UNIVERSITY OF LJUBLJANA INSTITUTE OF MATHEMATICS, PHYSICS AND MECHANICS DEPARTMENT OF MATHEMATICS JADRANSKA 19, 1000 LJUBLJANA, SLOVENIA

Preprint series, Vol. 47 (2009), 1093

## CROSSING-CRITICAL GRAPHS WITH LARGE MAXIMUM DEGREE

Zdeněk Dvořák Bojan Mohar

ISSN 1318-4865

Ljubljana, July 9, 2009

## Crossing-critical graphs with large maximum degree

Zdeněk Dvořák<sup>\*†</sup> Department of Mathematics Simon Fraser University Burnaby, B.C. V5A 1S6 email: rakdver@kam.mff.cuni.cz Bojan Mohar<sup>‡§</sup> Department of Mathematics Simon Fraser University Burnaby, B.C. V5A 1S6 email: mohar@sfu.ca

March 14, 2009

## Abstract

A conjecture of Richter and Salazar about graphs that are critical for a fixed crossing number k is that they have bounded bandwidth. A weaker well-known conjecture is that their maximum degree is bounded in terms of k. In this note we disprove these conjectures for every  $k \ge 171$ , by providing examples of k-crossing-critical graphs with arbitrarily large maximum degree.

A graph is k-crossing-critical (or simply k-critical) if its crossing number is at least k, but every proper subgraph has crossing number smaller than k. Using the Excluded Grid Theorem of Robertson and Seymour [8], it is not hard to argue that k-crossing-critical graphs have bounded tree-width [2]. However, all known constructions of crossing-critical graphs suggested that their structure is "path-like". Salazar and Thomas conjectured (cf. [2]) that they have bounded path-width. This problem was solved by Hliněný [3], who proved that the path-width of k-critical graphs is bounded above by  $2^{f(k)}$ , where  $f(k) = (432 \log_2 k + 1488)k^3 + 1$ .

In the late 1990's, two other conjectures were proposed (see [7] or [6]).

<sup>\*</sup>Supported in part through a postdoctoral position at Simon Fraser University.

 $<sup>^{\</sup>dagger}\mathrm{On}$  leave from: Institute of Theoretical Informatics, Charles University, Prague, Czech Republic.

<sup>&</sup>lt;sup>‡</sup>Supported in part by the Research Grant P1–0297 of ARRS (Slovenia), by an NSERC Discovery Grant (Canada) and by the Canada Research Chair program.

 $<sup>^{\$}</sup>$  On leave from: IMFM & FMF, Department of Mathematics, University of Ljubljana, Ljubljana, Slovenia.

**Conjecture 1.** For every positive integer k, there exists an integer D(k) such that every k-crossing-critical graph has maximum degree less than D(k).

The second conjecture was proposed as an open problem in the 1990's by Carsten Thomassen and formulated as a conjecture by Richter and Salazar [7].

**Conjecture 2.** For every positive integer k, there exists an integer B(k) such that every k-crossing-critical graph has bandwidth at most B(k).

Conjecture 2 would be a strengthening of Hliněný's theorem about bounded path-width and would also imply Conjecture 1.

Hliněný and Salazar [5] recently made a step towards Conjecture 1 by proving that k-crossing-critical graphs cannot contain a subdivision of  $K_{2,N}$  with  $N = 30k^2 + 200k$ .

In this note we give examples of k-crossing-critical graphs of arbitrarily large maximum degree, thus disproving both Conjectures 1 and 2.

A special graph is a pair (G, T), where G is a graph and  $T \subseteq E(G)$ . The edges in the set T are called *thick edges* of the special graph. A *drawing* of a special graph (G, T) is a drawing of G such that the edges in T are not crossed. The crossing number  $\operatorname{cr}(G, T)$  of a special graph is the minimum number of edge crossings in a drawing of (G, T) in the plane. (We set  $\operatorname{cr}(G, T) = \infty$  if a thick edge is crossed in every drawing of G.) An edge  $e \in E(G) \setminus T$  is k-critical if  $\operatorname{cr}(G, T) \geq k$  and  $\operatorname{cr}(G-e, T) < k$ . Let  $\operatorname{crit}_k(G, T)$ be the set of k-critical edges of (G, T). If  $T = \emptyset$ , then we write just  $\operatorname{cr}(G)$ for the crossing number of G and  $\operatorname{crit}_k(G)$  for the set of k-critical edges of G. Note that the graph G is k-critical if  $\operatorname{crit}_k(G) = E(G)$ .

A standard result (see, e.g., [1]) is that we can eliminate the thick edges by replacing them with sufficiently dense subgraphs. (In fact, one can replace every edge xy by t = cr(G, T) + 1 parallel edges or by  $K_{2,t}$  if multiple edges are not desired.)

**Lemma 3.** For every special graph (G,T) with  $\operatorname{cr}(G,T) < \infty$  and for any k, there exists a graph  $\tilde{G} \supseteq G$  such that  $\operatorname{cr}(G,T) = \operatorname{cr}(\tilde{G})$  and  $\operatorname{crit}_k(G,T) \subseteq \operatorname{crit}_k(\tilde{G})$ .

Furthermore, note the following:

**Lemma 4.** Let k be an integer. Any graph G with  $cr(G) \ge k$  contains a k-crossing-critical subgraph H such that  $crit_k(G) \subseteq E(H)$ .

*Proof.* For a contradiction, suppose that G is a smallest counterexample. If G were k-critical, then we would set H = G, hence G contains a non-k-critical edge e. It follows that  $cr(G - e) \ge k$ . Let f be a k-critical edge in

G, i.e.,  $\operatorname{cr}(G - f) < k$ . As  $\operatorname{cr}((G - e) - f) \leq \operatorname{cr}(G - f) < k$ , f is a k-critical edge in G - e. Therefore,  $\operatorname{crit}_k(G) \subseteq \operatorname{crit}_k(G - e)$ . Since G is the smallest counterexample, G - e has a k-critical subgraph H with  $\operatorname{crit}_k(G - e) \subseteq E(H)$ . However,  $H \subseteq G$  and  $\operatorname{crit}_k(G) \subseteq E(H)$ , which is a contradiction.  $\Box$ 

Let us now proceed with the main result. Two paths  $P_1$  and  $P_2$  in a special graph are *almost edge-disjoint* if all the edges in  $E(P_1) \cap E(P_2)$  are thick.

**Lemma 5.** For any d, there exists a special graph (G, T) such that  $\operatorname{crit}_{171}(G, T)$  contains at least d edges incident with one of the vertices of G.

Proof. Let (G,T) be the special graph drawn as follows: we start with d+1 thick cycles  $C_0, C_1, \ldots, C_d$  intersecting in a vertex v, i.e.,  $C_i \cap C_j = \{v\}$  for  $0 \leq i < j \leq d$ . Their lengths are  $|C_0| = 28$ ,  $|C_d| = 24$  and  $|C_i| = 7$  for  $1 \leq i < d$ . They are drawn in the plane so that all their vertices are incident with the unbounded face and their clockwise order around v is  $C_0, C_1, \ldots, C_d$ . See Figure 1 illustrating the case d = 5. Let  $C_0 = va_1a_2 \ldots a_{19}b_1b_2b_3c_1^0c_2^0 \ldots c_5^0$ ,  $C_d = vt^db'_3b'_2b'_1a'_1a'_2 \ldots a'_{19}$  and  $C_i = vt^ic_1^ic_2^j \ldots c_5^i$  for  $1 \leq i < d$ . Furthermore, add d vertices  $s^1, \ldots, s^d$  adjacent to v. The clockwise cyclic order of the neighbors of v is  $a_1, c_5^0, s^1, t^1, c_5^1, s^2, t^2, c_5^2, \ldots, s^{d-1}, t^{d-1}, c_5^{d-1}, s^d, t^d, a'_{19}$ . For  $1 \leq i \leq d$ , add thick cycles  $K_i$  whose vertices in the clockwise order are  $t^i, s^i$ , and five new vertices  $\tilde{c}_5^{i-1}, \tilde{c}_4^{i-1}, \ldots, \tilde{c}_1^{i-1}$ . Finally, add the following edges:  $c_j^i \tilde{c}_j^i$  for  $0 \leq i < d$  and  $1 \leq j \leq 5$ ,  $a_i a'_i$  for  $1 \leq i \leq 19$  and  $b_i b'_i$  for  $1 \leq i \leq 3$ . As described,  $T = \bigcup_{i=0}^d E(C_i) \cup \bigcup_{i=1}^d E(K_i)$ . Let  $M = \{a_1a'_1, a_2a'_2, \ldots, a_{19}a'_{19}, b_1b'_1, b_2b'_2, b_3b'_3\}$ . This drawing  $\mathcal{G}$  of (G, T) has  $\binom{19}{2} = 171$  crossings, as the edges  $a_i a'_i$  and

This drawing  $\mathcal{G}$  of (G, T) has  $\binom{12}{2} = 171$  crossings, as the edges  $a_i a'_i$  and  $a_j a'_j$  intersect for each  $1 \leq i < j \leq 19$ , and there are no other crossings. Let us show that  $\operatorname{cr}(G, T) = 171$ . Let  $\mathcal{G}'$  be an arbitrary drawing of (G, T), and for a contradiction assume that it has less than 171 crossings. Let us first observe that every thick cycle  $C_i$  and  $K_j$  is an induced nonseparating cycle of G. Therefore it bounds a face of  $\mathcal{G}'$ . Consider the cyclic clockwise order of the neighbors of v according to the drawing  $\mathcal{G}'$ . For each cycle  $C_i$   $(0 \leq i \leq d)$ , the two edges of  $C_i$  incident with v are consecutive in this order, since  $C_i$  bounds a face. Without loss on generality, we assume that each cycle  $C_i$  is drawn clockwise, otherwise it is drawn anti-clockwise. We may assume that  $C_0$  is drawn clockwise. If  $C_d$  were drawn clockwise as well, then each pair of edges  $a_i a'_i$  and  $a_j a'_j$  with  $1 \leq i < j \leq 19$  would intersect, and the drawing  $\mathcal{G}'$  would have at least 171 crossings. Therefore,  $C_d$  is drawn anti-clockwise. It

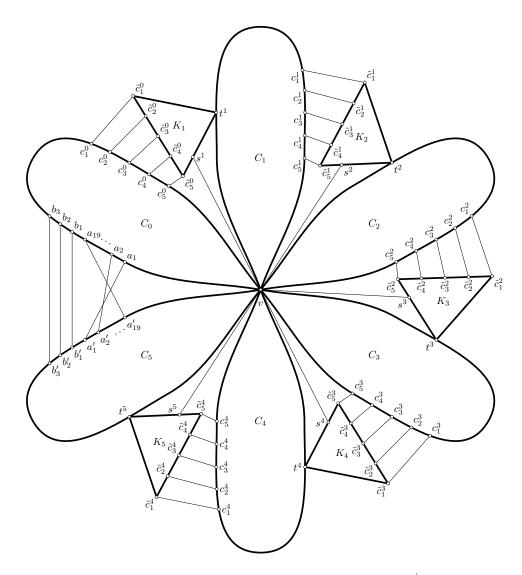


Figure 1: A special graph with critical edges  $vs^i$ 

follows that the edges  $a_i a'_i$  and  $b_j b'_j$  intersect for  $1 \le i \le 19$  and  $1 \le j \le 3$ , and the edges  $b_i b'_i$  and  $b_j b'_j$  intersect for  $1 \le i < j \le 3$ , giving 60 crossings. For  $1 \le i \le 5$ , let  $P_i$  be the path  $c_i^0 \tilde{c}_i^0 \tilde{c}_{i-1}^0 \dots \tilde{c}_1^0 t^1 c_1^1 c_2^1 \dots c_i^1 \tilde{c}_i^1 \dots \tilde{c}_1^1 t^2 \dots t^d$ . These paths are mutually almost edge-disjoint and each of them intersects all edges of M in the drawing  $\mathcal{G}'$ , thus contributing at least 110 crossings all together. Therefore, the drawing  $\mathcal{G}'$  has at least 170 crossings. Since we assume that this drawing has less than 171 crossings, we conclude that there are no other crossings.

The cycle  $va_1a'_1a'_2 \ldots a'_{19}$  splits the plane into two regions  $R_1$  and  $R_2$ , such that  $R_1$  contains the face bounded by  $C_0$  and  $R_2$  contains the face bounded by  $C_d$ . For j = 1, 2, let  $A_j$  be the set of cycles  $C_i$   $(0 \le i \le d)$ such that the face bounded by  $C_i$  lies in the region  $R_j$ . As  $P_1$  intersects the edge  $a_1a'_1$  only once,  $A_1 = \{C_0, C_1, \ldots, C_{k-1}\}$  and  $A_2 = \{C_k, C_{k+1}, \ldots, C_d\}$ for some k with  $1 \le k \le d$ . As the path  $P_1$  does not intersect itself, all cycles in  $A_1$  are drawn clockwise and their clockwise order around v is  $C_0$ ,  $C_1, \ldots, C_{k-1}$ . Similarly, all cycles in  $A_2$  are drawn anti-clockwise and their clockwise order around v is  $C_d$ ,  $C_{d-1}, \ldots, C_k$ .

Let us now consider the cycle  $K_k$ . Since the edges  $c_4^{k-1}\tilde{c}_4^{k-1}$  and  $c_5^{k-1}\tilde{c}_5^{k-1}$ do not intersect, the thick path  $c_5^{k-1}vt^ks^k\tilde{c}_5^{k-1}$  is not intersected, and  $C_{k-1}$  is drawn clockwise,  $K_k$  is drawn clockwise as well. Since  $C_k$  lies in the region  $R_2$ , the vertex  $t^k$  and thus the whole thick cycle  $K_k$  lie in  $R_2$ . However, that means that the edge  $s^k v$  intersects either the path  $P_1$  or the edge  $a_1a'_1$ , which is a contradiction. We conclude that cr(G,T) = 171.

On the other hand,  $\operatorname{cr}(G - vs^k, T) < 171$ , for  $1 \leq k \leq d$  (in fact,  $\operatorname{cr}(G - vs^k, T) = 170$ ). To see that, consider the drawing of  $(G - vs^k, T)$  in which the cycles  $C_0, C_1, \ldots, C_{k-1}$  are drawn clockwise, the cycles  $C_k, C_{k+1}, \ldots, C_d$  are drawn anti-clockwise, and the cyclic order of the neighbors of v is  $a_1c_5^0s^1t^1c_5^1\ldots s^{k-1}t^{k-1}c_5^{k-1}a'_{19}t^dc_5^{d-1}s^{d-1}t^{d-1}\ldots c_5^kt^k$ . The intersections of this drawing are of edges  $a_ia'_i$  with  $b_jb'_j$  for  $1 \leq i \leq 19$  and  $1 \leq j \leq 3$ , the edges  $b_ib'_i$  with  $b_jb'_j$  for  $1 \leq i < j \leq 3$ , and the edges  $c_i^{k-1}\tilde{c}_i^{k-1}$  with all edges of M for  $1 \leq i \leq 5$ . Therefore, the edge  $vs^k$  is 171-critical for each k, so v is incident with d critical edges.

We are ready for our main result.

**Theorem 6.** For every  $k \ge 171$  and every d, there exists a k-crossingcritical graph H containing a vertex of degree at least d.

*Proof.* Let (G, T) be the special graph constructed in Lemma 5. By Lemma 3, there exists a graph  $H' \supseteq G$  such that  $\operatorname{cr}(H') = \operatorname{cr}(G, T) \ge 171$  and

 $\operatorname{crit}_{171}(G,T) \subseteq \operatorname{crit}_{171}(H')$ . Let H be the 171-critical subgraph of H' obtained by Lemma 4. As  $\operatorname{crit}_{171}(G,T) \subseteq \operatorname{crit}_{171}(H') \subseteq E(H)$ , H contains at least d edges incident with one vertex, hence  $\Delta(H) \geq d$ . For k > 171 we add to H k - 171 copies of the graph  $K_5$  in order to get a k-crossing-critical graph.

Actually, in the proof of Theorem 6, we can take  $t = \lfloor \frac{k}{171} \rfloor$  copies of the graph H and k - 171t copies of  $K_5$ . This gives rise to a k-critical graph with  $t = \Omega(k)$  vertices of (arbitrarily) large degree. We conjecture that this is best possible in the following sense:

**Conjecture 7.** For every positive integer k there exists an integer D = D(k) such that every k-crossing-critical graph contains at most k vertices whose degree is larger than D.

It is not even obvious if there exist k-crossing-critical graphs with arbitrarily many vertices of degree more than 6. Surprisingly, such examples have been constructed recently by Hliněný [4]. His examples may contain arbitrarily many vertices of any even degree smaller than 2k - 1.

## References

- [1] M. DeVos, B. Mohar, R. Samal, Unexpected behaviour of crossing sequences, submitted.
- [2] J.F. Geelen, R.B. Richter, G. Salazar, *Embedding grids in surfaces*, European J. Combin. 25 (2004) 785–792.
- [3] P. Hliněný, Crossing-number critical graphs have bounded path-width, J. Combin. Theory Ser. B 88 (2003) 347–367.
- [4] P. Hliněný, New infinite families of almost-planar crossing-critical graphs, Electr. J. Combin. 15 (2008) #R102.
- [5] P. Hliněný, G. Salazar, Stars and bonds in crossing-critical graphs, preprint, 2008.
- [6] B. Mohar, J. Pach, B. Richter, R. Thomas, C. Thomassen, *Topological graph theory and crossing numbers*, Report on the BIRS 5-Day Workshop, 2007, 18 pages. http://www.birs.ca/workshops/2006/06w5067/report06w5067.pdf

- [7] R.B. Richter, G. Salazar, A survey of good crossing number theorems and questions, to appear.
- [8] N. Robertson and P. D. Seymour, Graph minors. V. Excluding a planar graph, J. Combin. Theory, Ser. B 41 (1986) 92–114.