

Planar Graphs on Nonplanar Surfaces

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It is shown that embeddings of planar graphs in arbitrary surfaces other than the 2-sphere have a special structure. It turns out that these embeddings can be described in terms of noncontractible curves in the surface, meeting the graph in at most two points (which may taken to be vertices of the graph). The close connection between the homology group of the surface and the planar graph embeddings is perhaps the most interesting aspect of this study. Some important consequences follow from these results. For example, any two embeddings of a planar graph in the same surface can be obtained from each other by means of simple local reembeddings very similar to Whitney's switchings. © 1996 Academic Press, Inc.

1. INTRODUCTION

Let Π be a (2-cell) embedding of a graph G into a *nonplanar surface* S , i.e., a closed surface distinct from the 2-sphere. Then we define the *face-width* $\text{fw}(\Pi)$ (also called the *representativity*) of the embedding Π as the smallest number of (closed) faces of G in S whose union contains a noncontractible curve.

One of the first results about the face-width, due to Robertson and Vitray [5] (cf. also [6]), considers the face-width of nonplanar embeddings of planar graphs. They proved that a planar graph embedded in a nonplanar surface has face-width at most two. Our main concern is to strengthen this result to obtain, essentially, a simple description of the

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structure of planar graphs embedded in nonplanar surfaces. Indirectly, this characterizes embedded graphs whose dual graphs are planar. It turns out that these embeddings can be described in terms of noncontractible curves in the surface, meeting the graph in at most two points (which may be taken to be vertices of the graph). We prove, in particular, that such curves pass through all vertices whose local rotation differs from that in a planar embedding of the graph; see Theorems 4.4 and 4.9. The close connection between the first homology group of the surface and the planar graph embeddings is perhaps the most interesting aspect of this study. Several important consequences follow from these results. For example, any two embeddings of a 2-connected planar graph in the same surface can be obtained from each other by using very simple elementary changes, called generalized Whitney switchings (see Theorem 7.1 and Corollary 7.3). This generalizes Whitney's Theorem [7] stating that any two embeddings of a 2-connected graph in the plane can be obtained from each other using a sequence of Whitney's 2-switchings.

The structure of embeddings of planar graphs in the projective plane is analysed in more details in [4]. In this paper we give such a description for the case of the torus and the Klein bottle under some additional restrictions on the embedding; see Corollary 6.3. It turns out that planar graphs embed in the torus in a particularly simple way. The local switches necessary for the generalized Whitney's theorem are explicitly described.

In the rest of the paper we assume that G is a 2-connected planar graph. Most of the results can easily be extended to the 1-connected case but sometimes some technical conditions should be added. We shall repeat this assumption only in the statements of the main results to warn the readers that have skipped reading this part.

2. EMBEDDINGS

Let G be a connected graph. 2-cell embeddings of G in closed surfaces can be described in a purely combinatorial way by specifying:

(1) A *rotation system* $\pi = (\pi_v; v \in V(G))$; for each vertex v of G we have a cyclic permutation π_v of edges incident with v , representing their circular order around v on the surface. The cyclic sequence $e, \pi_v(e), \pi_v^2(e), \pi_v^3(e), \dots$ is called *Π -clockwise ordering* around v .

(2) A *signature* $\lambda: E(G) \rightarrow \{-1, 1\}$. Suppose that $e = uv$. Following the edge e on the surface, we see if the local rotations π_v and π_u are chosen consistently or not. If yes, then we have $\lambda(e) = 1$, otherwise we have $\lambda(e) = -1$.

The reader is referred to [1] for more details. We will use this description as a definition: An *embedding* of a connected graph G is a pair $\Pi = (\pi, \lambda)$ where π is a rotation system and λ is a signature. Having an embedding Π of G , we say that G is Π -*embedded*. The embedding Π is *nonorientable* if there is a cycle with an odd number of edges e having $\lambda(e) = -1$. Such a cycle is Π -*onesided*. Other cycles are Π -*twosided*. We define Π -*facial walks* as closed walks in the graph that correspond to face boundaries of the corresponding topological embedding. If W is a walk, any subwalk of W is a *segment* of W . If $e = uv \in E(G)$, the pair $\{e, \pi_v(e)\}$ forms a Π -*angle*. A pair of edges that forms a Π -angle is Π -*consecutive*. We define the *genus* and the *Euler characteristic* $\chi(\Pi)$ of Π as the genus and the Euler characteristic of the corresponding topological embedding, respectively.

If $\Pi = (\pi, \lambda)$ is an embedding of a graph G and H is a subgraph of G , then the *restriction* of Π to H is the embedding of H whose rotation system is obtained from π by ignoring all edges in $E(G) \setminus E(H)$ and whose signature is the restriction of λ to $E(H)$.

If G is a Π -embedded graph and C is a Π -twosided cycle of G , then we define the *left graph* and the *right graph* of C as follows. Select a vertex $v \in V(C)$, and let e and e' be the edges of C incident with v . If $e' = \pi_v^k(e)$, then all edges $e, \pi_v(e), \pi_v^2(e), \dots, \pi_v^k(e)$ are said to be *on the left side* of C . Then we traverse C and determine left edges at each vertex of C in the same way as at v . However, after traversing an edge f of C with $\lambda(f) = -1$, we change clockwise orientation to anticlockwise, and vice versa. In particular, traversing the edge $e' = vu$ from v to u , the left edges at u are $e', \pi_u(e'), \pi_u^2(e'), \dots, \pi_u^l(e')$ (where $\pi_u^l(e') \in E(C)$) if we have the clockwise orientation. Having the anticlockwise orientation, the left edges are $\pi_u^l(e'), \pi_u^{l+1}(e'), \dots, e'$. Since C is Π -twosided, the clockwise orientation is the same as at the beginning when we come back to the initial vertex v after traversing the entire cycle C . An arbitrary edge e (possibly not incident with C) is also said to be *on the left side* of C if one of its ends is connected by a path in $G - C$ to an end of an edge on the left side of C (and incident with C). Now the left graph $G_l = G_l(\Pi, C)$ is defined as the graph induced by all edges on the left side of C . The right graph $G_r = G_r(\Pi, C)$ is defined analogously. Note that $C \subseteq G_l \cap G_r$.

Let C be a Π -twosided cycle and G_l and G_r its left and right graph. If $G_l \cap G_r = C$, then C is said to be Π -*bounding*. A Π -bounding cycle C is Π -*contractible* if the embedding of G restricted to $G_l(\Pi, C)$ (or to $G_r(\Pi, C)$) is an embedding of genus 0. In this case, $G_l(\Pi, C)$ (or $G_r(\Pi, C)$, respectively) is called the Π -*interior* of C , and the rest of G , $G_r(\Pi, C) \setminus E(C)$ ($G_l(\Pi, C) \setminus E(C)$, respectively) is the Π -*exterior* of C . By definition, Π -onesided cycles are Π -nonbounding.

3. PATCHES

Let G be a Π -embedded 2-connected planar graph. Suppose that G contains a Π -contractible cycle C such that only two vertices u, v of C have incident edges that are in the Π -exterior of C . The replacement of the Π -interior of C by the edge uv is called a *2-reduction*. Note that G admits an embedding Π' in the 2-sphere such that the Π' -interior of the cycle C is the same as the Π -interior of C and hence the same 2-reduction can be performed to Π' . If v is a vertex of degree 2 in G and the neighbors of v are u and w , then the replacement of v and its incident edges by the edge uw is also called a *2-reduction*.

Let us now suppose that no 2-reductions are possible. Let Π' be an embedding of G in the 2-sphere. Then we define $\text{CS}_0(\Pi, \Pi')$ to be the set of all Π -facial walks, that are also Π' -facial, together with all paths P in G such that P is simultaneously a segment of a Π -facial walk and a segment of a Π' -facial walk. Since G is 2-connected, Π' -facial walks are cycles and hence $\text{CS}_0(\Pi, \Pi')$ contains only paths and cycles. Denote by $\text{CS}(\Pi, \Pi')$ the subset of $\text{CS}_0(\Pi, \Pi')$ consisting of all cycles and paths that are not contained in another element of $\text{CS}_0(\Pi, \Pi')$. If W is a Π -facial walk from $\text{CS}(\Pi, \Pi')$, then we replace W by a graph \tilde{W} as shown in Fig. 1. Similarly, if $W \in \text{CS}(\Pi, \Pi')$ is a maximal common segment of a Π -facial walk and a Π' -facial walk, then we replace W by \tilde{W} as shown in Fig. 2. (As a special case, when W is just a path consisting of a single edge e of G , this operation is just a subdivision of e obtained by inserting five vertices of degree 2 on e .) When we do such replacements for all cycles and paths from $\text{CS}(\Pi, \Pi')$, we obtain a 2-connected planar graph \bar{G} containing the (subdivided) graph G as a subgraph. The embeddings Π and Π' can be naturally extended to embeddings of \bar{G} .

In general, starting with G we first perform all possible 2-reductions and then we construct, from the obtained graph G' , the graph \bar{G}' as described above. If the graph \bar{G}' contains a path whose interior vertices are all of degree 2, then we replace such a path by a single edge. After all such

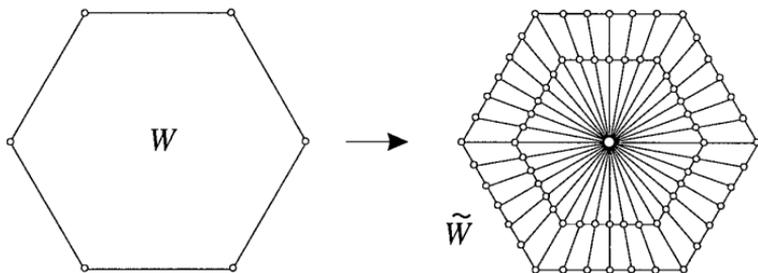


FIG. 1. Filling up common faces.

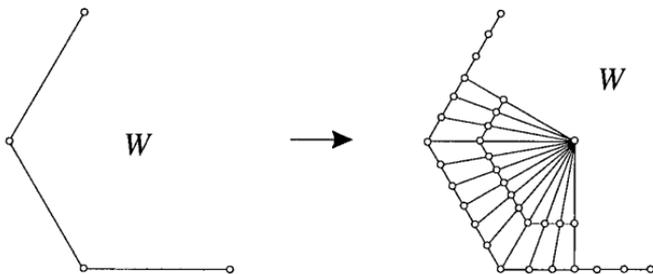


FIG. 2. Filling up maximal common facial segments.

replacements, if one half of an original edge e of G' becomes just an edge \bar{e} in the resulting graph (and the other half of e contains vertices of degrees 3 and 4), then we contract \bar{e} to a point so that the “middle vertex” of e is identified with an end of e . After doing all such changes, we obtain an embedded graph \tilde{G} that is called a *patch extension* of the Π -embedded graph G . The patch extension \tilde{G} contains a subdivision of G as a subgraph.

It is clear that the embeddings Π' and Π can be extended to embeddings of \tilde{G} in the same surfaces such that all triangles and quadrangles shown in Figs. 1 and 2, respectively, are facial. In particular, \tilde{G} is a planar graph. It is also easy to see that if G is 2-connected (3-connected, respectively), then so is \tilde{G} .

Let \tilde{G} be the patch extension of a 2-connected Π -embedded planar graph G . Denote by $\tilde{\Pi}$ the corresponding embedding of \tilde{G} . The $\tilde{\Pi}$ -facial walks that are not facial walks of the plane embedding of \tilde{G} are the *patch facial walks* and the corresponding faces are the *patch faces*. Vertices of \tilde{G} that belong to two or more patch facial walks are *patch vertices*. Segments of patch facial walks joining patch vertices are also segments of facial walks of \tilde{G} embedded in the plane. They are called *patch edges*. Two $\tilde{\Pi}$ -consecutive patch edges incident with the same patch vertex v form a *patch angle* at v . The *patch degree* of a patch vertex v is the number of patch angles at v . Let us remark that each patch vertex v is either a vertex of G or the middle vertex of a subdivided edge of G . In the latter case, the patch degree of v is equal to two.

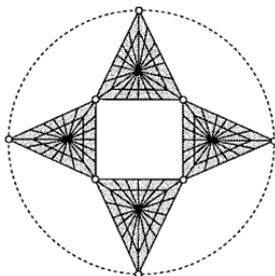


FIG. 3. Patch structure of the octahedron in the projective plane.

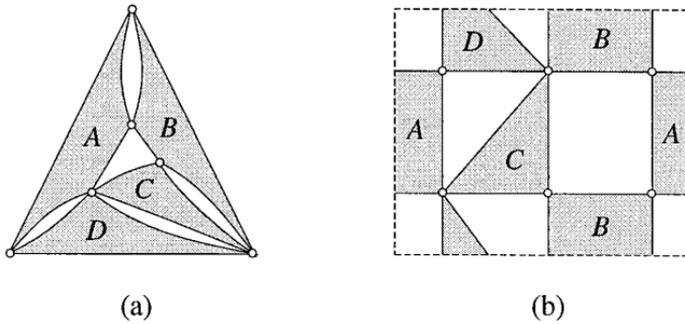


FIG. 4. Patch structure of an embedding in the torus.

Edges e and f of \tilde{G} are *similar* if they both lie on the same patch edge or if they both lie on the same $\tilde{\Pi}$ -facial walk that is not a patch facial walk. The smallest equivalence relation on $E(\tilde{G})$ containing the similarity relation determines a partition of edges of \tilde{G} into subgraphs of \tilde{G} . They are called *patches* of G (with respect to the embedding Π). It is convenient to consider the patches as being subsets of the surface of the embedding Π consisting of the corresponding subgraph of \tilde{G} together with all non-patch $\tilde{\Pi}$ -faces that they contain. As such, distinct patches have disjoint interiors and they meet only in common patch vertices. They partition the complement of the interiors of patch faces in S . The interior (in S) of every patch is homeomorphic to an open disk in the plane with $p \geq 0$ holes.

Geometrically we will represent patches as shaded areas on the surface and will refer to the combinatorial structure of patches of an embedding Π as the *patch structure* of Π . For example, Fig. 3 shows the patch extension of the octahedron embedded in the projective plane. The shaded areas in Fig. 3 are the patches. Another example of a planar graph (Fig. 4(a)) and its embedding in the torus (Fig. 4(b)) shows a more complicated patch structure.

4. PATCH ANGLES AND 2-CURVES

Suppose that Π is an embedding of a 2-connected planar graph G in a nonplanar surface S . Let k be an integer. A simple closed curve γ in S is a *k-curve* if it satisfies the following conditions:

- (C1) γ intersects the graph G in exactly k points and all these points are patch vertices of G .
- (C2) γ uses only patch vertices and patch faces of G .
- (C3) γ is Π -noncontractible.

In (C3) we have used the term Π -noncontractible for a curve γ that is not a cycle of G . However, if we consider the segments of γ in the Π -faces of G as edges between the corresponding patch vertices, the embedding Π can be naturally extended to the union of G and these edges. Then γ corresponds to a cycle in the new graph, and contractibility of γ refers to contractibility of this cycle. Similarly, we can use other concepts that were introduced for cycles in embedded graphs also for k -curves.

We will be interested mainly in 1-curves and 2-curves. Clearly, these curves coincide in G and in the patch extension \tilde{G} of G . In the case of the patch extension \tilde{G} , condition (C2) is automatically satisfied for any curve γ for which (C1) and (C3) hold. Therefore it is more convenient to work with \tilde{G} instead of the original graph G . Since \tilde{G} has the same patch structure as G , we can assume from now on that our planar graph G is the patch extension graph of some planar graph. In particular, we assume that no 2-reductions are possible.

Two 1-curves are *equivalent* if they use the same patch face Φ , the same patch vertex v and the same pair of patch angles of Φ at v . Two 2-curves are *equivalent* if they pass through the same pair of patch faces and patch vertices and use the same patch angles. We will distinguish 1-curves and 2-curves only up to equivalence.

By a *1/2-curve* we shall refer to 1-curves and 2-curves. In this section we shall prove that there are 1/2-curves “everywhere in the surface” except inside the patches where the embedding of the graph G locally matches planar embeddings of G .

LEMMA 4.1. *Let $\{\alpha, \beta\}$ be a patch angle at a patch vertex v in a patch facial walk Φ . If the patch edge α does not occur on Φ twice, then there is a cycle $C(\alpha, \beta) \subseteq \Phi$ that contains α . If v appears on Φ just once, then $C(\alpha, \beta)$ also contains β .*

Proof. Those patch edges of Φ that occur on Φ just once form an Eulerian graph. Hence the claim. ■

Disjoint pairs $\{\alpha, \beta\}$ and $\{\gamma, \delta\}$ of (patch) edges incident to the same patch vertex v Π' -interlace if α and β (and hence also γ and δ) are not consecutive under the rotation system of Π' restricted to $\{\alpha, \beta, \gamma, \delta\}$.

LEMMA 4.2. *Suppose that $\{\alpha, \beta\}$ and $\{\gamma, \delta\}$ are Π' -interlacing patch angles at a patch vertex v . Then either there is a 1-curve through $\{\alpha, \beta\}$ or through $\{\gamma, \delta\}$, or there is a 2-curve through v that uses both patch angles $\{\alpha, \beta\}$ and $\{\gamma, \delta\}$ at v .*

Proof. If v appears more than once on the patch facial walk containing the angle $\{\alpha, \beta\}$, then there is a simple closed curve ψ in the corresponding

patch face through the angle $\{\alpha, \beta\}$. Since G is 2-connected, ψ is Π -nonbounding and hence a 1-curve. Similarly for the angle $\{\gamma, \delta\}$. Otherwise, let $C(\alpha, \beta)$ and $C(\gamma, \delta)$ be cycles from Lemma 4.1. In the plane, these two cycles cross at v , and hence they have another point τ in common. We may assume that τ is a patch vertex. Let ψ be a simple closed curve through the angles $\{\alpha, \beta\}$ and $\{\gamma, \delta\}$ that intersects G only at v and τ . If ψ is contractible, then it bounds a disk containing more than just one edge in its interior. Because of our initial 2-reductions, this is not possible, and hence ψ is a 2-curve. ■

Let γ and γ' be simple closed curves on a surface S that intersect in finitely many points. Suppose that $z \in \gamma \cap \gamma'$. Then γ and γ' *cross* at z if z has an open neighborhood U homeomorphic to the plane such that the homeomorphism maps $U \cap \gamma$ onto the x -axis and $U \cap \gamma'$ onto the y -axis in the plane. Otherwise they *touch* at z . The curves are *noncrossing* if they touch at each of their points of intersection.

A similar proof as above yields the following result:

LEMMA 4.3. *Let F and F' be Π -facial walks with an edge $e = uv$ in common. Suppose that $F = \alpha e v \beta \dots$ and that $F' = \delta e v \gamma \dots$. Let H be the subgraph of G consisting of edges $e, \alpha, \beta, \gamma,$ and δ . Suppose that in the rotation system of Π' restricted to H , the local rotation at u is $(\alpha e \delta)$ and the local rotation at v is $(\beta e \gamma)$, and that the signature of e is 1. If there is no 1-curve in F or in F' through v or through u , then there is a 2-curve through F and F' that uses the Π -angles $\{\alpha, e\}$ and $\{\delta, e\}$*

Lemma 4.2 yields an important result about existence of 1/2-curves.

THEOREM 4.4. *Let B be a 2-connected planar graph that is Π -embedded in a nonplanar surface. Then for every patch face Φ and every patch vertex v of Φ , there is either a 1-curve through v or there is a 2-curve through Φ and v .*

Proof. Let v be a patch vertex at a patch angle $\{\alpha, \beta\}$ of Φ . Split the patch edges that are incident to v and different from α and β into two classes, depending in which Π' -subinterval from α to β they are. Since α and β are not Π' -consecutive, there are Π -consecutive patch edges γ, δ that are in different classes. They determine a patch angle that Π' -interlaces with $\{\alpha, \beta\}$, and Lemma 4.2 can be applied. ■

The following result is a simple corollary of Theorem 4.4.

COROLLARY 4.5. *Suppose that G has k patch edges. If $\text{fw}(\Pi) = 2$, then there is a set of at least $k/4$ nonequivalent 2-curves such that each of these*

2-curves passes through a patch angle that is not used by any of the other 2-curves in the set.

Proof. Denote by v_1, \dots, v_s the patch vertices of G . For $i = 1, \dots, s$, let d_i be the number of patch edges incident with v_i . The number of patch angles is at least $\sum_{i=1}^s d_i/2 = k$. Since every 2-curve uses four patch angles, Theorem 4.4 implies that the number of nonequivalent 2-curves satisfying the “minimality” condition of the corollary is at least $k/4$. ■

We cannot argue in the same way as above if $\text{fw}(\Pi) = 1$. The result that we get is slightly weaker.

COROLLARY 4.6. *If G has k patch edges, then there is a set of at least $k/(76 - 48\chi(S))$ nonequivalent 1/2-curves, where $\chi(S)$ denotes the Euler characteristic of Π . Each of these curves passes through a patch angle that is not used by any of the other curves in the set.*

Proof. A patch angle is said to be *nonsimple* if it is used by a 1-curve. It is *bad* if no 1/2-curve uses it. For every patch angle $\{\alpha, \beta\}$ there is an angle $\{\gamma, \delta\}$ that Π' -interlaces with it (cf. the proof of Theorem 4.4.). If an angle $\{\alpha, \beta\}$ at a patch vertex v is bad, the $\{\gamma, \delta\}$ is nonsimple. Let W be the patch facial walk containing the angle $\{\gamma, \delta\}$. Write $W: \gamma_1 Q_1 \delta_2 \gamma_2 Q_2 \delta_3 \dots \gamma_s Q_s \delta_1$ where subwalks Q_1, \dots, Q_s do not contain v and each angle $\{\gamma_i, \delta_i\}$ contains v . Suppose that the two angles $\{\rho, \alpha\}$ and $\{\beta, \sigma\}$ adjacent to the bad angle $\{\alpha, \beta\}$ are not nonsimple. Then $\gamma_i, \delta_i \notin \{\alpha, \beta\}$ for $i = 1, \dots, s$. It is easy to see that the number of angles $\{\gamma_i, \delta_i\}$ that Π' -interlace with $\{\alpha, \beta\}$ is even. In the same way as we proved Lemma 4.2 we see that $C(\alpha, \beta)$ and the cycle $\gamma_i Q_i \delta_{i+1}$ intersect only once, and therefore no consecutive pair γ_i, δ_{i+1} on W Π' -interlaces with $\{\alpha, \beta\}$ ($i = 1, \dots, s$; the index $i + 1$ taken modulo s). These properties imply that the total number of angles at v that Π' -interlace with $\{\alpha, \beta\}$ is even. Hence we have:

(P1) If the two angles $\{\rho, \alpha\}$ and $\{\beta, \sigma\}$ adjacent to the bad angle $\{\alpha, \beta\}$ are not nonsimple, then ρ, σ do not Π' -interlace with α, β .

Let us now consider the patch angles (of Π) at a patch vertex v . Let p be the number of nonsimple ones. Since G is 2-connected, no two 1-curves through v are homotopic. Therefore the number of 1-curves through v is at most $4 - 3\chi(S)$ (see, e.g. [2, Proposition 3.6]). This implies that

$$p \leq 8 - 6\chi(S). \tag{1}$$

Let r be the number of bad angles at v and let q be the number of the remaining angles at v (simple and not bad). Denote by $\{\alpha_1, \beta_1\}, \dots, \{\alpha_r, \beta_r\}$ the bad angles in the order determined by Π .

Suppose that these are s consecutive bad angles, say $\{\alpha_1, \beta_1\}, \dots, \{\alpha_s, \beta_s\}$, such that for $i=1, \dots, s-1$, there is no patch angle between $\{\alpha_i, \beta_i\}$ and $\{\alpha_{i+1}, \beta_{i+1}\}$. Property (P1) implies that for each $i=2, \dots, s-1$ all edges $\alpha_1, \beta_1, \dots, \alpha_{i-1}, \beta_{i-1}, \alpha_{i+1}, \beta_{i+1}, \dots, \alpha_s, \beta_s$ are in the same Π' -part between α_i and β_i . A nonsimple angle $\{\gamma_i, \delta_i\}$ that Π' -overlaps with $\{\alpha_i, \beta_i\}$ has one of its edges, say γ_i , in the other Π' -part between β_i and α_i . Then the edges $\gamma_2, \dots, \gamma_{s-1}$ are all distinct and hence $s-2 \leq 2p$. By (1), $s \leq 18 - 12\chi(S)$. This implies that

$$r \leq (18 - 12\chi(S))(p + q). \quad (2)$$

Now, (2) implies a bound on the number of patch angles at v

$$p + q + r \leq (19 - 12\chi(S))(p + q)$$

in terms of the number of 1-curves and 2-curves through v . Now the same conclusion as used in the proof of Corollary 4.5 yields the bound of the corollary. ■

A more careful application of methods in the above proof yields a better bound than presented above. However, this bound still depends on the genus of S and we do not see a way how to improve it to a similar bound as obtained in Corollary 4.5 in the case of face-width two.

Let P be path in G and let v_1, \dots, v_k be interior vertices of P . The edges incident with v_1, \dots, v_k can be classified as edges on the left side of P (or on the right side of P) in the same way as in the definition of the left (and the right) graph of a Π -twosided cycle. Denote by E'_P the set of all edges of G incident to v_1, \dots, v_k that are distinct from edges on P . By splitting E'_P into the set of edges on the left side and the set of edges on the right side of P , respectively, we obtain the Π -splitting at v_1, \dots, v_k with respect to P . (It may happen that the splitting is not a partition of E'_P .)

LEMMA 4.7. *Let v be a patch vertex and e, f edges (possibly inside patches) incident with v . If no 1-curve contains v and no 2-curve crosses the path $P = evf$ at v , then the Π -splitting and the Π' -splitting at v with respect to P coincide.*

Proof. Suppose that no 1-curve passes through v . Let C_1, C_2 and C'_1, C'_2 be the Π -splitting and the Π' -splitting, respectively. If they are not the same, there is a pair α, β of Π -consecutive patch edges in C_1 (say) such that $\alpha \in C'_1$ and $\beta \in C'_2$. Since e and f Π' -interlace with $\{\alpha, \beta\}$, there is a pair of Π -consecutive edges $\gamma, \delta \in C_2 \cup \{e, f\}$ that Π' -interlace with $\{\alpha, \beta\}$. We are done by applying Lemma 4.2. ■

The condition of Lemma 4.7 that there are no 1-curves through v cannot be omitted. It may happen that the Π and Π' -splittings at v with respect to P do not coincide and that there are neither 1-curves nor 2-curves crossing P at v . An example in the torus is shown in Fig. 5 where P is the vertical path at v .

LEMMA 4.8. *Suppose that $P = eu\sigma v$ is a path in G where σ is a patch edge joining patch vertices u and v . Suppose that no 1-curve passes through u or v and that no 2-curve crosses P at u or v . Then the Π -splitting and the Π' -splitting at u and v with respect to P are the same.*

Proof. By Lemma 4.7, the Π -splitting at u coincides with the Π' -splitting at u with respect to the path $eu\sigma$. Similarly at v . If the patch containing σ is not just an edge, the two pairs of splittings are clearly the same as the splittings with respect to P . Otherwise, we simply apply Lemma 4.3. ■

Above results imply the following.

THEOREM 4.9. *Let G be a Π -embedded 2-connected planar graph. Suppose that C is a Π -nonbounding cycle of G . If no 1-curve passes through a vertex of C , then C contains at least two vertices at which some 2-curve crosses C .*

Proof. Since C is Π' -bounding, the edges of $E(G) \setminus E(C)$ can be classified as interior or exterior edges, depending on whether they are in the Π' -interior or in the Π' -exterior of C , respectively.

Suppose that there is a vertex $u_0 \in V(C)$ and that no 2-curve crosses C at a vertex distinct from u_0 . Consider $C \setminus \{u_0\}$ as an open path P in G . Lemmas 4.7 and 4.8 imply that the Π -splitting and the Π' -splitting at the internal vertices of P are the same. Suppose that C is Π -twosided. Since there is no 1-curve through u_0 , u_0 cannot be the only vertex of C that has an incident edge $e \notin E(C)$ which is on the left side of C . Similarly on the right. Since C is Π -nonbounding, $G_\lambda(\Pi, C) \cap G_r(\Pi, C) \neq C$. This implies

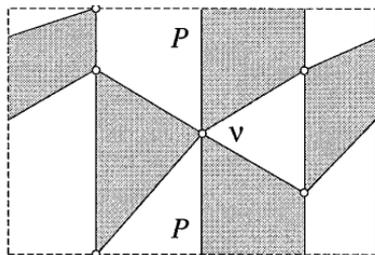


FIG. 5. No 1-curve or 2-curve crosses P .

that there is a Π -facial walk W that uses interior as well as exterior edges. The same conclusion holds also when C is Π -onesided.

The change from an interior to an exterior edge (or vice versa) in the facial walk W can only be achieved at u_0 . Consequently, W meets u_0 at least twice. Since G is 2-connected, the curve in the face of W connecting the two appearances of u_0 on W is Π -noncontractible, hence a 1-curve. A contradiction. ■

It is worth mentioning that for any k there are examples of embeddings Π of planar graphs with Π -noncontractible (but necessarily Π -bounding) cycles that are neither crossed nor touched by an s -curve for $s \leq k$.

A simple corollary of Theorem 4.9 is:

COROLLARY 4.10. *Let C be a Π -nonbounding cycle of G . Then there are patch faces Φ_1, Φ_2 that intersect at a vertex v of C such that Φ_1 is on the left side of C at v , and Φ_2 is on the right side of C at v .*

Proof. Let us first remark that the left and the right side of C are defined locally at each vertex of C also when C is Π -onesided. Suppose now that at each vertex of C , patch faces are only on one side of C . Then there is a cycle \tilde{C} homotopic to C that contains only vertices in the interiors of patches. In particular, no $1/2$ -curve intersects C . Theorem 4.9 now gives a contradiction. ■

Let $\Gamma = \{C_1, \dots, C_k\}$ be a set of cycles of G . If there is a set D of Π -facial walks such that any edge $e \in E(G)$ appears exactly once in facial walks from D exactly when e is contained in an odd number of cycles from Γ , then Γ is Π -bounding. We also say that Γ bounds D . If no nonempty subset of Γ is Π -bounding, then Γ is *homologically independent*. If $\{C_1, C_2\}$ is Π -bounding, then C_1 and C_2 are Π -homologic. The same definitions apply for sets of noncrossing $1/2$ -curves.

Theorems 4.4 and 4.9 show that there are many $1/2$ -curves. Based on these results we formulate the following conjecture that is, in a sense, a claim dual to Theorem 4.9.

Conjecture 4.11. Suppose that G is a 2-connected planar graph that is Π -embedded in an orientable surface of genus g with face-width 2. Then there is a set $\{\gamma_1, \dots, \gamma_g\}$ of pairwise noncrossing homologically independent 2-curves.

A corresponding conjecture for nonorientable surfaces S claims that there is a set $\Gamma = \{\gamma_1, \dots, \gamma_k\}$ of homologically independent 2-curves such that twice the number of twosided 2-curves plus the number of onesided 2-curves in Γ equals the nonorientable genus of S , i.e., $2 - \chi(S)$.

It may be true that even the following stronger property holds: If Γ is any maximal set of pairwise noncrossing 2-curves (i.e., any 2-curve that is equivalent to no curve in Γ crosses some curve from Γ), then Γ contains a set of 2-curves satisfying Conjecture 4.11.

5. CHOICE OF Π'

Let G be a Π -embedded planar graph, and let Π' be an embedding of G in the 2-sphere. The embedding Π' *maximally coincides* with Π if for every embedding Π'' of G in the 2-sphere, $\text{CS}_0(\Pi, \Pi') \subseteq \text{CS}_0(\Pi, \Pi'')$ implies that $\text{CS}_0(\Pi, \Pi') = \text{CS}_0(\Pi, \Pi'')$.

Although the results of the previous sections hold for an arbitrary embedding Π' , an additional assumption that Π' maximally coincides with Π makes some of the results “stronger” since if $\text{CS}_0(\Pi, \Pi'') \subset \text{CS}(\Pi, \Pi'')$, then every patch angle with respect to Π'' is also a patch angle with respect to Π' , but there is a patch angle with respect to Π'' that disappears if we take Π' instead of Π'' .

Given a planar embedding Π'' of G , it is easy to find an embedding Π' that maximally coincides with Π and such that $\text{CS}_0(\Pi, \Pi'') \subseteq \text{CS}_0(\Pi, \Pi')$. The procedure is as follows. Take an arbitrary patch angle of Π'' , subdivide the edges of this angle and connect the inserted vertices by a new edge. If the resulting supergraph of the patch extension \tilde{G} of G is planar, it determines a planar embedding Π_1 of G such that $\text{CS}_0(\Pi, \Pi'') \subset \text{CS}_0(\Pi, \Pi_1)$. By repeating the same with other angles, we eventually stop with an embedding that maximally coincides with Π . Note that this gives a good characterization of embeddings Π' that maximally coincide with Π .

The choice of Π' that maximally coincides with Π has another advantage: our results easily carry over to graphs that are not 2-connected by applying the following proposition.

PROPOSITION 5.1. *Let G be a connected Π -embedded planar graph and let Π' be an embedding of G in the 2-sphere that maximally coincides with Π . Then the patch extension \tilde{G} of G with respect to Π' is 2-connected.*

Proof. Suppose that v is a cutvertex of \tilde{G} . Then v is also a cutvertex of G (viewed as a subgraph of \tilde{G}). There are Π -consecutive edges e_1, e_2 belonging to distinct $\{v\}$ -bridges B_1, B_2 in \tilde{G} . It is easy to see that Π' can be changed so that e_1 and e_2 become Π' -consecutive. The embedding in the plane cannot change the triangular and quadrangular faces of \tilde{G} inside the patches. Hence the new embedding contradicts the assumption that Π' maximally coincides with Π . ■

6. A CHESSBOARD PATTERN

Let γ_1, γ_2 be noncrossing nonbounding twosided 2-curves. Suppose that γ_1 and γ_2 are homologic and let D be the subset of the surface that they bound. For $i=1, 2$ we denote by x_i and y_i the vertices of $\gamma_i \cap G$, and by f_i, g_i the Π -faces used by γ_i . We have the following corollary of Theorem 4.9.

PROPOSITION 6.1. *Suppose that $\text{fw}(\Pi)=2$. Let γ_1 and γ_2 be homologic 2-curves as introduced above. If there is no 2-curve crossing both γ_1 and γ_2 and if every 2-curve in D is equivalent either to γ_1 or to γ_2 , then $\{f_1, g_1\} \cap \{f_2, g_2\} \neq \emptyset$, say $f_1=f_2$, and $\partial f_1 \cap D$ consists of two segments joining x_1, x_2 (say) and y_1, y_2 , respectively.*

Proof. Suppose that $\varphi = \partial f_1 \cap D$ is connected, and so is $\psi = \partial g_1 \cap D$. The closed walk W composed of φ and ψ is homotopic to γ_1 . Thus W contains a Π -nonbounding cycle C . By Theorem 4.9, C is crossed by a 2-curve, say γ . Since γ crosses C , it is equivalent neither to γ_1 nor to γ_2 . Thus γ is not entirely in D , and hence it crosses exactly one of γ_1, γ_2 , say γ_1 . Since C is in D , γ crosses γ_1 (at least) twice. Let x, y and f, g be the vertices and Π -faces, respectively, used by γ . Then one of the following two cases occurs.

Case 1: $x = x_1$ and $y = y_1$. Let γ' be the curve consisting of the segment of γ in D and of the segment of γ_1 in f_1 , and let γ'' be the curve consisting of the segment of γ in D and the segment of γ_1 in g_1 . Then either γ' or γ'' is a 2-curve in D that crosses C . As we have proved above, this is a contradiction.

Case 2: $x, y \neq x_1$ (say). Then γ crosses γ_1 in the interior of f_1 , say, so we may assume that $f=f_1$. If x, y are both in D , the part of γ that is in f can be redrawn inside f so that $\gamma \subseteq D$, a contradiction. (If γ would become contractible after this change, f_1 would contain a 1-curve, a contradiction with $\text{fw}(\Pi)=2$.) So $y \notin D$, say. Since γ crosses C , we have $x \in D$ and $x \neq x_1, y_1$. It follows that $g = g_1$. We now conclude as in Case 1. ■

Let K be a subgraph of G . A K -bridge in G is a subgraph of G which is either an edge of $E(G) \setminus E(K)$ with its endpoints in K , or it is a connected component of $G - V(K)$ together with all edges (and their endpoints) between this component and K . We say that a K -bridge B is attached to a vertex x of K if $x \in V(B \cap K)$. For $X \subseteq V(G)$, an X -bridge is a K -bridge where K is the edgeless graph with vertex set equal to X .

Suppose that C_1, C_2 are disjoint Π -homologic cycles that bound a subset D of the faces of Π . If D is a cylinder (i.e., its Euler characteristic is 0), we say that C_1 and C_2 are Π -homotopic. Similar definition applies if C_1 and C_2

touch. In that case they bound a *degenerate cylinder*. Instead of cycles we can also use simple closed curves γ_1, γ_2 on the surface of Π . Then we say that D is a (*degenerate*) *cylinder* between γ_1 and γ_2 and that γ_1, γ_2 are *homotopic*.

If $\gamma_1, \dots, \gamma_k$ are disjoint homotopic simple closed curves in S , then any two of them bound a cylinder, and they can be enumerated such that for $i = 1, \dots, k - 1$, the cylinder between γ_i and γ_{i+1} contains none of the other curves $\gamma_j, j \neq i, i + 1$. Such enumeration is *natural*. The same definition can be used when $\gamma_1, \dots, \gamma_k$ intersect but none of their intersection is a crossing.

Let D be a cylinder between two homotopic nonbounding curves in the surface of Π . Suppose that γ' and γ'' are noncrossing 1/2-curves in the interior of D . Then γ', γ'' are homotopic and they bound a cylinder $D' \subseteq D$ (possibly degenerate if they touch). Let $\gamma_1, \dots, \gamma_k$ be a maximal family of pairwise noncrossing 1/2-curves in D' such that $\gamma_1 = \gamma'$ and $\gamma_k = \gamma''$. Then we have:

THEOREM 6.2. *Suppose that the planar embedding Π' of the 2-connected graph G maximally coincides with Π . Let $\gamma_1, \dots, \gamma_k$ be as above and suppose that they are naturally enumerated. Let D_i be the (*degenerate*) cylinder between γ_i and $\gamma_{i+1}, i = 1, \dots, k - 1$. If every 1/2-curve that intersects $V(G) \cap D'$ is contained entirely in D , then there are patch faces F_1, \dots, F_{k-1} such that for $i = 1, \dots, k - 1, F_i$ contains a segment of γ_i and a segment of γ_{i+1} (where in case of 1-curves the segment could be just the vertex of G crossed by the curve) and such that $\partial F_i \cap D_i$ consists of two Π -facial segments joining vertices of $\gamma_i \cap G$ with vertices of $\gamma_{i+1} \cap G$.*

Proof. Let $\Gamma = \{\gamma_1, \dots, \gamma_k\}$. Consider an arbitrary consecutive pair of curves in Γ , say γ_1, γ_2 . We shall use the notation $f_i, g_i, x_i, y_i (i = 1, 2)$ introduced for Lemma 6.1. In case when γ_1 (or γ_2) is a 1-curve, we have $x_1 = y_1$ and in that case we can take as g_1 (say) any face containing x_1 .

First, we claim that the only patch vertices in D_1 are x_1, y_1, x_2, y_2 . If not, let x be another one. By Theorem 4.4, there is a 1/2-curve γ through x . By maximality of Γ, γ is not entirely contained in D_1 and hence it is a 2-curve. Since γ cannot escape out of D , it is homotopic to γ_1 and it must intersect γ_1 (or γ_2) twice. Let x' be the other patch vertex used by γ . If $x' \notin \{x_1, y_1\}$, then a segment of γ_1 between the points of $\gamma_1 \cap \gamma$ can be replaced by a segment of γ in D_1 , yielding a 2-curve in D_1 through x . A contradiction. If $x' = x_1$ (say), let f_1 and g be the faces used by γ . As above in the case when $x' \notin \{x_1, y_1\}$ we can try to change a segment of γ by a segment of γ_1 and get a contradiction. If this is not possible, then one of the possible replacements yields a 3-curve (and so $y_1 \neq x_1$), the other one a contractible curve γ'' through x and x_1 . Because of 2-reductions that we have performed prior to defining the patches, the interior of γ'' contains

just an edge $a = xx_1$ of G . By the above we may assume that no 1-curve intersects G in x . Therefore $g \neq f_1$. Let $\{a, b\}$ be the patch angle in g and $\{a, d''\}$ the patch angle in f_1 used by γ at x . Since these are patch angles, a and b are not Π' -consecutive and neither are a and d' . This implies that there is a patch angle $\{c, d\}$ such that $\{a, b\}$ and $\{c, d\}$ Π' -interlace and such that $\{a, b\} \cap \{c, d\} = \emptyset$. By Lemma 4.2, there is a 2-curve γ_0 through the patch angles $\{a, b\}$ and $\{c, d\}$. By using this 2-curve instead of γ , we obtain a contradiction since both patch faces of γ_0 are distinct from f_1, g_1, f_2, g_2 .

Suppose now that there is no face F_1 as claimed. Then $C = (\partial f_1 \cap D_1) \cup (\partial g_1 \cap D_1)$ (or $C = \partial f_1 \cap D_1$ if $x_1 = y_1$) is a Π -noncontractible cycle in D_1 that is composed of one or two patch edges between x_1 and y_1 . Similarly we have a Π -noncontractible cycle $C' = (\partial f_2 \cap D_1) \cup (\partial g_2 \cap D_1)$ (or $C' = \partial f_2 \cap D_1$) on the other side.

Let us first assume that $C \cap C' = \emptyset$. By Corollary 4.10, there exists a patch face Φ between C and C' . If $x_1 \neq y_1$, Φ cannot contain x_1 and y_1 since otherwise a segment of γ_1 could be replaced by an arc in Φ , contradicting maximality of Γ . Similarly for x_2, y_2 . Since G is 2-connected, Φ contains at least two patch angles. By maximality of Γ , these angles are at distinct vertices. Thus we may assume that $x_1, x_2 \in \Phi$, and these are the only patch vertices of Φ . Let α_i be the patch angle of Φ at x_i , $i = 1, 2$. If there is a 1/2-curve γ through α_1 , it must exit Φ through α_2 . Since γ cannot escape D , it is entirely in D_1 , a contradiction. Consequently, Theorem 4.4 implies that for $i = 1, 2$, there is a 1-curve γ'_i through x_i . The cylinder D'_1 between γ'_1, γ'_2 contains a 3-connected block of G . There is a planar embedding of G that coincides with Π in D'_1 , and coincides with Π' elsewhere. Since $D_1 \subseteq D'_1$, this contradicts the assumption that Π' maximally coincides with Π if $y_1 \neq x_1$ or $y_2 \neq x_2$. Otherwise, $F_1 = \Phi$ is the required face.

The other case is when C and C' intersect. By symmetry we may assume that $x_2 \in V(C \cap C')$. Then f_1 (say) contains x_2 . Suppose first that $\{x_1, y_1\} \cap \{x_2, y_2\} = \emptyset$. If $x_2 = y_2$, we can take f_1 for F_1 . Otherwise, denote by α the patch angle of f_1 at x_2 . A 1/2-curve γ through α crosses C at x_2 and thus it crosses C in another patch vertex. Let f_1 and f be the patch faces used by γ . If $f = f_2$ or $f = g_2$ or f is between C and C' , then γ can be taken to be entirely in D_1 , a contradiction with maximality of Γ . Otherwise, γ intersects C at y_2 and its re-routing through f_2 or g_2 gives rise to the previous case. The conclusion is that no 1/2-curve uses α . By Theorem 4.4, there is a 1-curve γ'_2 through x_2 . Note that γ'_2 is not in D_1 and is not equivalent to γ_2 . If $y_2 \in C \cap C'$, we repeat the same procedure at y_2 and obtain a contradiction. If $x_1 \in C \cap C'$, we see in the same way that there is a 1-curve γ'_1 through x_1 , and we conclude as in the previous case by reaching a contradiction with the assumption that Π' maximally coincides with Π . Hence $C \cap C' = \{x_2\}$. The cycle C' consists of two patch

edges that are Π' -facial segments. Since G is 2-connected, the only C' -bridge embedded in the Π' -interior of C' is the subgraph H of G bounded by γ and γ_2 . It is clear that Π' can be chosen so that the induced embedding of $H \cup C'$ coincides with its embedding induced by Π . This yields a contradiction with the fact that y_2 is a patch angle.

If $\{x_1, y_1\} \cap \{x_2, y_2\} \neq \emptyset$, we may assume that $x_1 = x_2$. If $x_1 = y_1$ (or $x_2 = y_2$), then $F_1 = f_2$ ($F_1 = f_1$, respectively) is the required face. Hence $x_1 \neq y_1$ and $x_2 \neq y_2$. If $y_1 \in V(C')$, then we get a new 2-curve in D_1 through the face f_2 (say) containing x_1 and y_1 and the corresponding face f_1 or g_1 on the other side. Thus we may assume that $y_1 \notin V(C')$ and $y_2 \notin V(C)$. We shall distinguish two subcases.

(a) D_1 contains a patch face Φ distinct from f_1, g_1, f_2, g_2 . By maximality of Γ , Φ contains just one vertex from $\{x_1, y_1\}$ and one vertex from $\{x_2, y_2\}$. In particular, Φ contains y_1 and y_2 and does not contain $x_1 = x_2$. We see as before that no 1/2-curve passes through the patch angles of Φ . Hence, there is a 1-curve δ_i through $y_i, i = 1, 2$. Let H be the subgraph of G between δ_1 and δ_2 . Then H is a $\{y_1, y_2\}$ -bridge in G . Since Π' maximally coincides with Π , it is easy to see that Π' coincides with Π everywhere on H except at y_1 and y_2 . This contradicts the fact that x_1 is a patch vertex.

(b) There is no such face Φ as in (a). Let H be the patch of G in D_1 . The restriction of Π' to H has two faces sharing x_1 but without any other vertices in common, and these two faces contain y_1 and y_2 , respectively. Therefore, every H -bridge in G is either attached to x_1 and y_1 or to x_1 and y_2 (and no other vertices). Since C' is a Π -nonbounding cycle, one of these components has an edge on the left side of C and an edge on the right side of C' (where the right side of C and the left side of C' point in D_1). This implies that there is a Π -facial walk Ψ that starts at x_1 , say, on the left of C , leaves D , and returns to C' at x_1 or y_2 on the right side. Then Ψ and f_2 contain a 1/2-curve that leaves D and crosses C , a contradiction.

This concludes our case analysis and establishes existence of faces F_1, \dots, F_{k-1} . ■

Suppose that G contains disjoint homotopic Π -nonbounding cycles C_1, C_2, C_3 such that the cylinder D bounded by C_1 and C_2 contains C_3 . The assumptions of Theorem 6.2 are clearly satisfied in D for any two non-crossing 1/2-curves γ' and γ'' in the interior of D . Theorem 6.2 shows that a general patch structure of G in D follows a variety of patterns, whose most general examples are shown in Fig. 6.

Mohar, Robertson, and Vitray [4] described the general patch structure of planar graphs in the projective plane. They either look as a “double wheel” or as the octahedron structure shown in Figure 3. As a corollary of

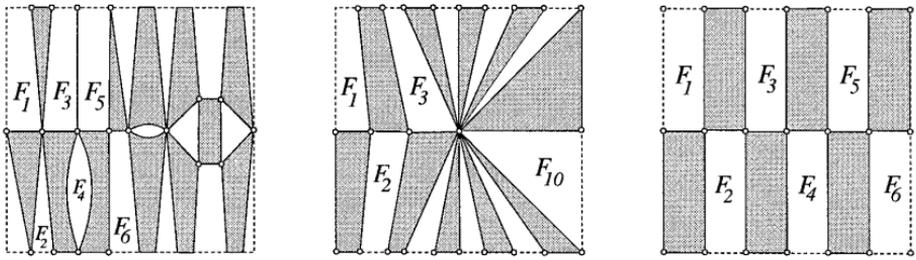


FIG. 6. A general patch structure.

Theorem 6.2 we will derive the corresponding result for the torus and the Klein bottle in case when there are three disjoint Π -nonbounding cycles and the face-width is two.

Suppose that G is Π -embedded in the torus with face-width 2 and that $\gamma_1, \dots, \gamma_k$ are noncrossing 2-curves. For $i=1, \dots, k$ denote by x_i, y_i and f_i, g_i the vertices and patch faces (respectively) used by γ_i . Curves γ_i and γ_{i+1} (index $i+1$ taken modulo k) are homotopic and hence they bound a (degenerate) cylinder D_i . We may assume that the curves are naturally enumerated such that the union of D_1, \dots, D_k and $\gamma_1, \dots, \gamma_k$ is the entire surface. If for each i , a segment of γ_i can be joined to a segment of γ_{i+1} by a curve in D_i that is disjoint from G , then we say that the embedding of G has the *chessboard structure*. This structure is *nondegenerate* if for each i , $\{x_i, y_i\} \cap \{x_{i+1}, y_{i+1}\} = \emptyset$. Such a structure with $k=6$ is represented in Fig. 7. Examples of degenerate chessboard structures are shown in Fig. 6. The same definition applies for the Klein bottle in which case we also require that the 2-curves γ_i ($1 \leq i \leq k$) are nonbounding.

The following is a corollary of Theorem 6.2.

COROLLARY 6.3. *Let G be a 2-connected planar graph embedded in the torus or the Klein bottle with face-width 2. If G contains three disjoint Π -nonbounding cycles, then the embedding has the chessboard structure.*

Proof. Disjoint Π -nonbounding cycles C_0, C_1, C_2 on the torus or the Klein bottle have the property that the removal of each of them leaves a

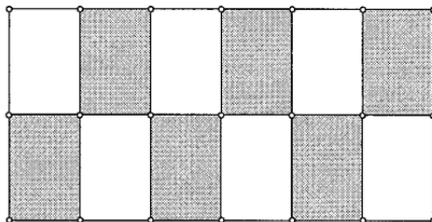


FIG. 7. A chessboard patch structure on the torus.

cylinder. By Theorem 4.9 there are 2-curves γ_i crossing C_i ($i = 0, 1, 2$). Now we apply Theorem 6.2 to get the chessboard structure between γ_i and γ_{i+1} by using the cylinder D obtained by cutting the surface along C_{i+2} (indices modulo 3). ■

Theorem 6.3 stimulated the following conjecture.

Conjecture 6.4. Suppose that G is a 2-connected planar graph that is Π -embedded with face-width 2, and that C_1, C_2, C_3 are disjoint homotopic Π -nonbounding cycles. Let k be the minimal number such that there exists a k -curve γ that intersects each of C_1, C_2, C_3 exactly once. Then k is equal to the maximal number t of pairwise disjoint cycles C'_1, \dots, C'_t homotopic to C_1 .

Clearly, $k \geq t$. By Corollary 6.3 we know that Conjecture 6.4 holds for the torus and the Klein bottle. Embeddings Π in the torus that have face-width 2 and no two disjoint Π -noncontractible cycles are classified in [3]. Examples show that the requirement of Conjecture 6.4 about existence of three disjoint cycles C_1, C_2, C_3 cannot be entirely omitted.

7. GENERALIZED WHITNEY'S THEOREM

The patch degree of a patch vertex can be arbitrarily large. Examples on the projective plane or the torus are easy to construct. On the other hand, we will show in this section that the patches or the patch faces cannot be too complicated if we restrict our attention to a fixed surface S .

Suppose that G is a 2-connected Π -embedded graph and C is a Π -contractible cycle of G such that only two vertices of C , say v and w , have incident edges that are embedded in the Π -exterior of C . Denote by D the Π -interior of C . Then we define a *Whitney 2-switching* of Π (with respect to C) as a reembedding of G such that the local rotation of each vertex in $D \setminus \{v, w\}$ is reversed, and local rotations at v and w are changed as follows. If $\pi_v = (e_1 e_2 \cdots e_d)$ where e_1, \dots, e_j are edges in D , then we change π_v to $(e_j e_{j-1} \cdots e_1 e_{j+1} \cdots e_d)$, and similarly at w . This operation that preserves the underlying surface of the embedding generates an equivalence relation among embeddings of G . It was proved by Whitney [7] that any two embeddings of G in the 2-sphere are equivalent.

Our main goal in this section is to prove a Whitney-type result for embeddings of planar graphs in an arbitrary fixed surface. For that purpose we also define a *k-switching operation* where $k > 1$ is an integer. Suppose that we have a patch P whose interior in S is homeomorphic to an open disk. Let v_1, \dots, v_k be the consecutive patch vertices that appear on the boundary of P , including possible multiple occurrences of the same vertex.

If there is a patch facial walk W containing v_1, \dots, v_k (in this or in the reverse order), we can reembed P in that face. Note that we may need to change the orientation in P and that sometimes there is more than one possibility how to do the reembedding. However, we require that the faces of the patch extension \tilde{G} in the patch P are unchanged. More generally, if there is a closed curve γ in S without self-crossings (but possibly touching itself) that bounds an open disk D in S and intersects G only in vertices v_1, \dots, v_k , the same reembedding of $G \cap D$ into the face W can be performed. Such a reembedding is called a k -switching. The Whitney 2-switching is a special case of this operation. If a k -switching does not change the underlying surface of the embedding, then it is invertible.

Two embeddings of G in the same surface S are *Whitney equivalent* if there is a sequence of 2-switchings, 3-switchings and 4-switchings transforming one embedding into the other. If a k -switching changes the underlying surface of the embedding, the Euler characteristic strictly decreases. Therefore all intermediate embeddings are also embeddings in S , and thus Whitney equivalence is an equivalence relation among embeddings of G in S .

Embeddings Π_1 and Π_2 of planar graphs G_1 and G_2 , respectively, in the same surface S are *patch equivalent* if there is a homeomorphism of S onto itself that induces a bijection on patch vertices, patch edges and patch faces of Π_1 and Π_2 , respectively. We also say that Π_1 and Π_2 have the same *patch structure*. Our next result shows that up to Whitney equivalence, there are not too many patch structures of embeddings of planar graphs.

THEOREM 7.1. *For each surface S there is a finite number of patch structures such that any embedding of a planar 2-connected graph G in S is Whitney equivalent to an embedding having one of these structures as its patch structure.*

In the proof of Theorem 7.1 we shall use the following lemma.

LEMMA 7.2 [2]. *Let S be a closed surface of genus $g \geq 1$, and $p \geq 1$ an integer. Suppose that Γ is a set of noncontractible simple closed curves in S that are either pairwise disjoint or they all pass through a point $x \in S$ and are disjoint elsewhere. If $|\Gamma| \geq 3pg$, then Γ contains a subset of $p + 1$ homotopic curves.*

Proof. By [2, Proposition 3.7] every set of $3g$ disjoint curves from Γ contains a pair of homotopic curves. By [2, Proposition 3.6], every set of $3g$ curves passing through x (and disjoint elsewhere) also contains a pair of homotopic curves. Now the lemma is immediate. ■

Proof of Theorem 7.1. It suffices to show that every embedding Π of a 2-connected planar graph G in S is Whitney equivalent to an embedding of G in S that has only a bounded number of patch edges.

Suppose that an embedding Π of G is Whitney equivalent to no embedding of G in S with fewer patch edges. We shall assume that the planar embedding Π' of G that determines the patches maximally coincides with Π . For every patch face Φ and a patch vertex x that occurs $t \geq 2$ times on Φ , we connect consecutive appearances of x on Φ by $t - 1$ 1-curves. Let Γ_1 be the set of the obtained 1-curves. For each simple and not bad patch angle take a 2-curve through it, and let Γ_2 be the set of the obtained 2-curves. We may assume that each 2-curve from Γ_2 passes through an angle that is not used by other curves from $\Gamma_1 \cup \Gamma_2$. By Corollary 4.6, it suffices to see that $|\Gamma_1| + |\Gamma_2|$ is bounded. We may assume that an arbitrary pair of curves from $\Gamma_1 \cup \Gamma_2$ intersects in the interior of a patch face Φ if and only if their angles in Φ interlace.

Suppose that $|\Gamma_1| \geq 216g^2$ where g is the genus of S . By our selection of the 1-curves, any two curves from Γ_1 intersect at most once. If $\gamma_1, \gamma_2 \in \Gamma_1$ are homotopic and they intersect in a patch vertex x , then they bound an open disc and they are either equivalent, or x is a cutvertex of G . None of these is possible. Hence, by Lemma 7.2, less than $3g$ curves from Γ_1 intersect in x . Similarly, if three homotopic curves from Γ_1 intersect in a point z inside a patch face, two of them bound an open disc D that contains the third one. It follows that the third 1-curve passes through a vertex of G that is not a patch vertex since the embedding Π' can be changed so that it matches the embedding of G in D . Now Lemma 7.2 implies that Γ_1 contains a set of more than $|\Gamma_1|/(9g) \geq 24g$ pairwise disjoint 1-curves. The same lemma implies that this subset contains nine disjoint homotopic 1-curves $\gamma_1, \dots, \gamma_9$. Assume that they are naturally enumerated. Let D_1, \dots, D_8 be the corresponding cylinders between the consecutive curves. For $i = 1, \dots, 7$, $D_i \cup D_{i+1}$ contains a Π -noncontractible cycle C_i . (Otherwise the embedding Π' can be changed so that it matches the embedding of G in $D_i \cup D_{i+1}$ which would contradict the fact that γ_{i+1} crosses G in a patch vertex.) Let D be the cylinder between C_1 and C_7 , and let $\gamma' = \gamma_4, \gamma'' = \gamma_6$. Theorem 6.2 implies that we can change the embedding of G between γ' and γ'' by using a sequence of 2/3/4-switchings so that the patch vertex of γ_5 disappears. This contradicts the minimality of the patch structure of our embedding. Hence, $|\Gamma_1| < 216g^2$.

Suppose now that there are patch vertices x, y such that p 2-curves from Γ_2 intersect x and y . These curves give rise to p simple arcs $\alpha_0, \dots, \alpha_{p-1}$ from x to y such that every arc uses a patch angle that is not used by other arcs and is not used by any of the 1-curves from Γ_1 . This implies that no two of these arcs are in the same patch face and hence they are internally disjoint. Let δ_i be the 2-curve composed of α_0 and $\alpha_i, 1 \leq i < p$. By contracting

α_0 to a point, we get $p-1$ simple closed curves intersecting in a single point. If $p > 6g$, then by Lemma 7.2, three of them, say $\delta_1, \delta_2, \delta_3$, are homotopic. This means that α_1, α_2 and α_2, α_3 , respectively, bound two discs containing distinct $\{x, y\}$ -bridges in G . Moreover, α_1 and α_3 (say) bound a disc that contains both of these components. Now, Π' can be changed so that these components merge into a single patch, a contradiction. Therefore $p \leq 6g$. Let Γ'_2 be a maximal subset of Γ_2 such that no two curves from Γ'_2 use the same pair of patch vertices. By the above, it suffices to see that $|\Gamma'_2|$ is bounded.

Suppose that there are patch faces Φ, Ψ such that p of the curves in Γ'_2 pass through Φ and Ψ . Let ϕ and ψ be points in the interior of Φ and Ψ , respectively. Since no two of the curves use the same pair of patch vertices, these curves determine at least $\lceil \sqrt{2p} \rceil$ internally disjoint simple arcs from ϕ to ψ consisting of segments of the curves. As above we see that no three of these arcs are homotopic, and hence $p \leq 18g^2$ by Lemma 7.2.

Suppose now that $\gamma_1, \dots, \gamma_p$ are 2-curves in Γ'_2 that all pass through a patch vertex x and a patch face Φ but any two of them use distinct second patch vertex and the patch face. As above we see that the vertex x appears on Φ at most $3g$ times. Hence $\gamma_1, \dots, \gamma_p$ contain a subset of at least $p/(3g)$ curves that intersect only in x . Let Γ''_2 be a maximal subset of Γ'_2 such that no two 2-curves from Γ''_2 intersect more than once. Let

$$r = \lfloor |\Gamma''_2|^{1/2} / (6g) \rfloor.$$

If a 2-curve γ from Γ''_2 intersects $12rg$ curves of Γ''_2 , then by Lemma 7.2 there is a subset $\Gamma = \{\gamma_1, \dots, \gamma_r\}$ of r pairwise homotopic nonequivalent 2-curves that all intersect γ in the same point. If there is no such curve γ , then Γ''_2 contains a subset of $|\Gamma''_2|^{1/2}/2$ disjoint curves. By Lemma 7.2, Γ''_2 has a subset $\Gamma = \{\gamma_1, \dots, \gamma_r\}$ of r pairwise homotopic nonequivalent 2-curves that are pairwise disjoint. By the above it suffices to see that r cannot be arbitrarily large.

Suppose that curves in Γ are pairwise disjoint. We claim that there is a sequence of switchings yielding an embedding of G in S with fewer patch angles if r is large enough. Suppose that $\gamma_1, \dots, \gamma_r$ are naturally enumerated. For $i = 1, \dots, r-1$, let x_i, y_i be the patch vertices used by γ_i and let D_i be the cylinder bounded by γ_i and γ_{i+1} . Denote by $D = D_1 \cup \dots \cup D_{r-1}$. If 18 of the cylinders D_i contain 1-curves, nine of these 1-curves are pairwise disjoint and homotopic. We get a contradiction as above. Otherwise, $q = \lfloor r/18 \rfloor$ consecutive cylinders, say D_1, \dots, D_q , contain no 1-curves. Menger's theorem implies that each D_i , $1 \leq i \leq q$, contains disjoint paths P_i, Q_i joining x_i with x_{i+1} (say) and y_i with y_{i+1} , respectively. At least every second cylinder D_i contains a path R_i joining P_i and Q_i (otherwise Π' could be changed and a patch vertex eliminated). Thus we can assume

that the paths R_i exist for all indices i and that they are disjoint (by taking a subset of our curves if necessary).

Suppose that the 2-curves γ_i are bounding. Then $\{x_i, y_i\}$ is a separating pair of G . Let G_1 be the subgraph of G on the left side of γ_2 (where we assume that D_2 is on the right side of γ_2), and let G_2 be the subgraph of G on the right side of γ_4 . Also, let $G_3 = G \cap (D_2 \cap D_3)$. Then G is edge-disjoint union of connected graphs G_1 , G_2 , and G_3 . The embedding Π restricted to G_3 is a planar embedding since $D_2 \cup D_3 \subset S$ is a cylinder. This embedding has a face F_1 containing x_2, y_2 and a face F_2 containing x_4, y_4 . If $F_1 = F_2$, then x_2, y_2 and x_4, y_4 do not interlace in this face. The embedding Π' restricted to G_1 contains x_2 and y_2 on the same face since there is a G_1 -bridge in G attached to x_2 and y_2 . The same conclusions apply for x_4, y_4 in Π' restricted to G_2 . Therefore we can change Π' into a planar embedding Π'' of G that coincides with Π' (or its inverse) on $G_1 \cup G_2$ and coincides with Π on G_3 . This is a contradiction with our choice of Π' .

We may now assume that γ_1 is nonbounding. We claim that $D'_i = D_i \cup \dots \cup D_{i+3}$ contains a Π -noncontractible cycle, $i = 2, 3, \dots, q - 4$. Let R be a path in the complement of D'_i from $\{x_i, y_i\}$ to $\{x_{i+4}, y_{i+4}\}$. Consider the planar embedding Π' restricted to the subgraph $H = \bigcup_{j=i+1}^{i+2} (P_j \cup Q_j \cup R_j)$. The H -bridge in G containing $R_{i-1} \cup R_{i+4} \cup R$ is attached to $x_{i+1}, y_{i+1}, x_{i+3}$, and y_{i+3} . This implies that H has all four vertices $y_{i+1}, x_{i+1}, x_{i+3}$ and y_{i+3} in the same face and in that order. The same holds for the embedding Π restricted to H if D'_i does not contain a Π -noncontractible cycle. Then Π' can be replaced by a planar embedding Π'' of G that coincides with Π in $G \cap (D_{i+1} \cup D_{i+2})$ and coincides with Π' elsewhere. This yields a contradiction with our choice of Π' .

Now, if q is sufficiently large, then D contains four disjoint Π -noncontractible cycles. We conclude as above in case of 1-curves by applying Theorem 6.2.

The above proof can also be used if $\gamma_1, \dots, \gamma_r$ intersect in the interior of the same patch face. We may now assume that there is a patch vertex x such that $\{\gamma_1, \dots, \gamma_r\}$ intersect in x . Suppose that the curves are naturally enumerated. (In case when γ_i are onesided, they cross at x , and then we extend the definition of naturally enumerated curves in the obvious way.) Let y_i be other patch vertex used by γ_i . If two of the disks D_i between γ_i and γ_{i+1} contain a 1-curve through x , then $G - x$ is disconnected. Therefore we may assume that for $i = 1, 2, \dots, q = \lfloor r/3 \rfloor$, $G \cap D_i \setminus \{x\}$ contains a path Q_i from y_i to y_{i+1} . If the curves in Γ are bounding, then we get a contradiction with the choice of Π' as above (if r is sufficiently large). The same arguments work if $\{x, y_1\}$ (and hence also $\{x, y_4\}$) is a separating set of the graph. Otherwise, there is a path R from y_1 to y_4 that is internally disjoint from x and from the discs D_1, D_2, D_3 . We argue as above to show existence of paths $R_i \subset D_i$ ($i = 1, 2, 3$) joining Q_i with x . After possible

reenumeration we may assume that $R_i \cap R_j = \{x\}$, $1 \leq i < j \leq 3$. The subgraph $H = \bigcup_{i=1}^3 (Q_i \cup R_i) \cup R$ is homeomorphic to K_4 and has unique embedding in the sphere. If $G \cap D_i$ ($1 \leq i \leq 3$) does not contain a Π -noncontractible cycle, then the embedding Π coincides with Π' on $G \cap D_i$ (by the choice of Π'). Then D_i contains a single patch attached to x , y_i and y_{i+1} . If D_i and D_{i+1} both have this property, then either y_{i+1} is not a patch vertex, or the patch in D_i can be reembedded by a 3-switching so that this patch and the patch in D_{i+1} merge into a single match. In the latter case, either the 3-switching in D_i or the 3-switching in D_{i+1} does not change the surface S . This cannot happen by the minimality of Π . On the other hand, if D_i contains a Π -noncontractible cycle of G , the path Q_i can be subdivided by taking the patch vertices in its interior. The refined partition into disks D_1, D_2, \dots is easily seen to contain patches without Π -noncontractible cycles, and then the above proof works.

Above results show that r is bounded by a constant depending only on g . This completes the proof. ■

Given a surface S , let Π_1, \dots, Π_N be the basic patch structures of Theorem 7.1. To perform the change of an embedding with patch structure Π_i into an embedding with patch structure Π_j ($1 \leq i, j \leq N$) we need to change only a bounded number of angles. By calling every such change a *generalized Whitney switching*, Theorem 7.1 can be formulated as follows.

COROLLARY 7.3. *Any two embeddings of a 2-connected planar graph in the same surface can be obtained from each other by performing a sequence of Whitney k -switchings for $k=2, 3, 4$ and applying (at most one) generalized Whitney switching operation.*

The following specific problem also occurred to us. Let G be a 5-connected planar triangulation. Change its planar rotation Π' at each vertex (to a local rotation Π) so that no two Π' -consecutive edges are Π -consecutive. (For example, if the local rotation in the plane at a vertex v of degree five is $\pi'_v = (e_1 e_2 \dots e_5)$, we can change it to $\pi_v = (e_1 e_3 e_5 e_2 e_4)$ and similarly at vertices of larger degrees.) Then every patch is just an edge. *Question:* Can we get an embedding Π of face-width two in this way? A negative answer to this question would support some further speculations that we have concerning Whitney equivalence of embeddings.

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