# CROSSING-CRITICAL GRAPHS <br> WITH LARGE MAXIMUM <br> DEGREE 

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# Crossing-critical graphs with large maximum degree 

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#### Abstract

A conjecture of Richter and Salazar about graphs that are critical for a fixed crossing number $k$ is that they have bounded bandwidth. A weaker well-known conjecture is that their maximum degree is bounded in terms of $k$. In this note we disprove these conjectures for every $k \geq 171$, by providing examples of $k$-crossing-critical graphs with arbitrarily large maximum degree.


A graph is $k$-crossing-critical (or simply $k$-critical) if its crossing number is at least $k$, but every proper subgraph has crossing number smaller than $k$. Using the Excluded Grid Theorem of Robertson and Seymour [8], it is not hard to argue that $k$-crossing-critical graphs have bounded tree-width [2]. However, all known constructions of crossing-critical graphs suggested that their structure is "path-like". Salazar and Thomas conjectured (cf. [2]) that they have bounded path-width. This problem was solved by Hliněný [3], who proved that the path-width of $k$-critical graphs is bounded above by $2^{f(k)}$, where $f(k)=\left(432 \log _{2} k+1488\right) k^{3}+1$.

In the late 1990's, two other conjectures were proposed (see [7] or [6]).

[^0]Conjecture 1. For every positive integer $k$, there exists an integer $D(k)$ such that every $k$-crossing-critical graph has maximum degree less than $D(k)$.

The second conjecture was proposed as an open problem in the 1990's by Carsten Thomassen and formulated as a conjecture by Richter and Salazar [7].

Conjecture 2. For every positive integer $k$, there exists an integer $B(k)$ such that every $k$-crossing-critical graph has bandwidth at most $B(k)$.

Conjecture 2 would be a strengthening of Hliněný's theorem about bounded path-width and would also imply Conjecture 1.

Hliněný and Salazar [5] recently made a step towards Conjecture 1 by proving that $k$-crossing-critical graphs cannot contain a subdivision of $K_{2, N}$ with $N=30 k^{2}+200 k$.

In this note we give examples of $k$-crossing-critical graphs of arbitrarily large maximum degree, thus disproving both Conjectures 1 and 2.

A special graph is a pair $(G, T)$, where $G$ is a graph and $T \subseteq E(G)$. The edges in the set $T$ are called thick edges of the special graph. A drawing of a special graph $(G, T)$ is a drawing of $G$ such that the edges in $T$ are not crossed. The crossing number $\operatorname{cr}(G, T)$ of a special graph is the minimum number of edge crossings in a drawing of $(G, T)$ in the plane. (We set $\operatorname{cr}(G, T)=\infty$ if a thick edge is crossed in every drawing of $G$.) An edge $e \in E(G) \backslash T$ is $k$-critical if $\operatorname{cr}(G, T) \geq k$ and $\operatorname{cr}(G-e, T)<k$. Let $\operatorname{crit}_{k}(G, T)$ be the set of $k$-critical edges of $(G, T)$. If $T=\emptyset$, then we write just $\operatorname{cr}(G)$ for the crossing number of $G$ and $\operatorname{crit}_{k}(G)$ for the set of $k$-critical edges of $G$. Note that the graph $G$ is $k$-critical if $\operatorname{crit}_{k}(G)=E(G)$.

A standard result (see, e.g., [1]) is that we can eliminate the thick edges by replacing them with sufficiently dense subgraphs. (In fact, one can replace every edge $x y$ by $t=\operatorname{cr}(G, T)+1$ parallel edges or by $K_{2, t}$ if multiple edges are not desired.)

Lemma 3. For every special graph $(G, T)$ with $\operatorname{cr}(G, T)<\infty$ and for any $k$, there exists a graph $\tilde{G} \supseteq G$ such that $\operatorname{cr}(G, T)=\operatorname{cr}(\tilde{G})$ and $\operatorname{crit}_{k}(G, T) \subseteq$ $\operatorname{crit}_{k}(\tilde{G})$.

Furthermore, note the following:
Lemma 4. Let $k$ be an integer. Any graph $G$ with $\operatorname{cr}(G) \geq k$ contains a $k$-crossing-critical subgraph $H$ such that $\operatorname{crit}_{k}(G) \subseteq E(H)$.

Proof. For a contradiction, suppose that $G$ is a smallest counterexample. If $G$ were $k$-critical, then we would set $H=G$, hence $G$ contains a non- $k$ critical edge $e$. It follows that $\operatorname{cr}(G-e) \geq k$. Let $f$ be a $k$-critical edge in
$G$, i.e., $\operatorname{cr}(G-f)<k$. As $\operatorname{cr}((G-e)-f) \leq \operatorname{cr}(G-f)<k, f$ is a $k$-critical edge in $G-e$. Therefore, $\operatorname{crit}_{k}(G) \subseteq \operatorname{crit}_{k}(G-e)$. Since $G$ is the smallest counterexample, $G-e$ has a $k$-critical subgraph $H$ with $\operatorname{crit}_{k}(G-e) \subseteq E(H)$. However, $H \subseteq G$ and $\operatorname{crit}_{k}(G) \subseteq E(H)$, which is a contradiction.

Let us now proceed with the main result. Two paths $P_{1}$ and $P_{2}$ in a special graph are almost edge-disjoint if all the edges in $E\left(P_{1}\right) \cap E\left(P_{2}\right)$ are thick.

Lemma 5. For any d, there exists a special graph $(G, T)$ such that $\operatorname{crit}_{171}(G, T)$ contains at least $d$ edges incident with one of the vertices of $G$.

Proof. Let $(G, T)$ be the special graph drawn as follows: we start with $d+1$ thick cycles $C_{0}, C_{1}, \ldots, C_{d}$ intersecting in a vertex $v$, i.e., $C_{i} \cap C_{j}=\{v\}$ for $0 \leq i<j \leq d$. Their lengths are $\left|C_{0}\right|=28,\left|C_{d}\right|=24$ and $\left|C_{i}\right|=7$ for $1 \leq$ $i<d$. They are drawn in the plane so that all their vertices are incident with the unbounded face and their clockwise order around $v$ is $C_{0}, C_{1}, \ldots, C_{d}$. See Figure 1 illustrating the case $d=5$. Let $C_{0}=v a_{1} a_{2} \ldots a_{19} b_{1} b_{2} b_{3} c_{1}^{0} c_{2}^{0} \ldots c_{5}^{0}$, $C_{d}=v t^{d} b_{3}^{\prime} b_{2}^{\prime} b_{1}^{\prime} a_{1}^{\prime} a_{2}^{\prime} \ldots a_{19}^{\prime}$ and $C_{i}=v t^{i} c_{1}^{i} c_{2}^{i} \ldots c_{5}^{i}$ for $1 \leq i<d$. Furthermore, add $d$ vertices $s^{1}, \ldots, s^{d}$ adjacent to $v$. The clockwise cyclic order of the neighbors of $v$ is $a_{1}, c_{5}^{0}, s^{1}, t^{1}, c_{5}^{1}, s^{2}, t^{2}, c_{5}^{2}, \ldots, s^{d-1}, t^{d-1}, c_{5}^{d-1}, s^{d}, t^{d}, a_{19}^{\prime}$. For $1 \leq i \leq d$, add thick cycles $K_{i}$ whose vertices in the clockwise order are $t^{i}, s^{i}$, and five new vertices $\tilde{c}_{5}^{i-1}, \tilde{c}_{4}^{i-1}, \ldots, \tilde{c}_{1}^{i-1}$. Finally, add the following edges: $c_{j}^{i} \tilde{c}_{j}^{i}$ for $0 \leq i<d$ and $1 \leq j \leq 5, a_{i} a_{i}^{\prime}$ for $1 \leq i \leq 19$ and $b_{i} b_{i}^{\prime}$ for $1 \leq i \leq 3$. As described, $T=\bigcup_{i=0}^{d} E\left(C_{i}\right) \cup \bigcup_{i=1}^{d} E\left(K_{i}\right)$. Let $M=\left\{a_{1} a_{1}^{\prime}, a_{2} a_{2}^{\prime}, \ldots, a_{19} a_{19}^{\prime}, b_{1} b_{1}^{\prime}, b_{2} b_{2}^{\prime}, b_{3} b_{3}^{\prime}\right\}$.

This drawing $\mathcal{G}$ of $(G, T)$ has $\binom{19}{2}=171$ crossings, as the edges $a_{i} a_{i}^{\prime}$ and $a_{j} a_{j}^{\prime}$ intersect for each $1 \leq i<j \leq 19$, and there are no other crossings. Let us show that $\operatorname{cr}(G, T)=171$. Let $\mathcal{G}^{\prime}$ be an arbitrary drawing of $(G, T)$, and for a contradiction assume that it has less than 171 crossings. Let us first observe that every thick cycle $C_{i}$ and $K_{j}$ is an induced nonseparating cycle of $G$. Therefore it bounds a face of $\mathcal{G}^{\prime}$. Consider the cyclic clockwise order of the neighbors of $v$ according to the drawing $\mathcal{G}^{\prime}$. For each cycle $C_{i}(0 \leq i \leq d)$, the two edges of $C_{i}$ incident with $v$ are consecutive in this order, since $C_{i}$ bounds a face. Without loss on generality, we assume that each cycle $C_{i}$ bounds a face distinct from the unbounded one. If the cyclic order of the vertices around the face $C_{i}$ is the same as in the drawing $\mathcal{G}$, we say that $C_{i}$ is drawn clockwise, otherwise it is drawn anti-clockwise. We may assume that $C_{0}$ is drawn clockwise. If $C_{d}$ were drawn clockwise as well, then each pair of edges $a_{i} a_{i}^{\prime}$ and $a_{j} a_{j}^{\prime}$ with $1 \leq i<j \leq 19$ would intersect, and the drawing $\mathcal{G}^{\prime}$ would have at least 171 crossings. Therefore, $C_{d}$ is drawn anti-clockwise. It


Figure 1: A special graph with critical edges $v s^{i}$
follows that the edges $a_{i} a_{i}^{\prime}$ and $b_{j} b_{j}^{\prime}$ intersect for $1 \leq i \leq 19$ and $1 \leq j \leq 3$, and the edges $b_{i} b_{i}^{\prime}$ and $b_{j} b_{j}^{\prime}$ intersect for $1 \leq i<j \leq 3$, giving 60 crossings. For $1 \leq i \leq 5$, let $P_{i}$ be the path $c_{i}^{0} \tilde{c}_{i}^{0} \tilde{c}_{i-1}^{0} \ldots \tilde{c}_{1}^{0} t^{1} c_{1}^{1} c_{2}^{1} \ldots c_{i}^{1} \tilde{c}_{i}^{1} \ldots \tilde{c}_{1}^{1} t^{2} \ldots t^{d}$. These paths are mutually almost edge-disjoint and each of them intersects all edges of $M$ in the drawing $\mathcal{G}^{\prime}$, thus contributing at least 110 crossings all together. Therefore, the drawing $\mathcal{G}^{\prime}$ has at least 170 crossings. Since we assume that this drawing has less than 171 crossings, we conclude that there are no other crossings.

The cycle $v a_{1} a_{1}^{\prime} a_{2}^{\prime} \ldots a_{19}^{\prime}$ splits the plane into two regions $R_{1}$ and $R_{2}$, such that $R_{1}$ contains the face bounded by $C_{0}$ and $R_{2}$ contains the face bounded by $C_{d}$. For $j=1,2$, let $A_{j}$ be the set of cycles $C_{i}(0 \leq i \leq d)$ such that the face bounded by $C_{i}$ lies in the region $R_{j}$. As $P_{1}$ intersects the edge $a_{1} a_{1}^{\prime}$ only once, $A_{1}=\left\{C_{0}, C_{1}, \ldots, C_{k-1}\right\}$ and $A_{2}=\left\{C_{k}, C_{k+1}, \ldots, C_{d}\right\}$ for some $k$ with $1 \leq k \leq d$. As the path $P_{1}$ does not intersect itself, all cycles in $A_{1}$ are drawn clockwise and their clockwise order around $v$ is $C_{0}$, $C_{1}, \ldots, C_{k-1}$. Similarly, all cycles in $A_{2}$ are drawn anti-clockwise and their clockwise order around $v$ is $C_{d}, C_{d-1}, \ldots, C_{k}$.

Let us now consider the cycle $K_{k}$. Since the edges $c_{4}^{k-1} \tilde{c}_{4}^{k-1}$ and $c_{5}^{k-1} \tilde{c}_{5}^{k-1}$ do not intersect, the thick path $c_{5}^{k-1} v t^{k} s^{k} \tilde{c}_{5}^{k-1}$ is not intersected, and $C_{k-1}$ is drawn clockwise, $K_{k}$ is drawn clockwise as well. Since $C_{k}$ lies in the region $R_{2}$, the vertex $t^{k}$ and thus the whole thick cycle $K_{k}$ lie in $R_{2}$. However, that means that the edge $s^{k} v$ intersects either the path $P_{1}$ or the edge $a_{1} a_{1}^{\prime}$, which is a contradiction. We conclude that $\operatorname{cr}(G, T)=171$.

On the other hand, $\operatorname{cr}\left(G-v s^{k}, T\right)<171$, for $1 \leq k \leq d$ (in fact, $\left.\operatorname{cr}\left(G-v s^{k}, T\right)=170\right)$. To see that, consider the drawing of $\left(G-v s^{k}, T\right)$ in which the cycles $C_{0}, C_{1}, \ldots, C_{k-1}$ are drawn clockwise, the cycles $C_{k}, C_{k+1}$, $\ldots, C_{d}$ are drawn anti-clockwise, and the cyclic order of the neighbors of $v$ is $a_{1} c_{5}^{0} s^{1} t^{1} c_{5}^{1} \ldots s^{k-1} t^{k-1} c_{5}^{k-1} a_{19}^{\prime} t^{d} c_{5}^{d-1} s^{d-1} t^{d-1} \ldots c_{5}^{k} t^{k}$. The intersections of this drawing are of edges $a_{i} a_{i}^{\prime}$ with $b_{j} b_{j}^{\prime}$ for $1 \leq i \leq 19$ and $1 \leq j \leq 3$, the edges $b_{i} b_{i}^{\prime}$ with $b_{j} b_{j}^{\prime}$ for $1 \leq i<j \leq 3$, and the edges $c_{i}^{k-1} \tilde{c}_{i}^{k-1}$ with all edges of $M$ for $1 \leq i \leq 5$. Therefore, the edge $v s^{k}$ is 171-critical for each $k$, so $v$ is incident with $d$ critical edges.

We are ready for our main result.
Theorem 6. For every $k \geq 171$ and every $d$, there exists a $k$-crossingcritical graph $H$ containing a vertex of degree at least $d$.

Proof. Let $(G, T)$ be the special graph constructed in Lemma 5. By Lemma 3, there exists a graph $H^{\prime} \supseteq G$ such that $\operatorname{cr}\left(H^{\prime}\right)=\operatorname{cr}(G, T) \geq 171$ and
$\operatorname{crit}_{171}(G, T) \subseteq \operatorname{crit}_{171}\left(H^{\prime}\right)$. Let $H$ be the 171-critical subgraph of $H^{\prime}$ obtained by Lemma 4. As $\operatorname{crit}_{171}(G, T) \subseteq \operatorname{crit}_{171}\left(H^{\prime}\right) \subseteq E(H), H$ contains at least $d$ edges incident with one vertex, hence $\Delta(H) \geq d$. For $k>171$ we add to $H k-171$ copies of the graph $K_{5}$ in order to get a $k$-crossing-critical graph.

Actually, in the proof of Theorem 6, we can take $t=\left\lfloor\frac{k}{171}\right\rfloor$ copies of the graph $H$ and $k-171 t$ copies of $K_{5}$. This gives rise to a $k$-critical graph with $t=\Omega(k)$ vertices of (arbitrarily) large degree. We conjecture that this is best possible in the following sense:

Conjecture 7. For every positive integer $k$ there exists an integer $D=D(k)$ such that every $k$-crossing-critical graph contains at most $k$ vertices whose degree is larger than $D$.

It is not even obvious if there exist $k$-crossing-critical graphs with arbitrarily many vertices of degree more than 6 . Surprisingly, such examples have been constructed recently by Hliněný [4]. His examples may contain arbitrarily many vertices of any even degree smaller than $2 k-1$.

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