we may choose a hyperplane H which separates z and P . Choose a normal vector y for H and suppose (for a contradiction) that there are indices i, j with $y_i > 0$ and $y_j < 0$. Choose a point $x \in P$ and suppose that $y^\top x = t$. If we start at x and then increase the coordinate x_i (x_j) we stay in P , and it follows that there are points in P whose dot product with y is equal to t' for any $t' \geq t$ ($t' \leq t$). But this contradicts the assumption that H is disjoint from P . So, by possibly replacing y by $-y$ we may assume that $y \geq 0$. Choose $\lambda \neq 0$ so that $H = \{x \in \mathbb{R}_+^S : x^\top y = \lambda\}$. It now follows from the fact that y, z, P all lie in the nonnegative orthant that $\lambda > 0$. This completes the proof. The proof of (ii) is similar except that H is chosen to be a hyperplane which intersects P only at x . \square

Proposition 2.5 $b(b(P)) = P$ for every up-monotone polyhedron $P \subseteq \mathbb{R}_+^S$.

Proof: Every $x \in P$ satisfies $x \geq 0$ and $x^\top y \geq 1$ for every $y \in b(P)$ by definition. It follows from this that $x \in b(b(P))$, so $P \subseteq b(b(P))$.

Suppose (for a contradiction) that $b(b(P)) \not\subseteq P$, and choose $z \in b(b(P)) \setminus P$. Apply the lemma to choose a vector $y \in \mathbb{R}_+^S$ and $\lambda \in \mathbb{R}_+$. Now, replacing y by the vector $\frac{1}{\lambda}y$ gives us $y^\top z < 1$ and $y^\top x \geq 1$ for all $x \in P$. It follows from the latter that $y \in b(P)$, but then the former condition contradicts $z \in b(b(P))$. \square

Let \mathcal{C} be a clutter with $V = V(\mathcal{C})$, set $A = M(\mathcal{C})$ and $B = M(b(\mathcal{C}))$.

Fractional Transversal: A *fractional transversal* of \mathcal{C} is a vector $x \in \mathbb{R}_+^V$ with the property that $Ax \geq 1$. So, in other words, a fractional transversal is a weighting of the vertices with the property that every edge gets a total weight of ≥ 1 . Note that a $0, 1$ vector is a fractional transversal if and only if it is the incidence vector of a transversal.

Packing Polyhedron: We define the *packing polyhedron* as follows:

$$\text{Pack}(\mathcal{C}) = \{x \in \mathbb{R}_+^V : Ax \geq 1\}$$

So, in words, the packing polyhedron is the set of all fractional transversals.

Ideal: The clutter \mathcal{C} is *ideal* if $\text{Pack}(\mathcal{C})$ is integral.

Observation 2.6

- (i) $Up(b(\mathcal{C})) \subseteq Pack(\mathcal{C})$
- (ii) $Up(b(\mathcal{C})) = Pack(\mathcal{C})$ if and only if \mathcal{C} is ideal.

Proof: Part (i) follows immediately from the observation that each incidence vector of an edge in $b(\mathcal{C})$ is contained in $Pack(\mathcal{C})$. Part (ii) follows from the observation that every integral point in $Pack(\mathcal{C})$ is also contained in $Up(b(\mathcal{C}))$. \square

Observation 2.7 $b(Up(\mathcal{C})) = Pack(\mathcal{C})$

Proof:

$$\begin{aligned} b(Up(\mathcal{C})) &= \{y \in \mathbb{R}_+^V : y^\top x \geq 1 \text{ for every } x \in Up(\mathcal{C})\} \\ &= \{y \in \mathbb{R}_+^V : y^\top x \geq 1 \text{ whenever } x \text{ is the incidence vector of an edge of } \mathcal{C}\} \\ &= Pack(\mathcal{C}) \quad \square \end{aligned}$$

For the remaining two results in this section, let us name the following polyhedra:

$$\begin{aligned} P &= Pack(\mathcal{C}) \\ Q &= Pack(b(\mathcal{C})) \\ P_I &= Up(b(\mathcal{C})) \\ Q_I &= Up(\mathcal{C}) \end{aligned}$$

So, by the previous observation $b(P) = Q_I$ and $b(Q) = P_I$. Observation 2.6 shows that $P_I \subseteq P$ and $Q_I \subseteq Q$ and further that \mathcal{C} is ideal if and only if $P = P_I$ and $b(\mathcal{C})$ is ideal if and only if $Q = Q_I$.

Theorem 2.8 (Lehman) \mathcal{C} is ideal if and only if $b(\mathcal{C})$ is ideal.

Proof:

$$\begin{aligned} \mathcal{C} \text{ is ideal} &\Leftrightarrow P = P_I \\ &\Leftrightarrow b(P) = b(P_I) \\ &\Leftrightarrow Q_I = Q \\ &\Leftrightarrow b(\mathcal{C}) \text{ is ideal.} \end{aligned}$$

\square

Theorem 2.9 (Lehman's width-length inequality) *The following are equivalent*

- (i) \mathcal{C} and $b(\mathcal{C})$ are ideal.
- (ii) $(\min_{e \in E(\mathcal{C})} w(e)) (\min_{f \in E(b(\mathcal{C}))} \ell(f)) \leq w^\top \ell$ for every $w, \ell \in \mathbb{R}_+^V$.

(here we let $x(S) = \sum_{s \in S} x(s)$ whenever $x \in \mathbb{R}^V$ and $S \subseteq V$).

Proof: The result is trivial whenever $\min_{e \in E(\mathcal{C})} w(e) = 0$ or $\min_{f \in E(b(\mathcal{C}))} \ell(f) = 0$. Furthermore, it is invariant under scaling w or ℓ by a positive constant, so we may assume that $\min_{e \in E(\mathcal{C})} w(e) = 1 = \min_{f \in E(b(\mathcal{C}))} \ell(f)$. We now have the following (the first identity is a consequence of these assumptions).

$$\begin{aligned}
 (ii) &\Leftrightarrow w^\top \ell \geq 1 \text{ whenever } w \in P \text{ and } \ell \in Q \\
 &\Leftrightarrow P \subseteq b(Q) \text{ and } Q \subseteq b(P) \\
 &\Leftrightarrow P \subseteq P_I \text{ and } Q \subseteq Q_I \\
 &\Leftrightarrow P = P_I \text{ and } Q = Q_I \\
 &\Leftrightarrow \mathcal{C} \text{ and } b(\mathcal{C}) \text{ are ideal} \\
 &\Leftrightarrow (i) \quad \square
 \end{aligned}$$

The Covering Polyhedron and Perfect Clutters

Antiblocker: If $P \subseteq \mathbb{R}_+^S$ is a down-monotone polyhedron, its *antiblocker* is

$$a(P) = \{y \in \mathbb{R}_+^S : x^\top y \leq 1 \text{ for every } x \in P\}$$

Lemma 2.10 *Let $P \subseteq \mathbb{R}_+^S$ be a down-monotone polyhedron.*

- (i) *if $z \notin P$ there exist $y \in \mathbb{R}_+^S$ and $\lambda \in \mathbb{R}_+$ s.t. $y^\top z > \lambda$ and $y^\top x \leq \lambda$ for every $x \in P$.*
- (ii) *if $x \in P$ is an essential vertex, there exist $y \in \mathbb{R}_+^S$ and $\lambda \in \mathbb{R}_+$ s.t. $\{x\} = \{z \in P : y^\top z \geq \lambda\}$.*

Proof: Homework

Theorem 2.11 $a(a(P)) = P$ for every down-monotone polyhedron $P \subseteq \mathbb{R}_+^S$.

Proof: Homework

Let \mathcal{C} be a clutter with $V = V(\mathcal{C})$ and set $A = M(\mathcal{C})$.

Fractional Independent Set: A *fractional independent set* of \mathcal{C} is a vector $x \in \mathbb{R}_+^V$ with the property that $Ax \leq 1$. So, in other words, a fractional independent set is a weighting of the vertices with the property that every edge gets a total weight of ≤ 1 . Note that a $0,1$ vector is a fractional independent set if and only if it is the incidence vector of an independent set.

Covering Polyhedron: We define the *covering polyhedron* as follows:

$$Cov(\mathcal{C}) = \{x \in \mathbb{R}_+^V : Ax \leq 1\}$$

So, in words, the covering polyhedron is the set of all fractional independent sets.

Perfect: The clutter \mathcal{C} is *perfect* if $Cov(\mathcal{C})$ is integral.

Observation 2.12

- (i) $Down(a(\mathcal{C})) \subseteq Cov(\mathcal{C})$
- (ii) $Down(a(\mathcal{C})) = Cov(\mathcal{C})$ if and only if \mathcal{C} is perfect.

Proof: Part (i) follows immediately from the observation that each incidence vector of an edge in $a(\mathcal{C})$ is contained in $Cov(\mathcal{C})$. Part (ii) follows from the observation that every integral point in $Cov(\mathcal{C})$ is also contained in $Down(a(\mathcal{C}))$. \square

Observation 2.13 $a(Down(\mathcal{C})) = Cov(\mathcal{C})$.

Proof:

$$\begin{aligned} a(Down(\mathcal{C})) &= \{y \in \mathbb{R}_+^V : y^\top x \leq 1 \text{ for every } x \in Down(\mathcal{C})\} \\ &= \{y \in \mathbb{R}_+^V : y^\top x \leq 1 \text{ whenever } x \text{ is the incidence vector of an edge in } \mathcal{C}\} \\ &= Cov(\mathcal{C}) \quad \square \end{aligned}$$

Define the following polyhedra:

$$P = \text{Cov}(\mathcal{C})$$

$$Q = \text{Cov}(a(\mathcal{C}))$$

$$R_I = \text{Down}(\mathcal{C})$$

$$P_I = \text{Down}(a(\mathcal{C}))$$

$$Q_I = \text{Down}(a(a(\mathcal{C})))$$

So, by the previous observation, $a(P) = R_I$ and $a(Q) = P_I$. It is immediate from the definitions that $R_I \subseteq Q_I$ and Observation 2.12 shows that we have $P_I \subseteq P$ and $R_I \subseteq Q_I \subseteq Q$. Further, this observation shows that \mathcal{C} is perfect if and only if $P = P_I$ and $a(\mathcal{C})$ is perfect if and only if $Q = Q_I$.

Theorem 2.14 *If \mathcal{C} is perfect, then $a(\mathcal{C})$ is perfect.*

Proof: By our observations

$$\mathcal{C} \text{ perfect} \Rightarrow P = P_I \Rightarrow a(P) = a(P_I) \Rightarrow R_I = Q \Rightarrow Q_I = Q \Rightarrow a(\mathcal{C}) \text{ perfect.} \quad \square$$

Theorem 2.15 (Chvátal) *If \mathcal{C} is perfect, then*

- (i) $a(a(\mathcal{C})) = \mathcal{C}$
- (ii) $\mathcal{C} = \text{CN}(G)$ for a graph G .

Proof: For (i) we have

$$\mathcal{C} \text{ perfect} \Rightarrow P = P_I \Rightarrow a(P) = a(P_I) \Rightarrow R_I = Q \Rightarrow R_I = Q_I \Rightarrow \mathcal{C} = a(a(\mathcal{C})).$$

The proof of (ii) follows from this and an earlier exercise. \square

Theorem 2.16 *The following are equivalent.*

- (i) \mathcal{C} is perfect.
- (ii) $(\max_{e \in E(\mathcal{C})} w(e)) (\max_{f \in E(a(\mathcal{C}))} \ell(f)) \geq w^\top \ell$ for every $w, \ell \in \mathbb{R}_+^V$.

Proof: Homework.