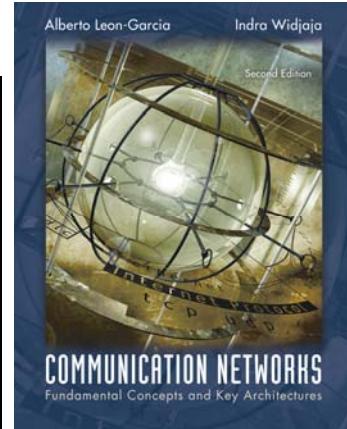


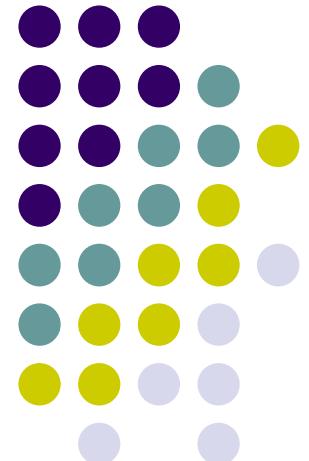
# Chapter 4

# Circuit-Switching

# Networks



Multiplexing  
SONET  
Transport Networks  
Circuit Switches  
The Telephone Network  
Signaling  
Traffic and Overload Control in Telephone Networks  
Cellular Telephone Networks

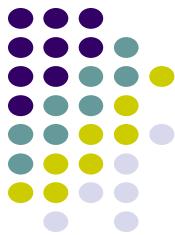


# Circuit Switching Networks

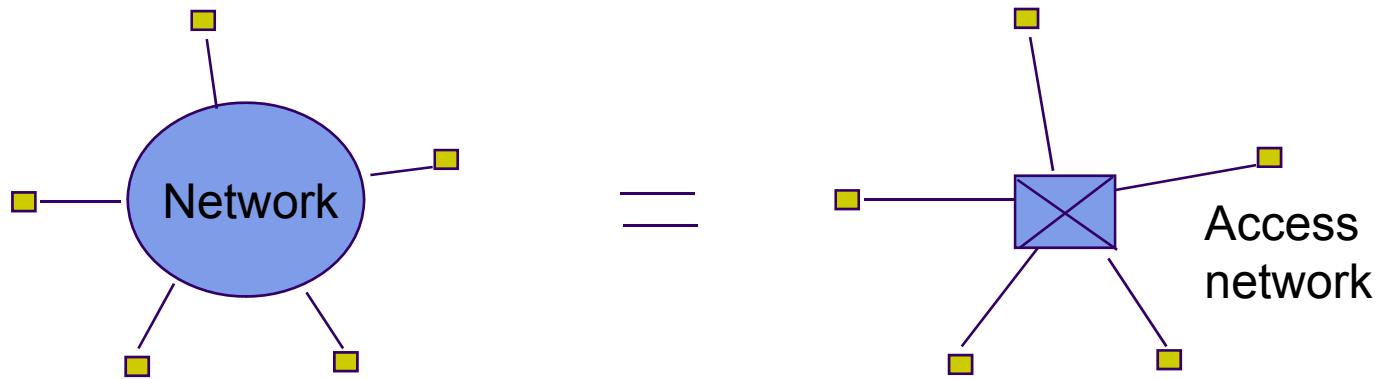


- End-to-end dedicated circuits between clients
  - Client can be a person or equipment (router or switch)
- Circuit can take different forms
  - Dedicated path for the transfer of electrical current
  - Dedicated time slots for transfer of voice samples
  - Dedicated frames for transfer of Nx51.84 Mbps signals
  - Dedicated wavelengths for transfer of optical signals
- Circuit switching networks require:
  - Multiplexing & switching of circuits
  - Signaling & control for establishing circuits
- These are the subjects covered in this chapter

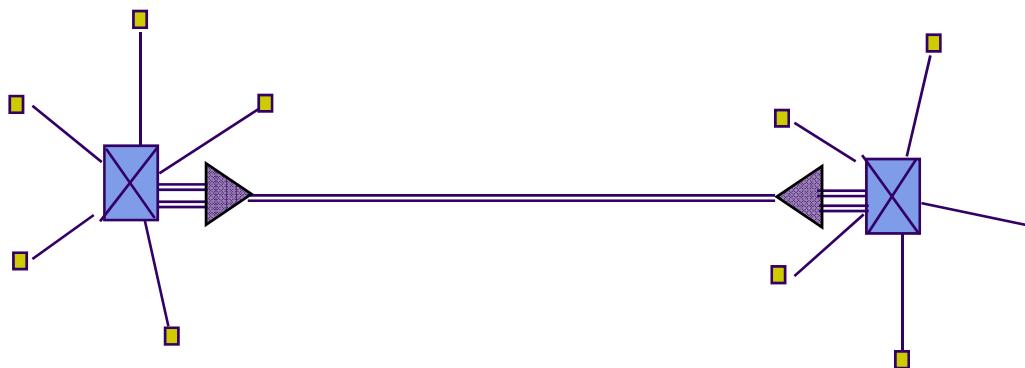
# How a network grows



(a) A switch provides the network to a cluster of users, e.g. a telephone switch connects a local community



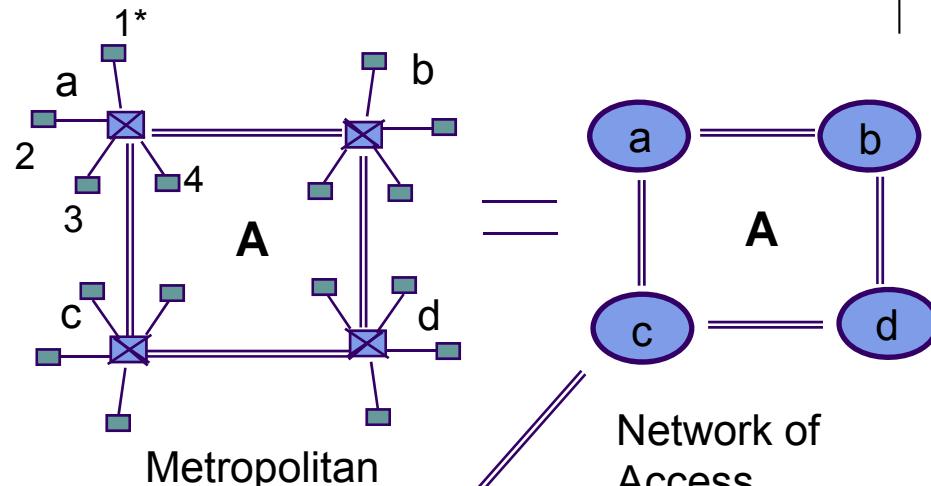
(b) A multiplexer connects two access networks, e.g. a high speed line connects two switches



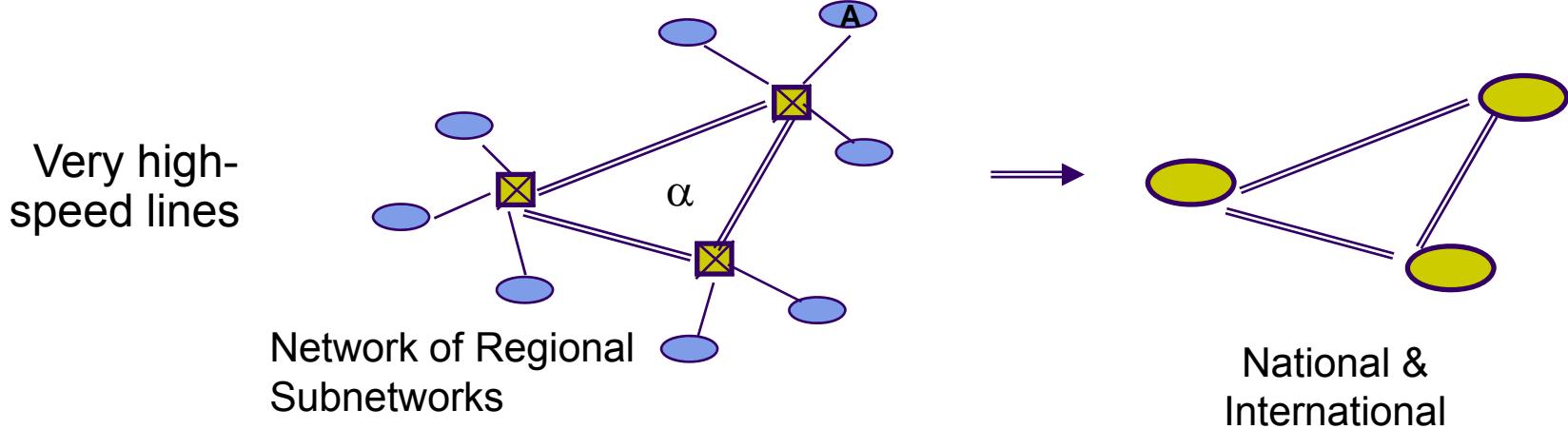
# A Network Keeps Growing



(a) Metropolitan network A viewed as Network A of Access Subnetworks

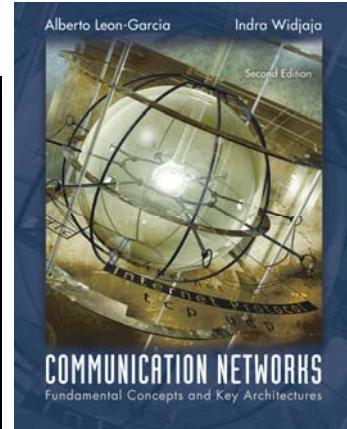


(b) National network viewed as Network of Regional Subnetworks (including A)

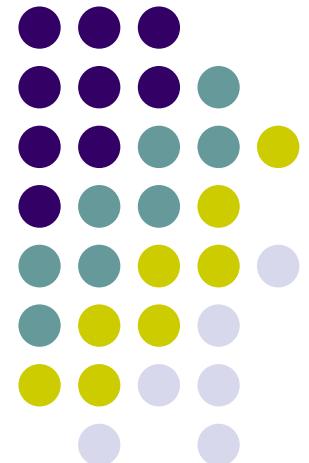


# Chapter 4

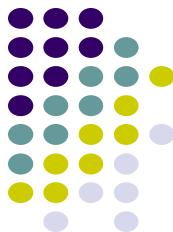
# Circuit-Switching Networks



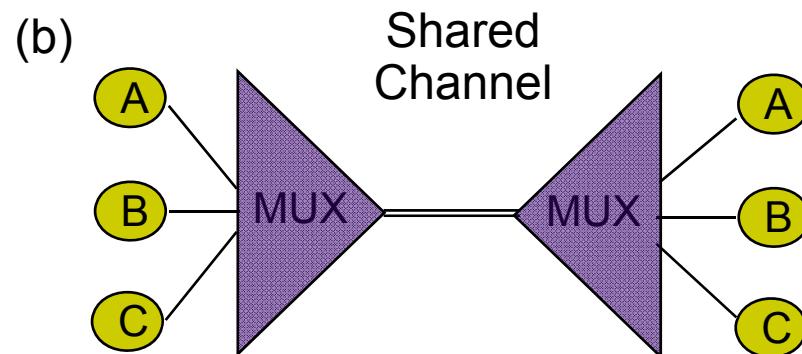
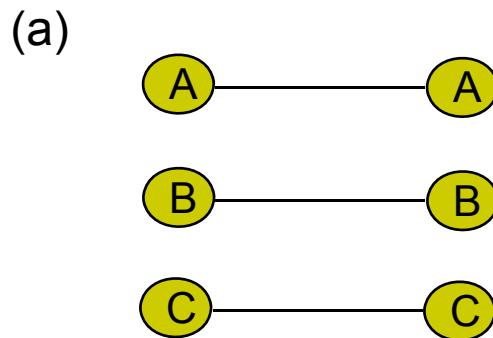
## *Multiplexing*

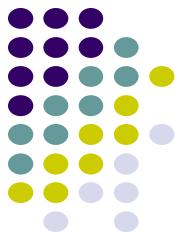


# Multiplexing



- Multiplexing involves the sharing of a transmission channel (resource) by several connections or information flows
  - Channel = 1 wire, 1 optical fiber, or 1 frequency band
- Significant economies of scale can be achieved by combining many signals into one
  - Fewer wires/pole; fiber replaces thousands of cables
- Implicit or explicit information is required to demultiplex the information flows.

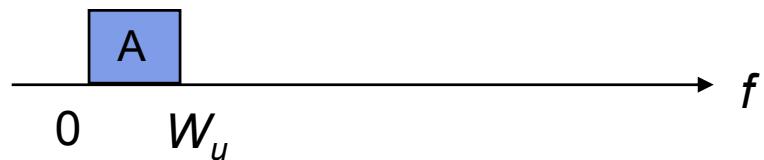




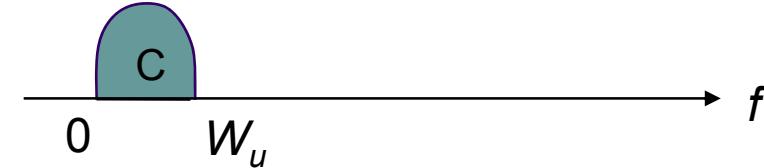
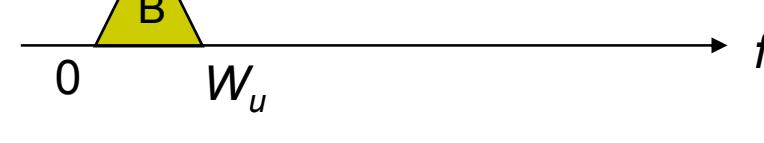
# Frequency-Division Multiplexing

- Channel divided into frequency slots

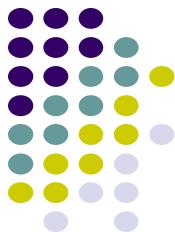
(a) Individual signals occupy  $W_u$  Hz



(b) Combined signal fits into channel bandwidth

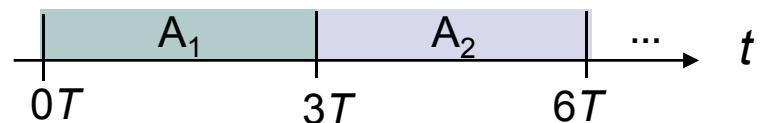


- Guard bands required
- AM or FM radio stations
- TV stations in air or cable
- Analog telephone systems

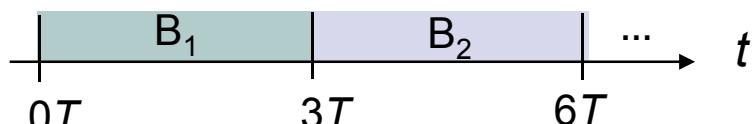


# Time-Division Multiplexing

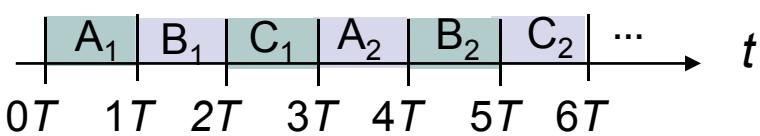
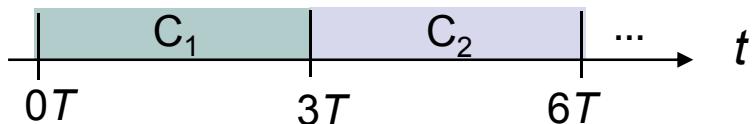
- High-speed digital channel divided into time slots



(a) Each signal transmits 1 unit every  $3T$  seconds



(b) Combined signal transmits 1 unit every  $T$  seconds

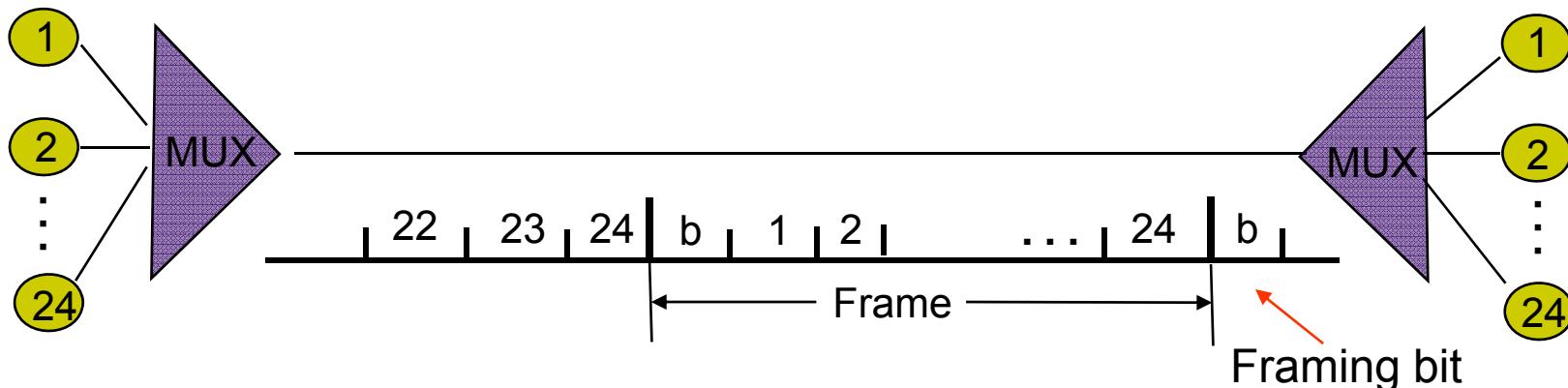


- Framing required
- Telephone digital transmission
- Digital transmission in backbone network



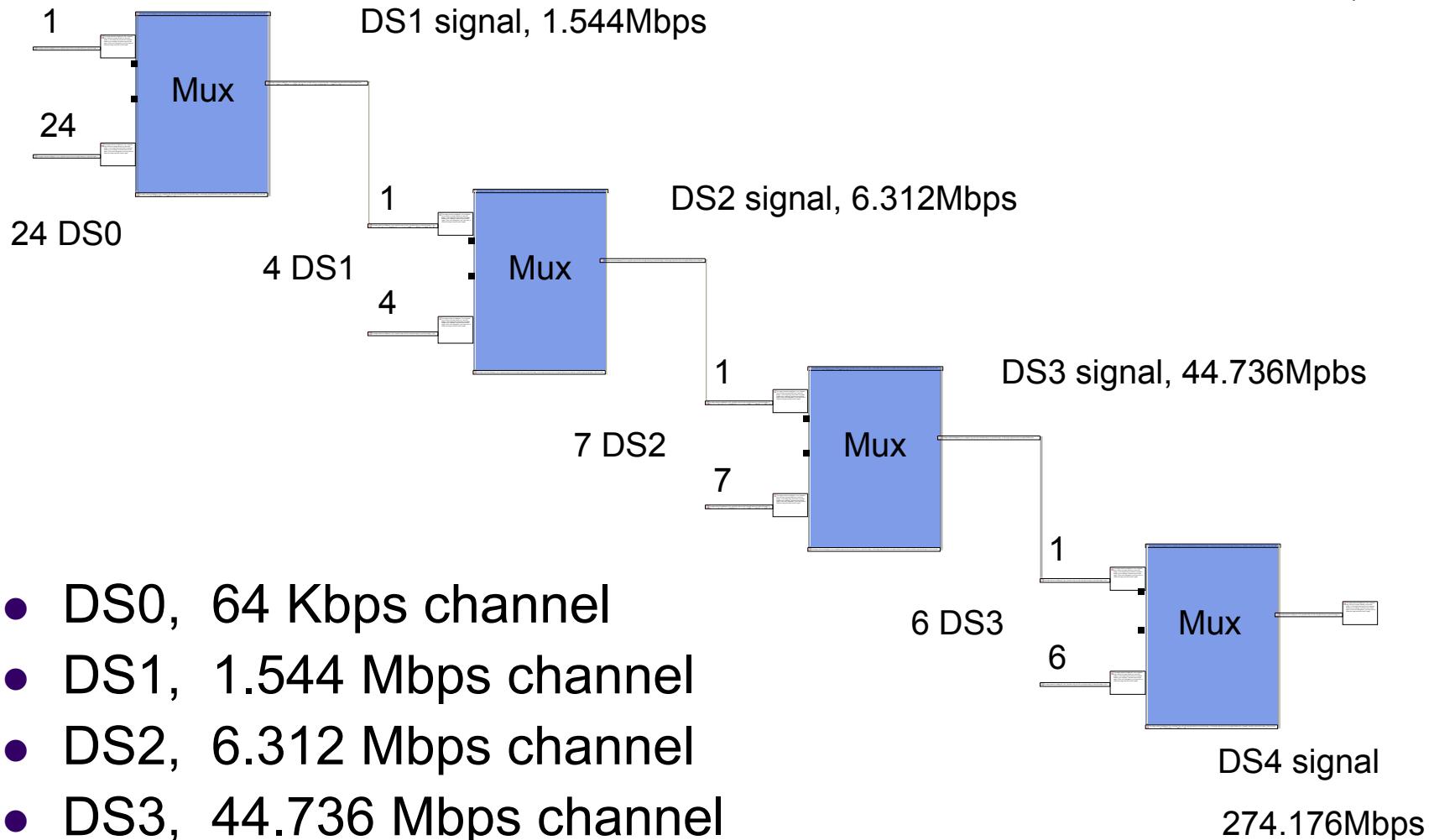
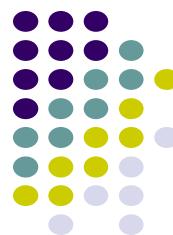
# T-Carrier System

- Digital telephone system uses TDM.
- PCM voice channel is basic unit for TDM
  - 1 channel = 8 bits/sample x 8000 samples/sec. = 64 kbps
- T-1 carrier carries Digital Signal 1 (DS-1) that combines 24 voice channels into a digital stream:



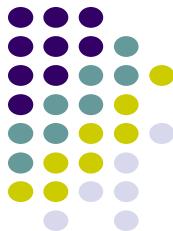
$$\begin{aligned}\text{Bit Rate} &= 8000 \text{ frames/sec.} \times (1 + 8 \times 24) \text{ bits/frame} \\ &= 1.544 \text{ Mbps}\end{aligned}$$

# North American Digital Multiplexing Hierarchy

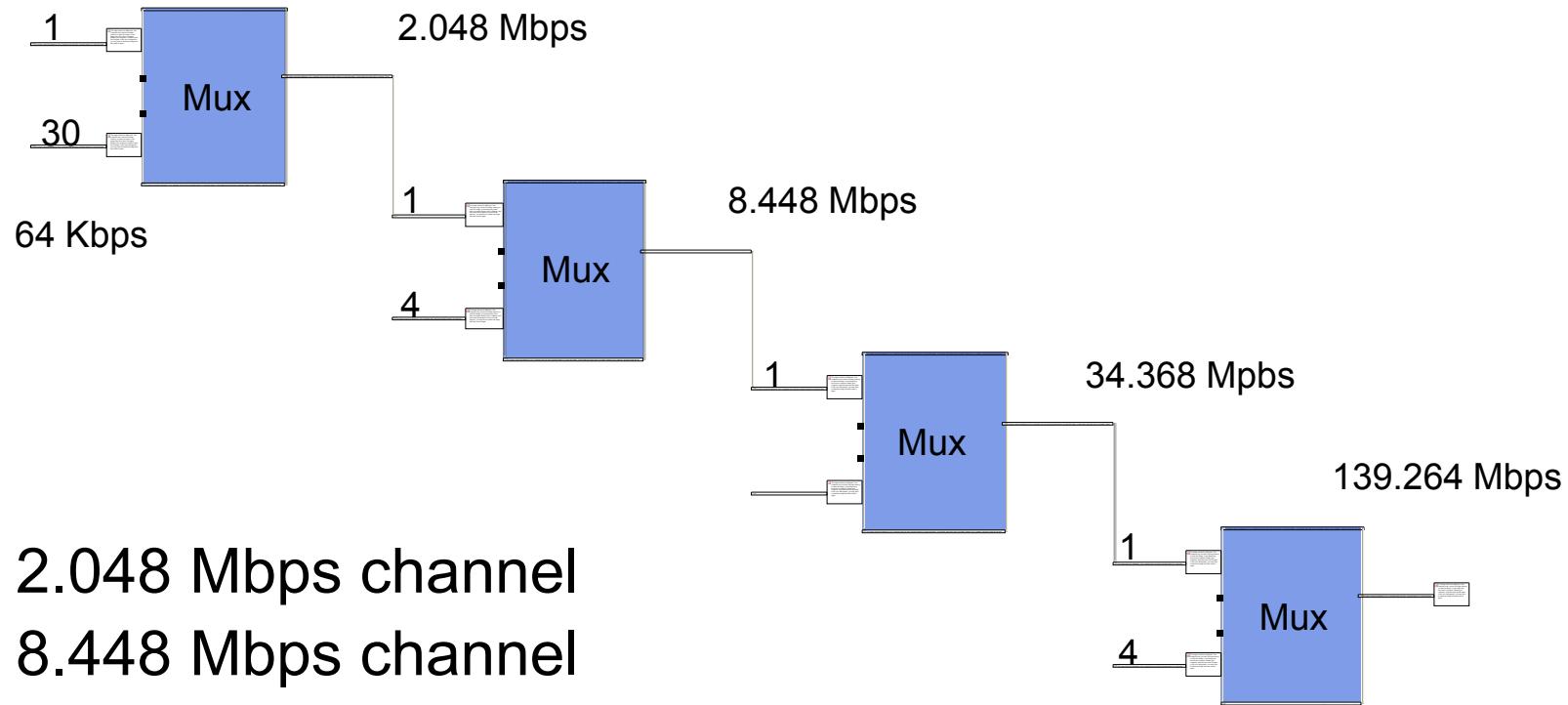


- DS0, 64 Kbps channel
- DS1, 1.544 Mbps channel
- DS2, 6.312 Mbps channel
- DS3, 44.736 Mbps channel
- DS4, 274.176 Mbps channel

# CCITT Digital Hierarchy



- CCITT digital hierarchy based on 30 PCM channels



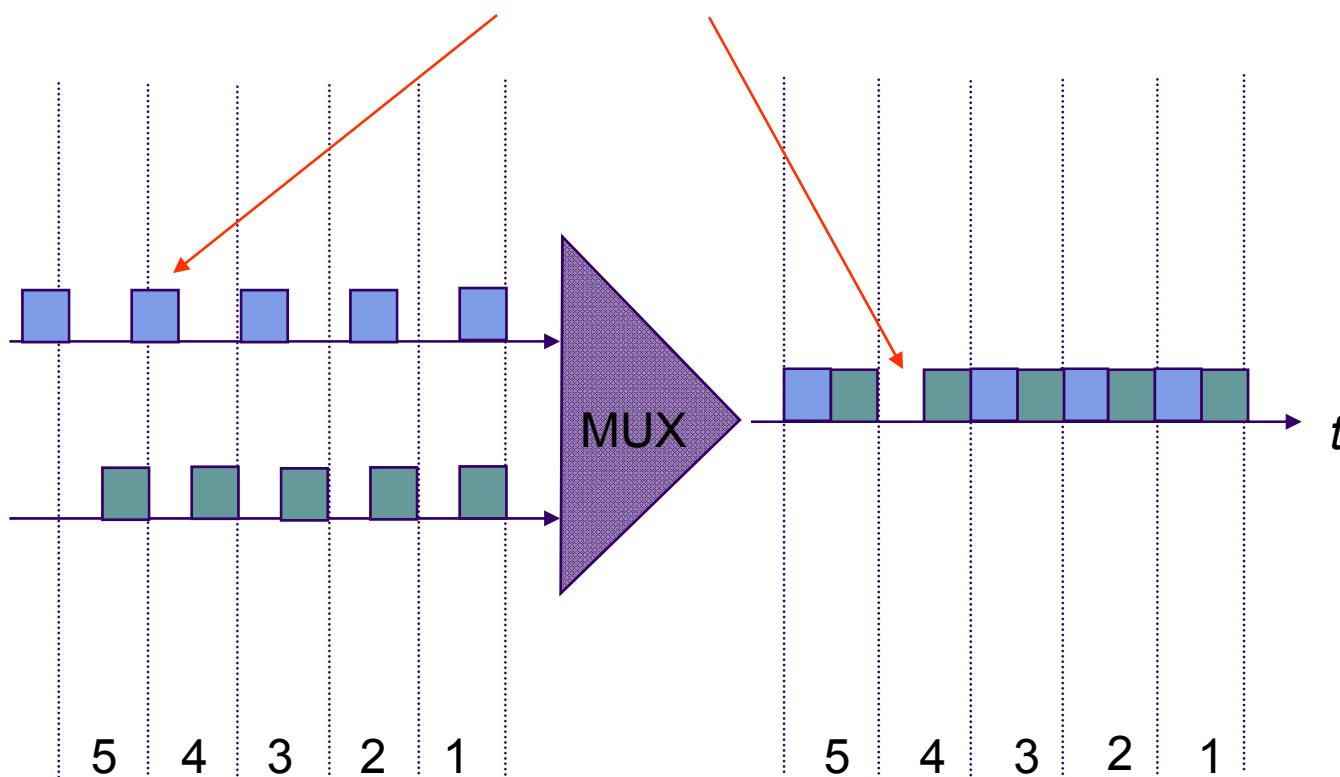
- E1, 2.048 Mbps channel
- E2, 8.448 Mbps channel
- E3, 34.368 Mbps channel
- E4, 139.264 Mbps channel

# Clock Synch & Bit Slips

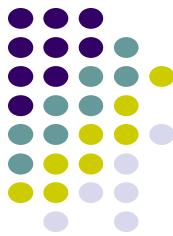


- Digital streams cannot be kept perfectly synchronized
- Bit slips can occur in multiplexers

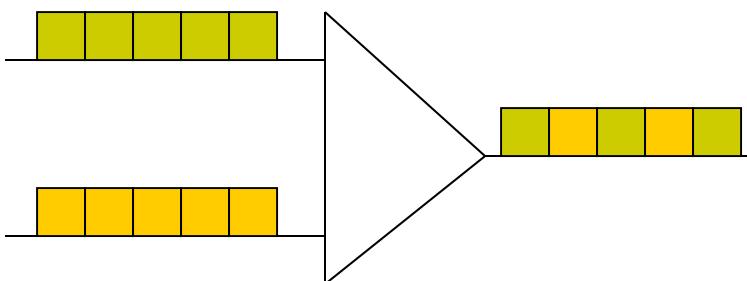
Slow clock results in late bit arrival and bit slip



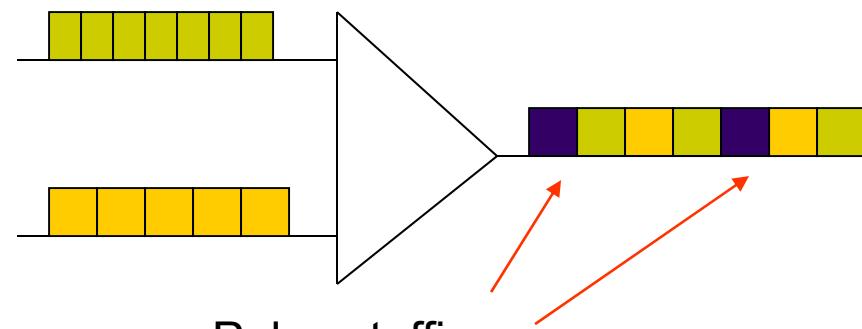
# Pulse Stuffing



- Pulse Stuffing: synchronization to avoid data loss due to slips
- Output rate  $> R_1 + R_2$ 
  - i.e. DS2,  $6.312\text{Mbps} = 4 \times 1.544\text{Mbps} + 136\text{ Kbps}$
- Pulse stuffing format
  - Fixed-length master frames with each channel allowed to stuff or not to stuff a single bit in the master frame.
  - Redundant stuffing specifications
  - signaling or specification bits (other than data bits) are distributed across a master frame.

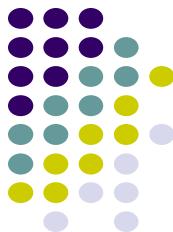


Muxing of equal-rate signals  
requires perfect synch

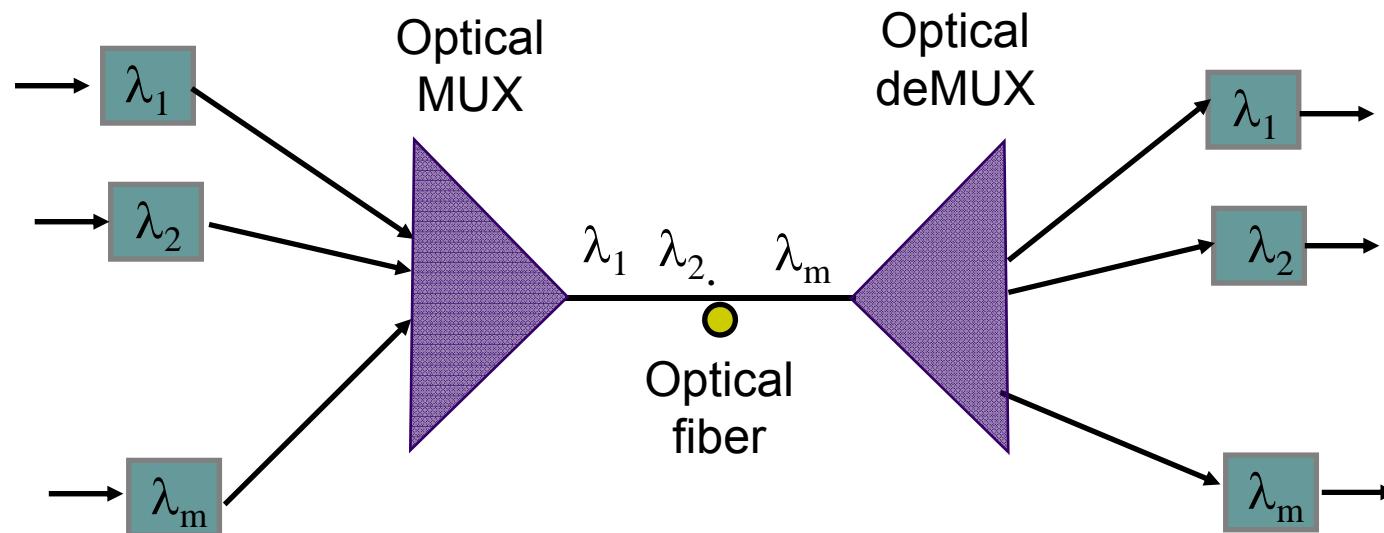


Pulse stuffing

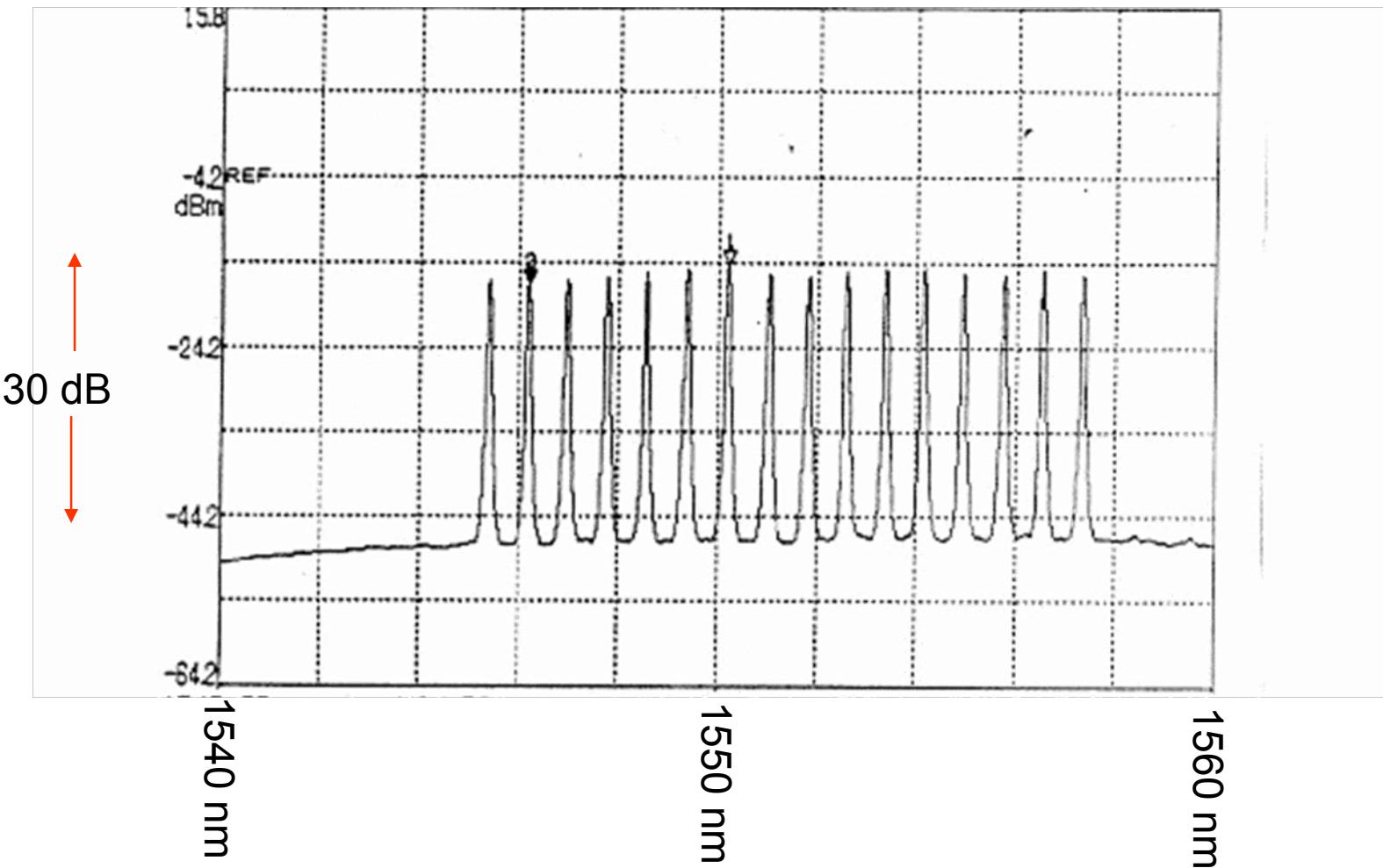
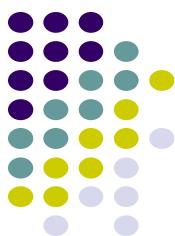
# Wavelength-Division Multiplexing



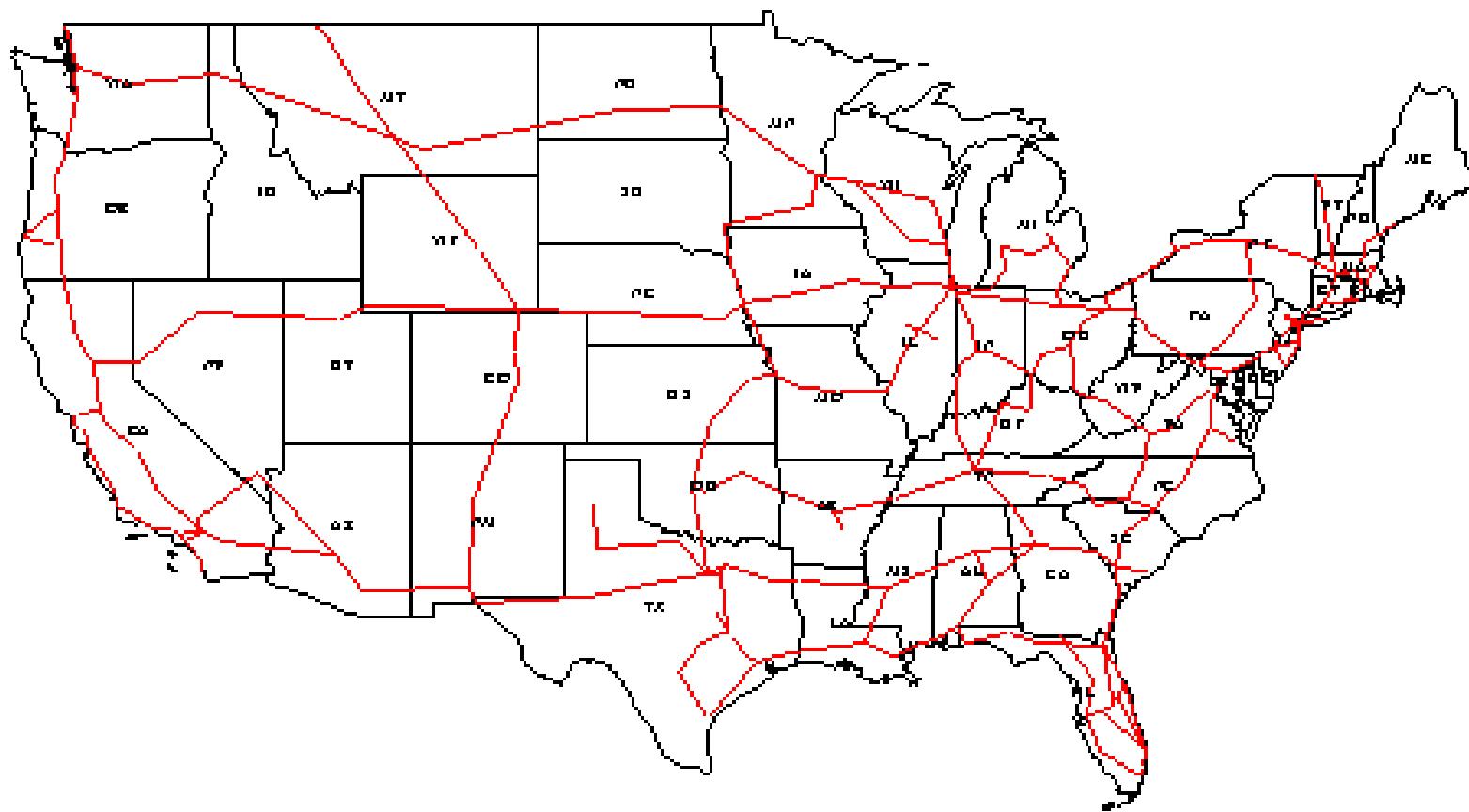
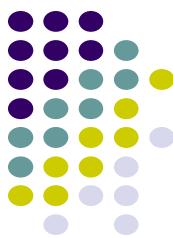
- Optical fiber link carries several wavelengths
  - From few (4-8) to many (64-160) wavelengths per fiber
- Imagine prism combining different colors into single beam
- Each wavelength carries a high-speed stream
  - Each wavelength can carry different format signal
  - e.g. 1 Gbps, 2.5 Gbps, or 10 Gbps



# Example: WDM with 16 wavelengths

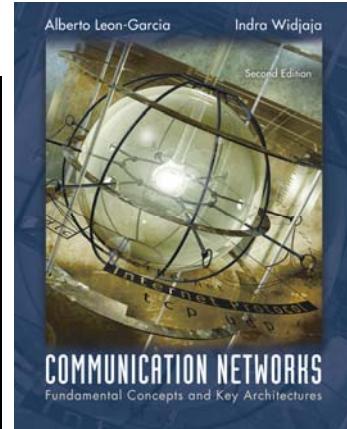


# Typical U.S. Optical Long-Haul Network

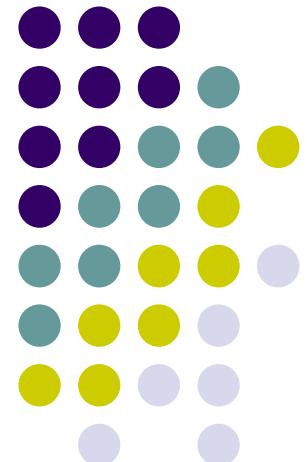


# Chapter 4

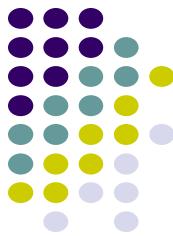
# Circuit-Switching Networks



***SONET***



# SONET: Overview

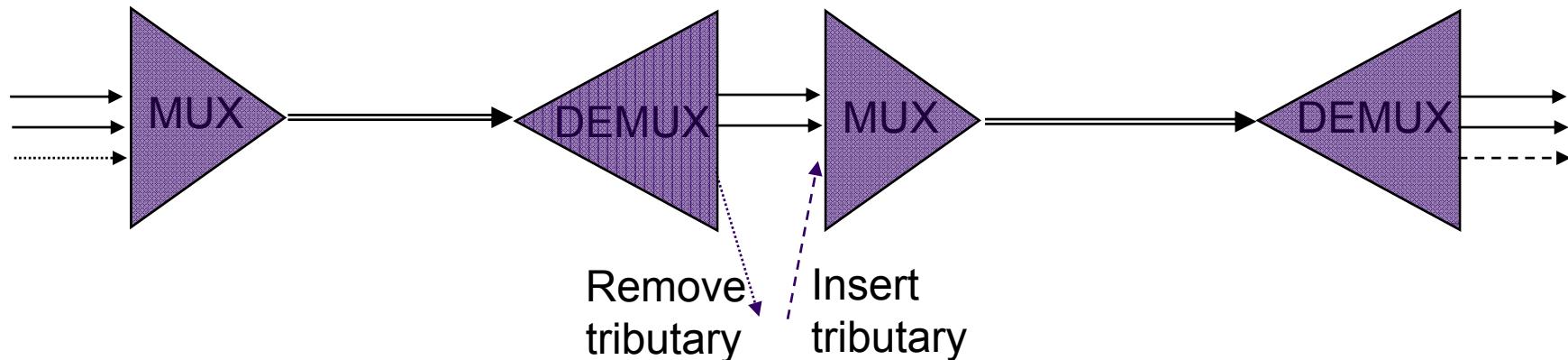


- *Synchronous Optical NETwork*
- North American TDM physical layer standard for optical fiber communications
- 8000 frames/sec. ( $T_{frame} = 125 \mu\text{sec}$ )
  - compatible with North American digital hierarchy
- SDH (Synchronous Digital Hierarchy) elsewhere
  - Needs to carry E1 and E3 signals
  - Compatible with SONET at higher speeds
- Greatly simplifies multiplexing in network backbone
- OA&M support to facilitate network management
- Protection & restoration

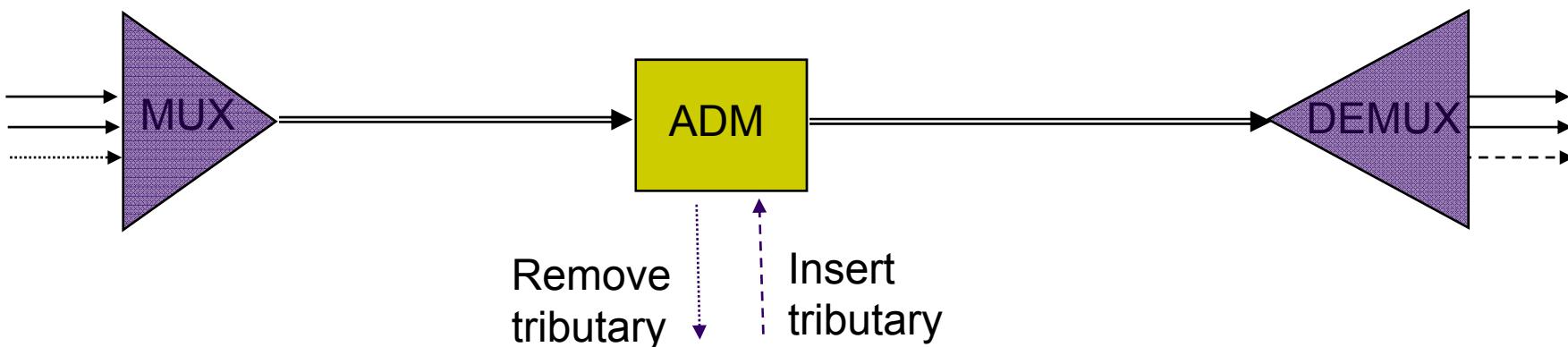
# SONET simplifies multiplexing



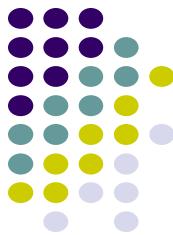
Pre-SONET multiplexing: Pulse stuffing required demultiplexing ***all*** channels



SONET Add-Drop Multiplexing: Allows taking individual channels in and out without full demultiplexing

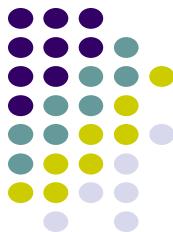


# SONET Specifications



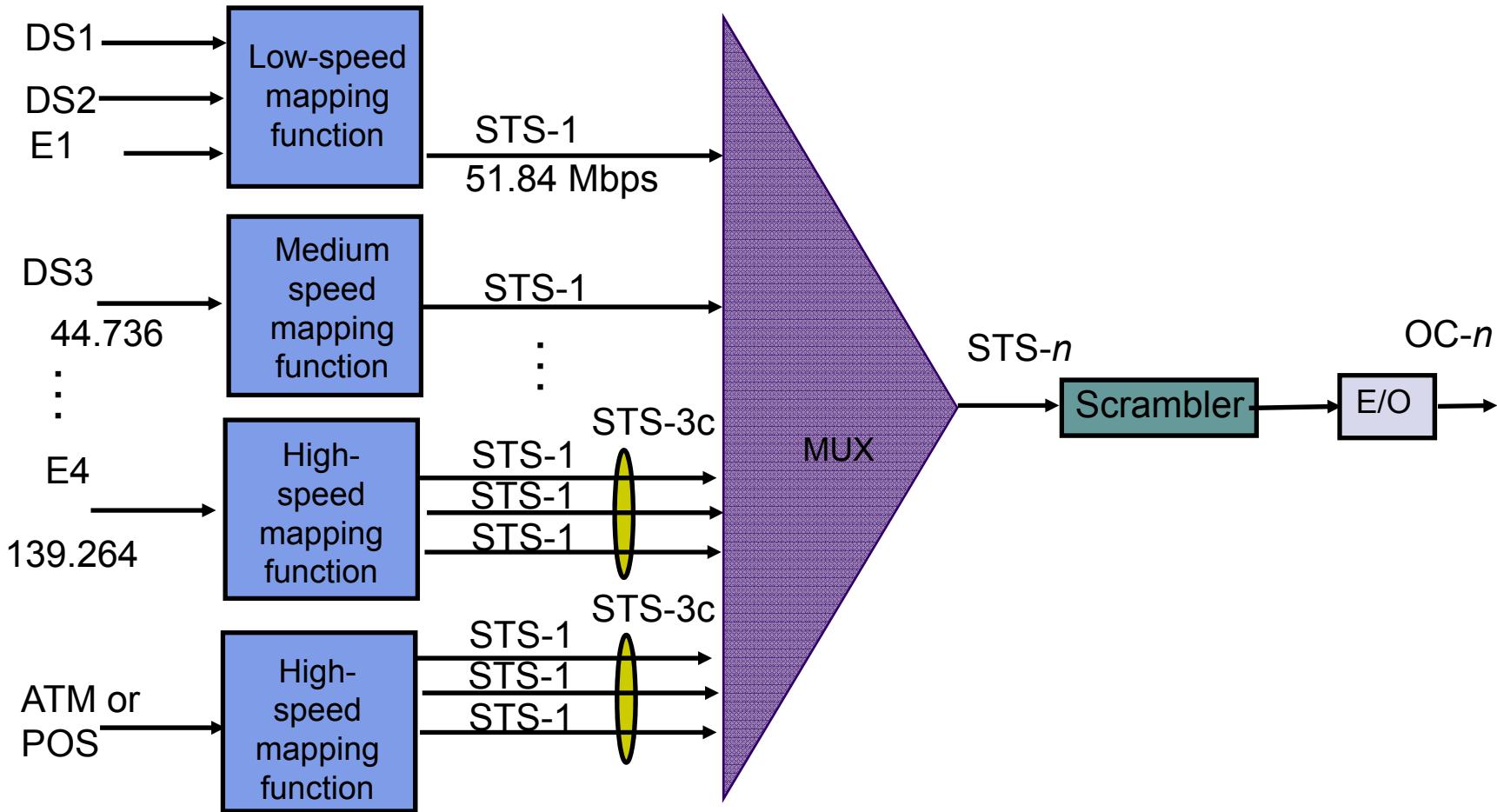
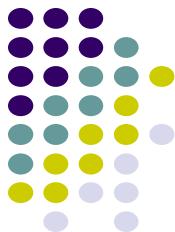
- Defines electrical & optical signal interfaces
- Electrical
  - Multiplexing, Regeneration performed in electrical domain
  - STS – Synchronous Transport Signals defined
  - Very short range (e.g., within a switch)
- Optical
  - Transmission carried out in optical domain
  - Optical transmitter & receiver
  - OC – Optical Carrier

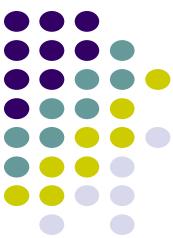
# SONET & SDH Hierarchy



<i><b>SONET Electrical Signal</b></i>	<i><b>Optical Signal</b></i>	<i><b>Bit Rate (Mbps)</b></i>	<i><b>SDH Electrical Signal</b></i>
STS-1	OC-1	51.84	N/A
STS-3	OC-3	155.52	STM-1
STS-9	OC-9	466.56	STM-3
STS-12	OC-12	622.08	STM-4
STS-18	OC-18	933.12	STM-6
STS-24	OC-24	1244.16	STM-8
STS-36	OC-36	1866.24	STM-12
STS-48	OC-48	2488.32	STM-16
STS-192	OC-192	9953.28	STM-64
<i>STS: Synchronous Transport Signal</i>	<i>OC: Optical Channel</i>		<i>STM: Synchronous Transfer Module</i>

# SONET Multiplexing

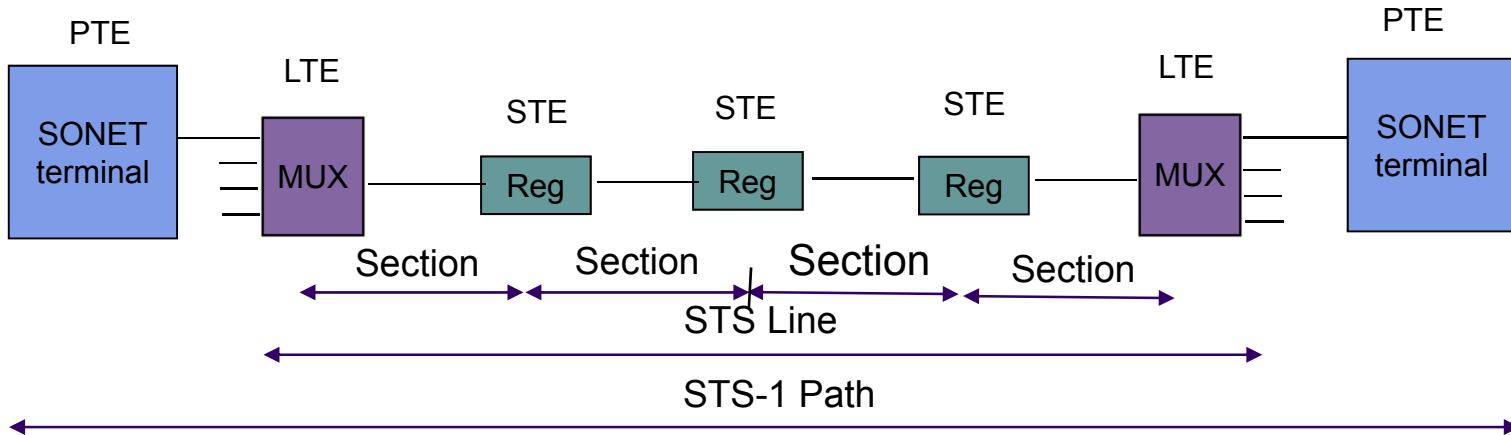




# SONET Equipment

- By Functionality
  - ADMs: dropping & inserting tributaries
  - Regenerators: digital signal regeneration
  - Cross-Connects: interconnecting SONET streams
- By Signaling between elements
  - Section Terminating Equipment (STE): span of fiber between adjacent devices, e.g. regenerators
  - Line Terminating Equipment (LTE): span between adjacent multiplexers, encompasses multiple sections
  - Path Terminating Equipment (PTE): span between SONET terminals at end of network, encompasses multiple lines

# Section, Line, & Path in SONET



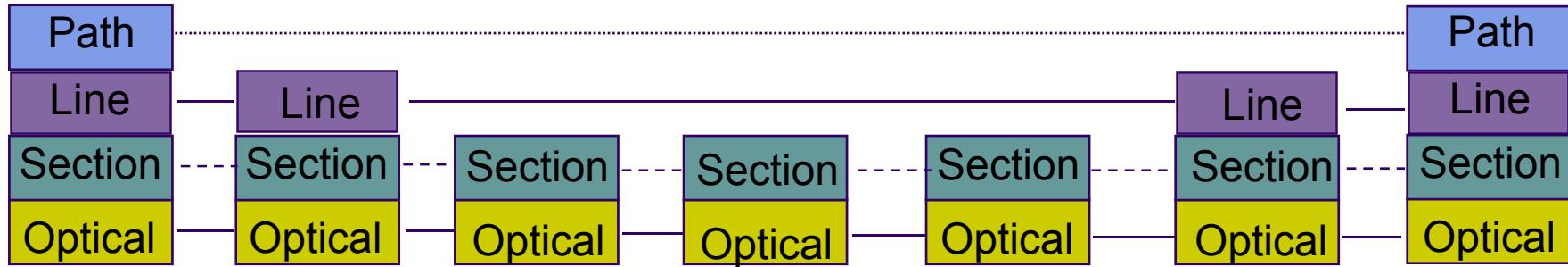
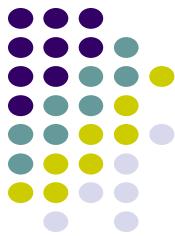
STE = Section Terminating Equipment, e.g., a repeater/regenerator

LTE = Line Terminating Equipment, e.g., a STS-1 to STS-3 multiplexer

PTE = Path Terminating Equipment, e.g., an STS-1 multiplexer

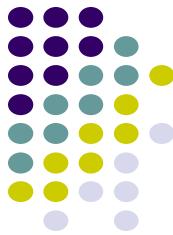
- Often, PTE and LTE equipment are the same
  - Difference is based on function and location
  - PTE is at the ends, e.g., STS-1 multiplexer.
  - LTE in the middle, e.g., STS-3 to STS-1 multiplexer.

# Section, Line, & Path Layers in SONET



- SONET has four layers
  - Optical, section, line, path
  - Each layer is concerned with the integrity of its own signals
- Each layer has its own protocols
  - SONET provides signaling channels for elements within a layer

# SONET STS Frame

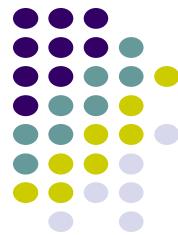


- SONET streams carry two types of overhead
- *Path overhead (POH)*:
  - inserted & removed at the ends
  - *Synchronous Payload Envelope (SPE)* consisting of Data + POH traverses network as a single unit
- *Transport Overhead (TOH)*:
  - processed at every SONET node
  - TOH occupies a portion of each SONET frame
  - TOH carries management & link integrity information

# STS-1 Frame

810x64kbps=51.84 Mbps

810 Octets per frame @ 8000 frames/sec



90 columns

1 Order of transmission  
2

A1	A2	J0	J1				
B1	E1	F1	B3				
D1	D2	D3	C2				
H1	H2	H3	G1				
B2	K1	K2	F2				
D4	D5	D6	H4				
D7	D8	D9	Z3				
D10	D11	D12	Z4				
S1	M0/1	E2	N1				

## Special OH octets:

**A1, A2** Frame Synch

**B1** Parity on Previous Frame  
(BER monitoring)

**J0** Section trace  
(Connection Alive?)

**H1, H2, H3** Pointer Action

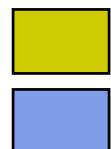
**K1, K2** Automatic Protection  
Switching

3 Columns of  
Transport OH

Synchronous Payload Envelope (SPE)  
1 column of Path OH + 8 data columns



Section Overhead



Path Overhead

Line Overhead

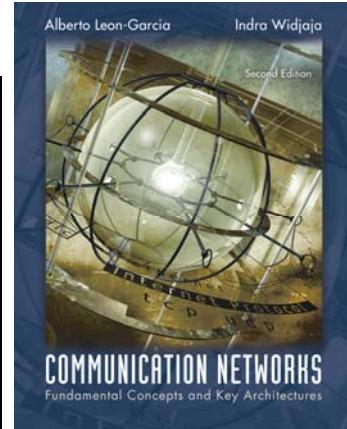


Data

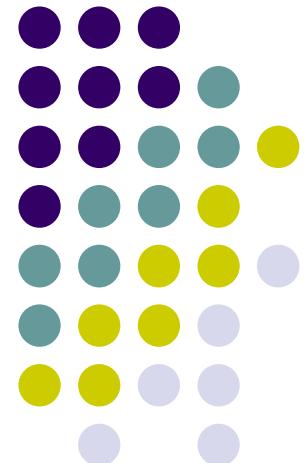
# Chapter 4

# Circuit-Switching

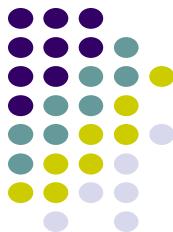
# Networks



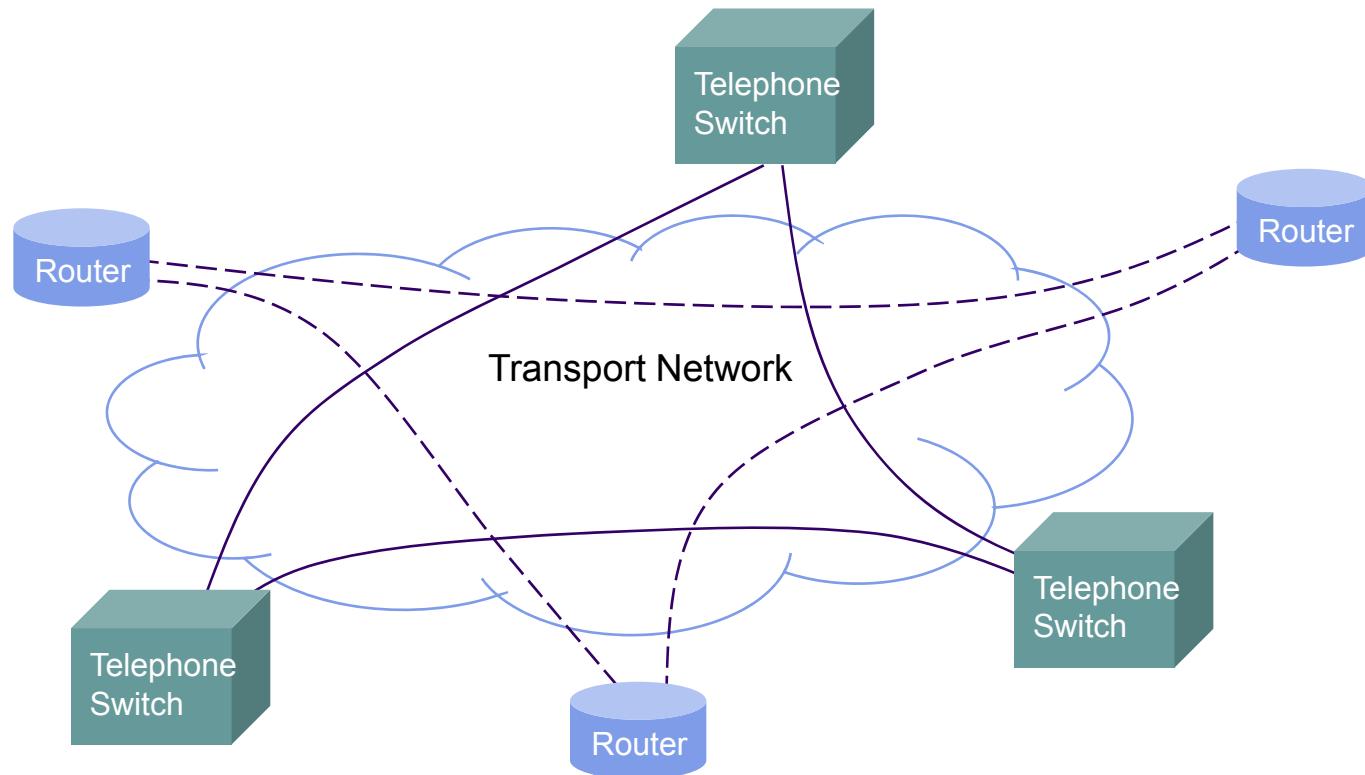
## *Transport Networks*



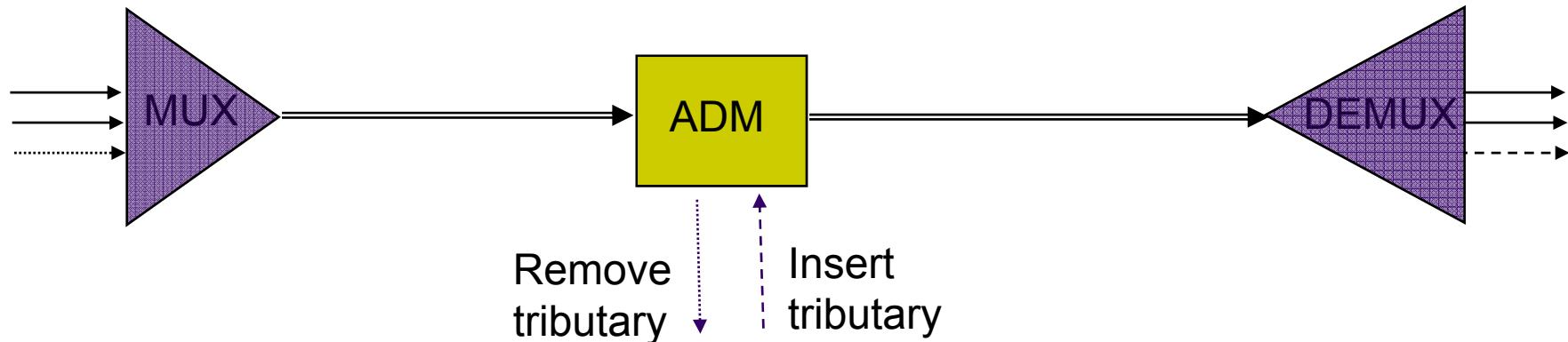
# Transport Networks



- Backbone of modern networks
- Provide high-speed connections: Typically STS-1 up to OC-192
- Clients: large routers, telephone switches, regional networks
- Very high reliability required because of consequences of failure
  - 1 STS-1 = 783 voice calls; 1 OC-48 = 32000 voice calls;



# SONET ADM Networks

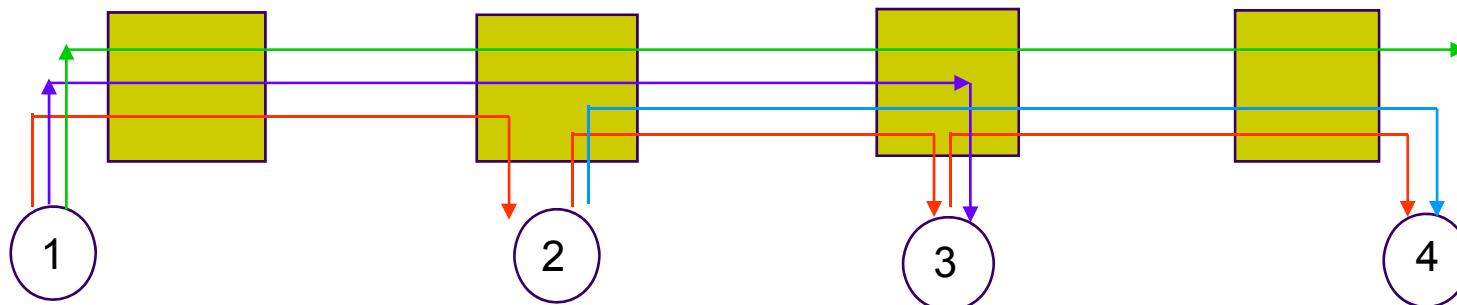


- SONET ADMs: the heart of existing transport networks
- ADMs interconnected in linear and ring topologies
- SONET signaling enables fast restoration (within 50 ms) of transport connections

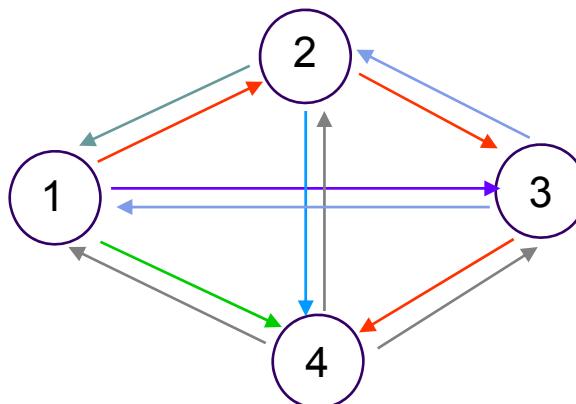
# Linear ADM Topology



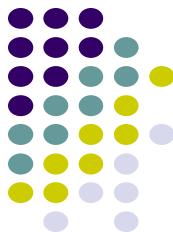
- ADMs connected in linear fashion
- Tributaries inserted and dropped to connect clients



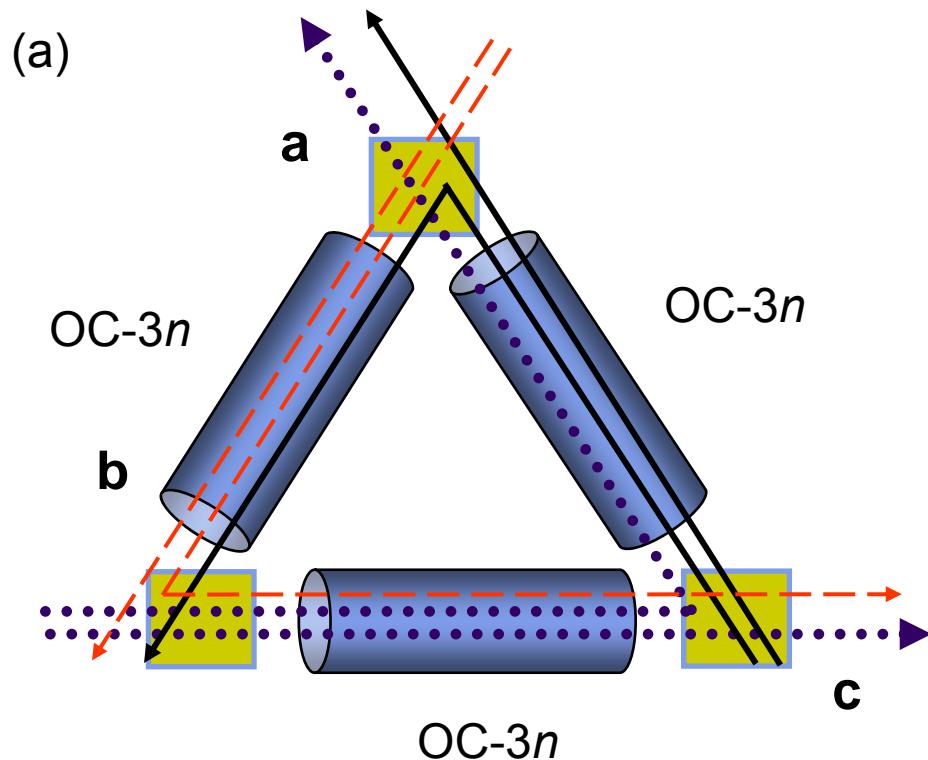
- Tributaries traverse ADMs transparently
- Connections create a *logical* topology seen by clients
- Tributaries from right to left are not shown



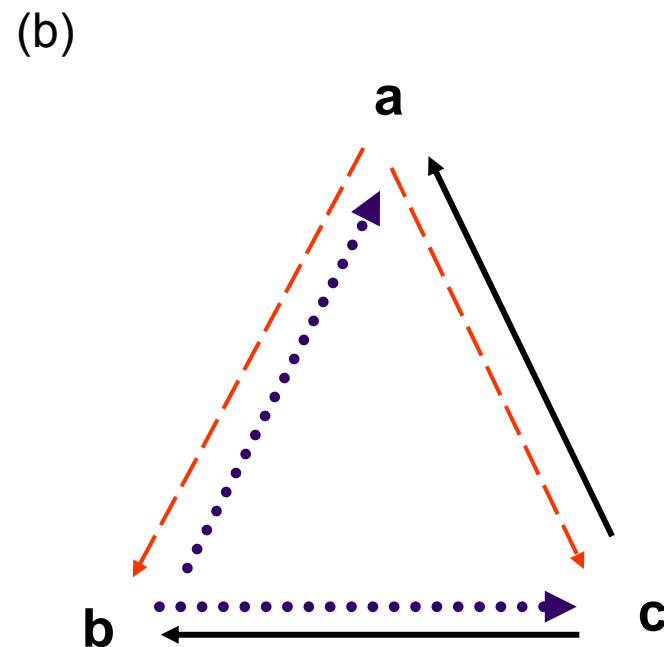
# SONET Rings



- ADMs can be connected in ring topology
- Clients see *logical* topology created by tributaries



Three ADMs connected in physical ring topology



Logical fully connected topology



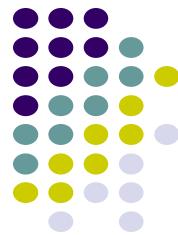
# SONET Ring Options

- 2 vs. 4 Fiber Ring Network
- Unidirectional vs. bidirectional transmission
- Path vs. Link protection

- Spatial capacity re-use & bandwidth efficiency
- Signalling requirements

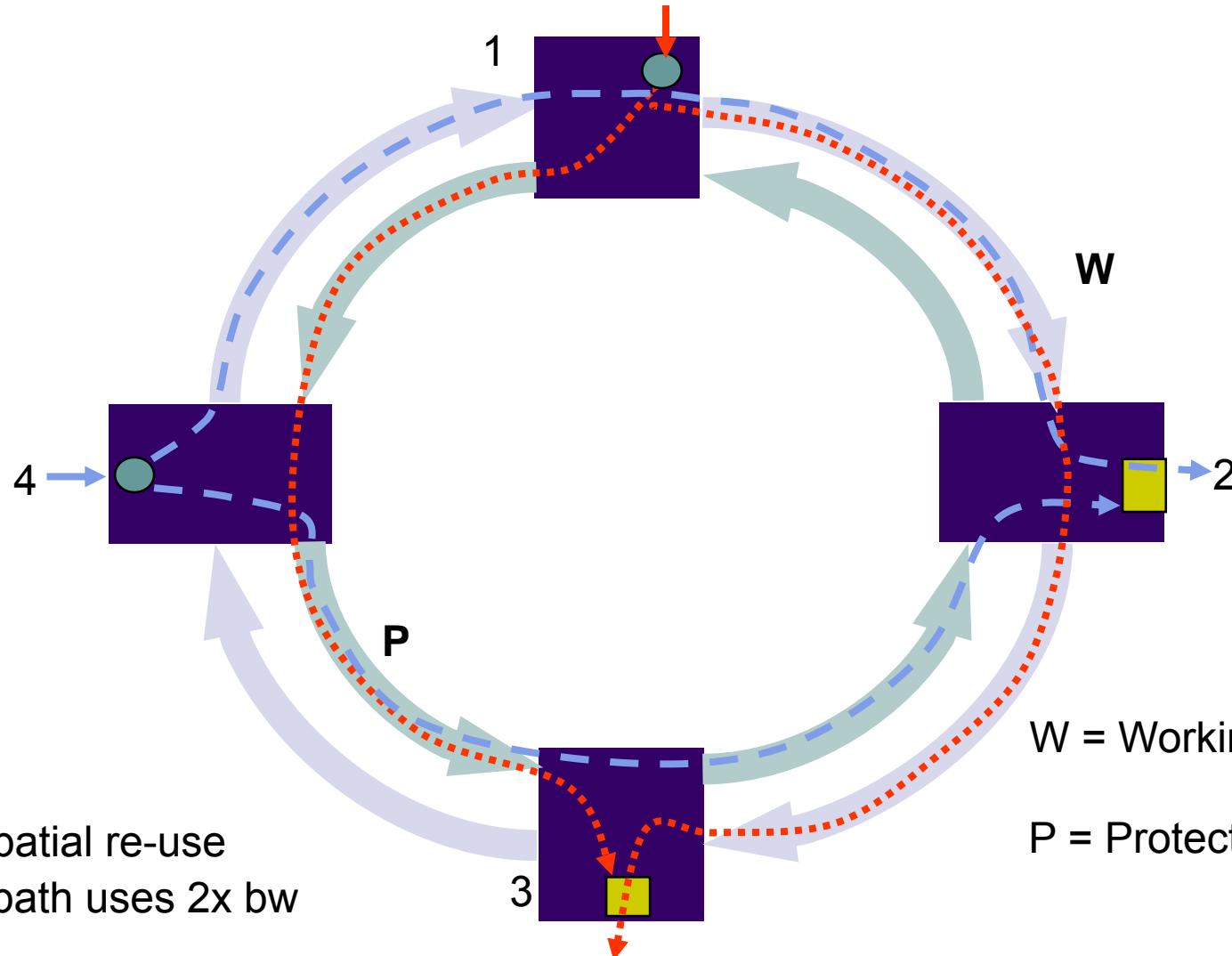
# Two-Fiber Unidirectional Path Switched Ring



Two fibers transmit in opposite directions

- Unidirectional
  - Working traffic flows clockwise
  - Protection traffic flows counter-clockwise
  - 1+1 like
- Selector at receiver does *path protection switching*

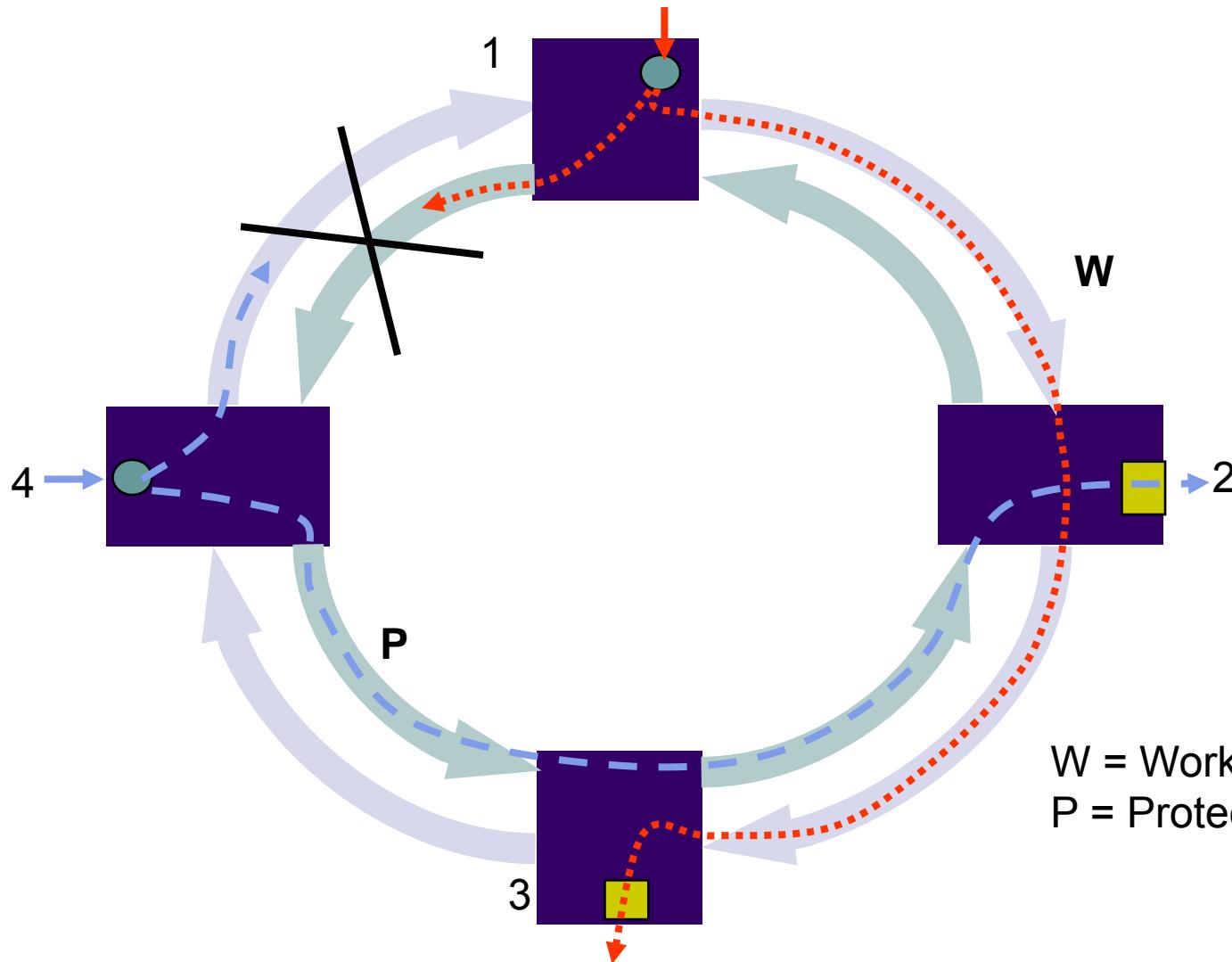
# UPSR



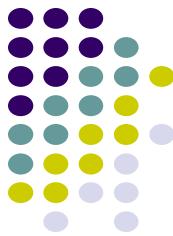
No spatial re-use  
Each path uses 2x bw

W = Working Paths  
P = Protection Paths

# UPSR path recovery

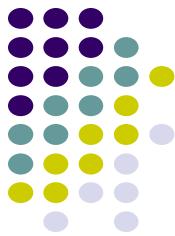


# UPSR Properties



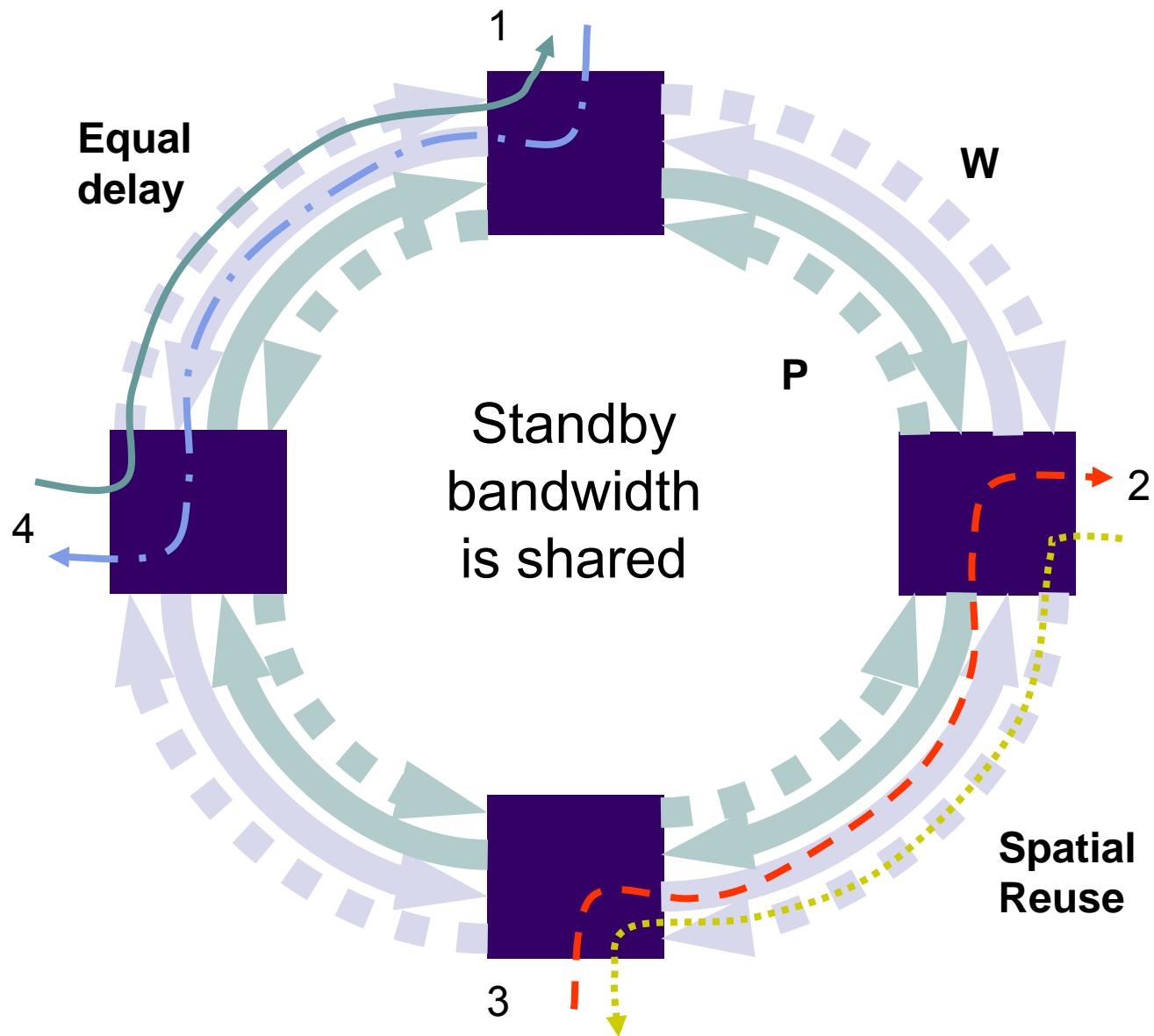
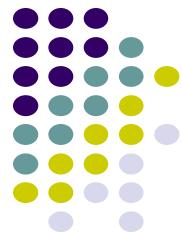
- Low complexity
- Fast path protection
- 2 TX, 2 RX
- No spatial re-use; ok for hub traffic pattern
- Suitable for lower-speed access networks
- Different delay between W and P path

# Four-Fiber Bidirectional Line Switched Ring

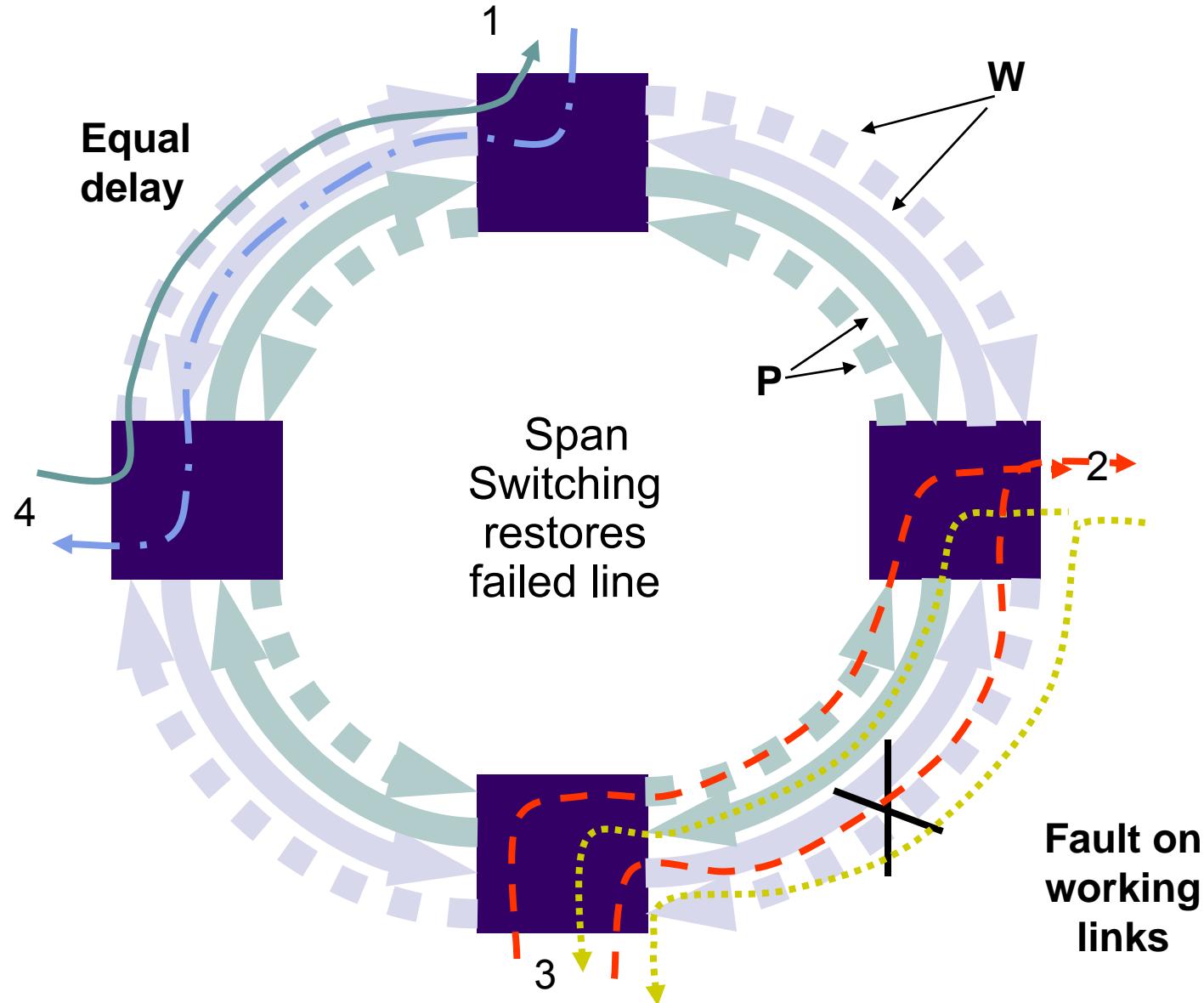


- 1 working fiber pair; 1 protection fiber pair
- Bidirectional
  - Working traffic & protection traffic use *same route* in working pair
  - 1:N like
- *Line* restoration provided by either:
  - Restoring a failed span
  - Switching the line around the ring

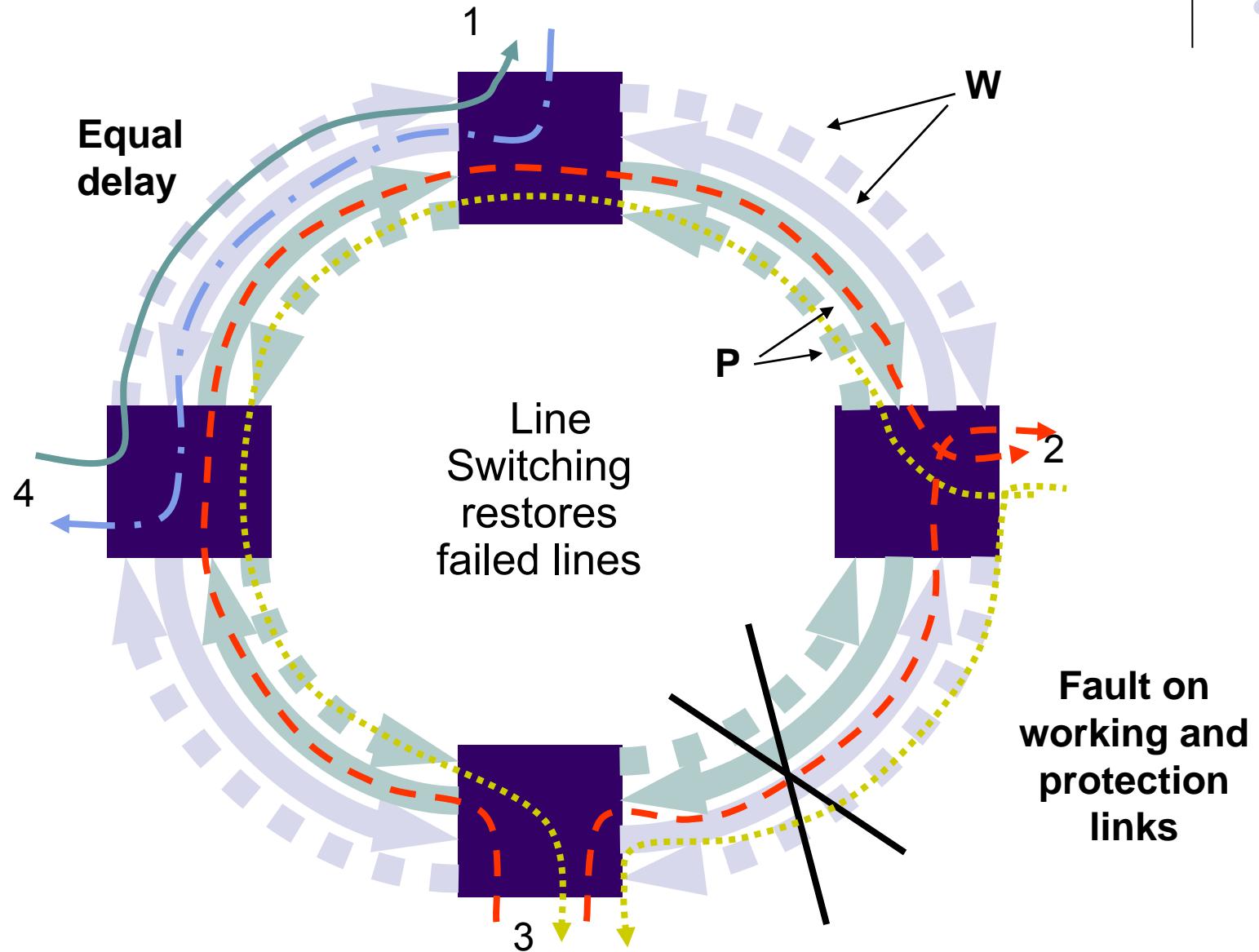
# 4-BLSR



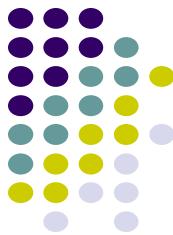
# BLSR Span Switching



# BLSR Span Switching

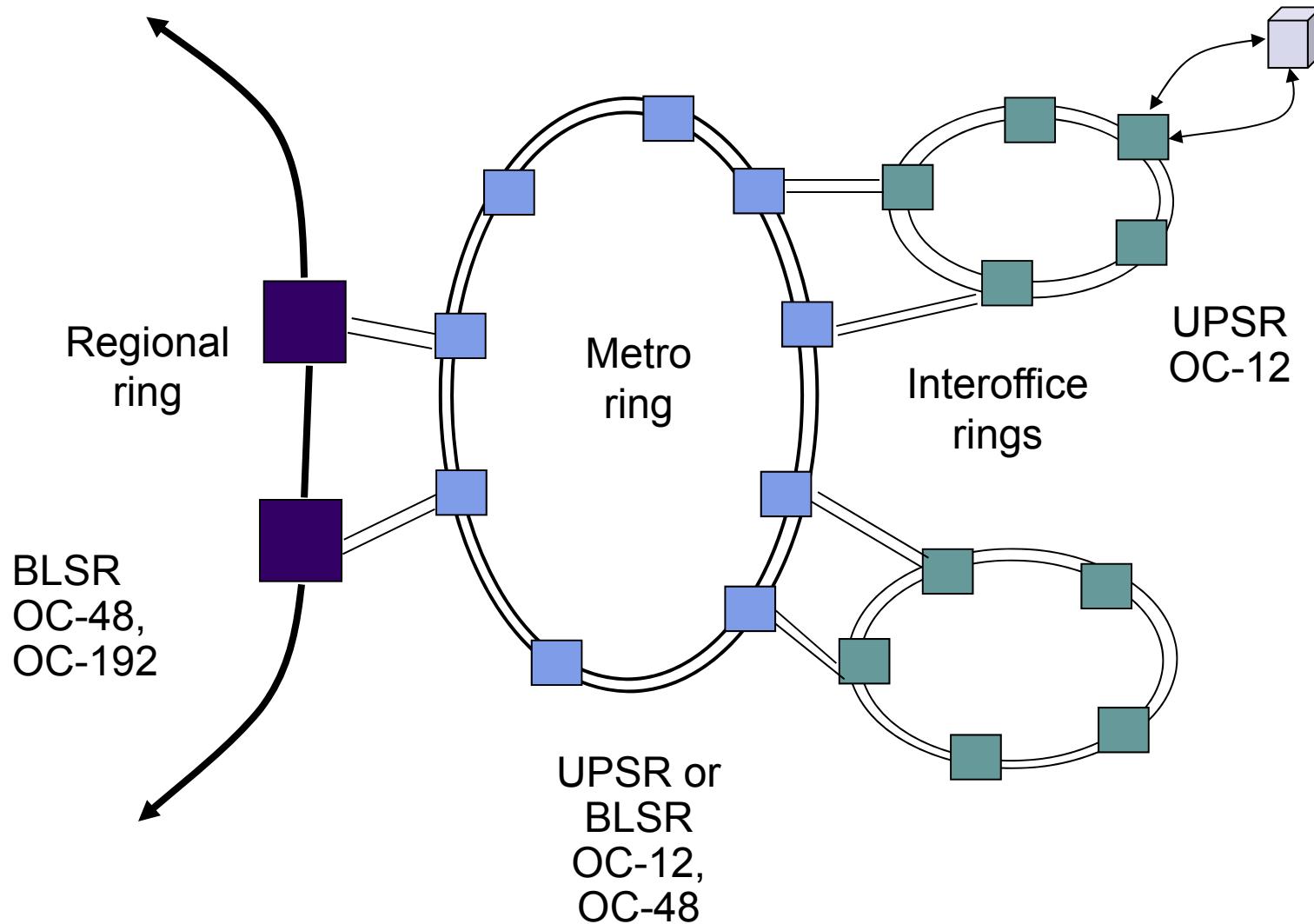
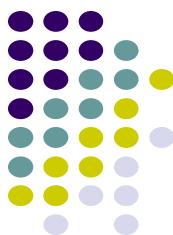


# 4-BLSR Properties

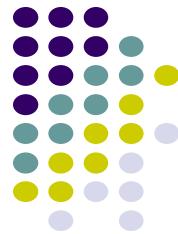


- High complexity: signalling required
- Fast line protection for restricted distance (1200 km) and number of nodes (16)
- 4 TX, 4 RX
- Spatial re-use; higher bandwidth efficiency
- Good for uniform traffic pattern
- Suitable for high-speed backbone networks
- Multiple simultaneous faults can be handled

# Backbone Networks consist of Interconnected Rings



# From SONET to WDM



## SONET

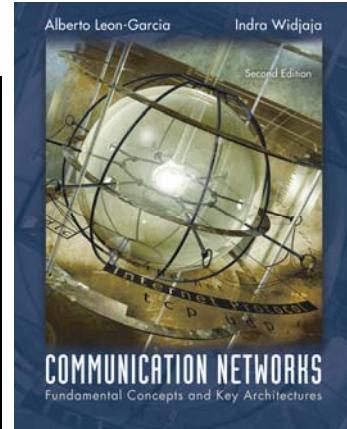
- combines multiple SPEs into high speed digital stream
- ADMs and crossconnects interconnected to form networks
- SPE paths between clients from logical topology
- High reliability through protection switching

## WDM

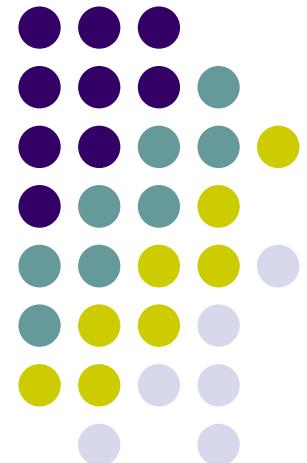
- combines multiple wavelengths into a common fiber
- Optical ADMs can be built to insert and drop wavelengths in same manner as in SONET ADMs
- Optical crossconnects can also be built
- All-optical backbone networks will provide end-to-end wavelength connections
- Protection schemes for recovering from failures are being developed to provide high reliability in all-optical networks

# Chapter 4

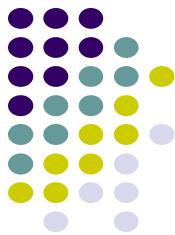
# Circuit-Switching Networks



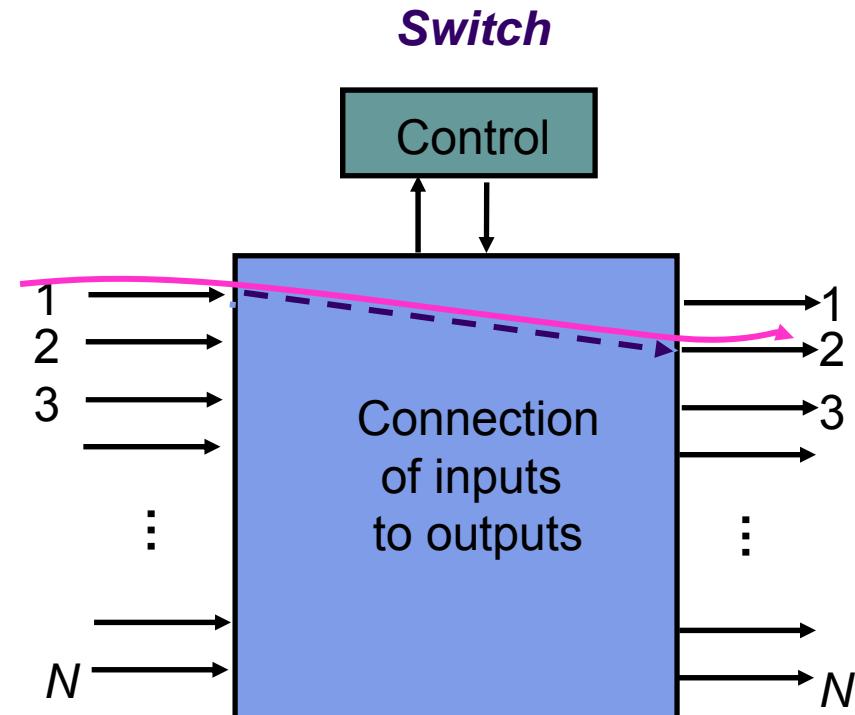
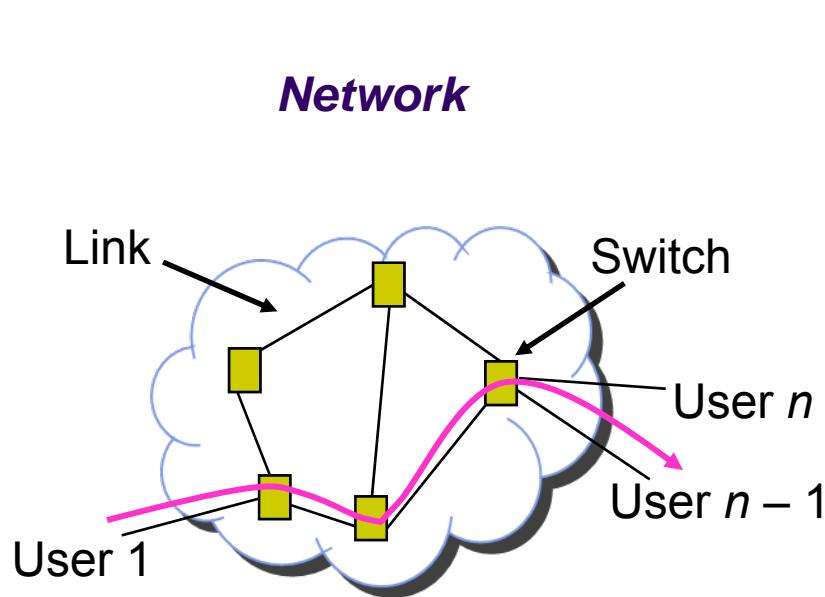
## *Circuit Switches*

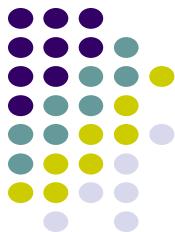


# Network: Links & switches



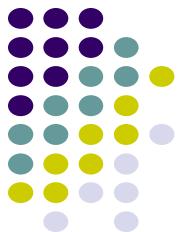
- Circuit consists of dedicated resources in sequence of links & switches across network
- *Circuit switch* connects input links to output links





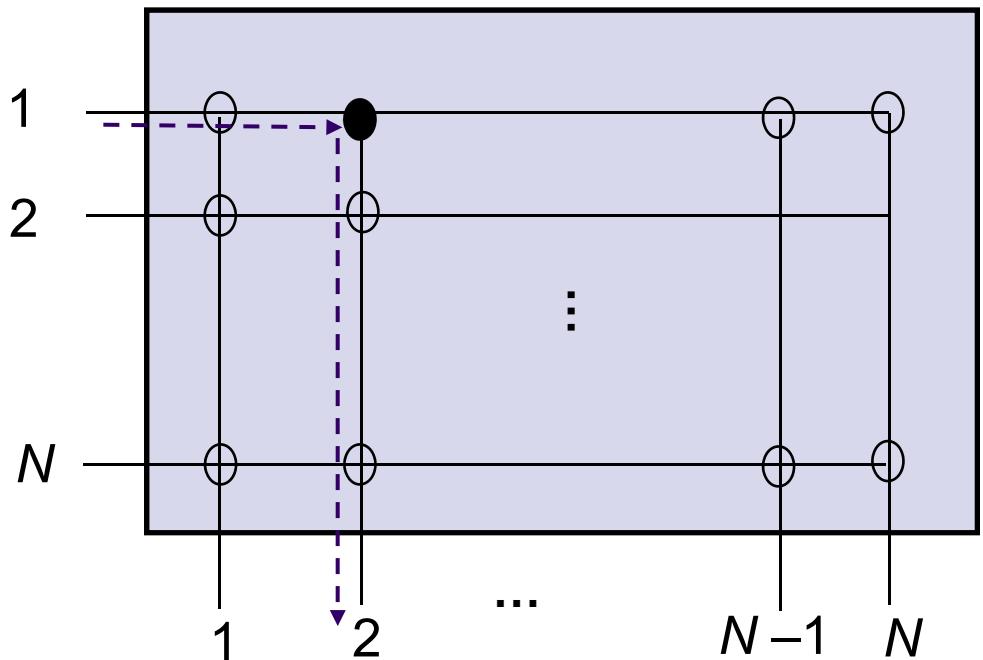
# Circuit Switch Types

- Space-Division switches
  - Provide separate physical connection between inputs and outputs
  - Crossbar switches
  - Multistage switches
- Time-Division switches
  - Time-slot interchange technique
  - Time-space-time switches
- Hybrids combine Time & Space switching



# Crossbar Space Switch

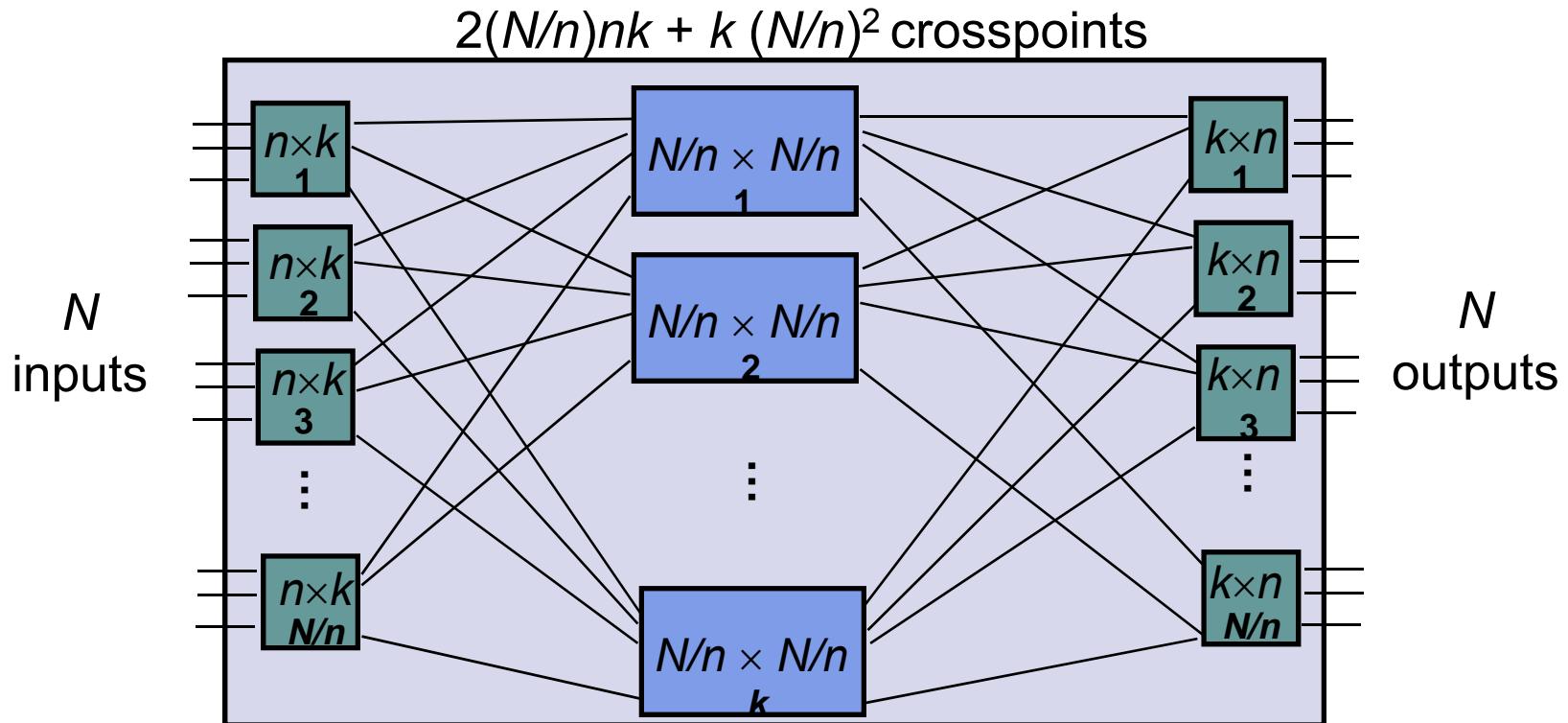
- $N \times N$  array of crosspoints
- Connect an input to an output by closing a crosspoint
- Nonblocking: Any input can connect to idle output
- Complexity:  $N^2$  crosspoints





# Multistage Space Switch

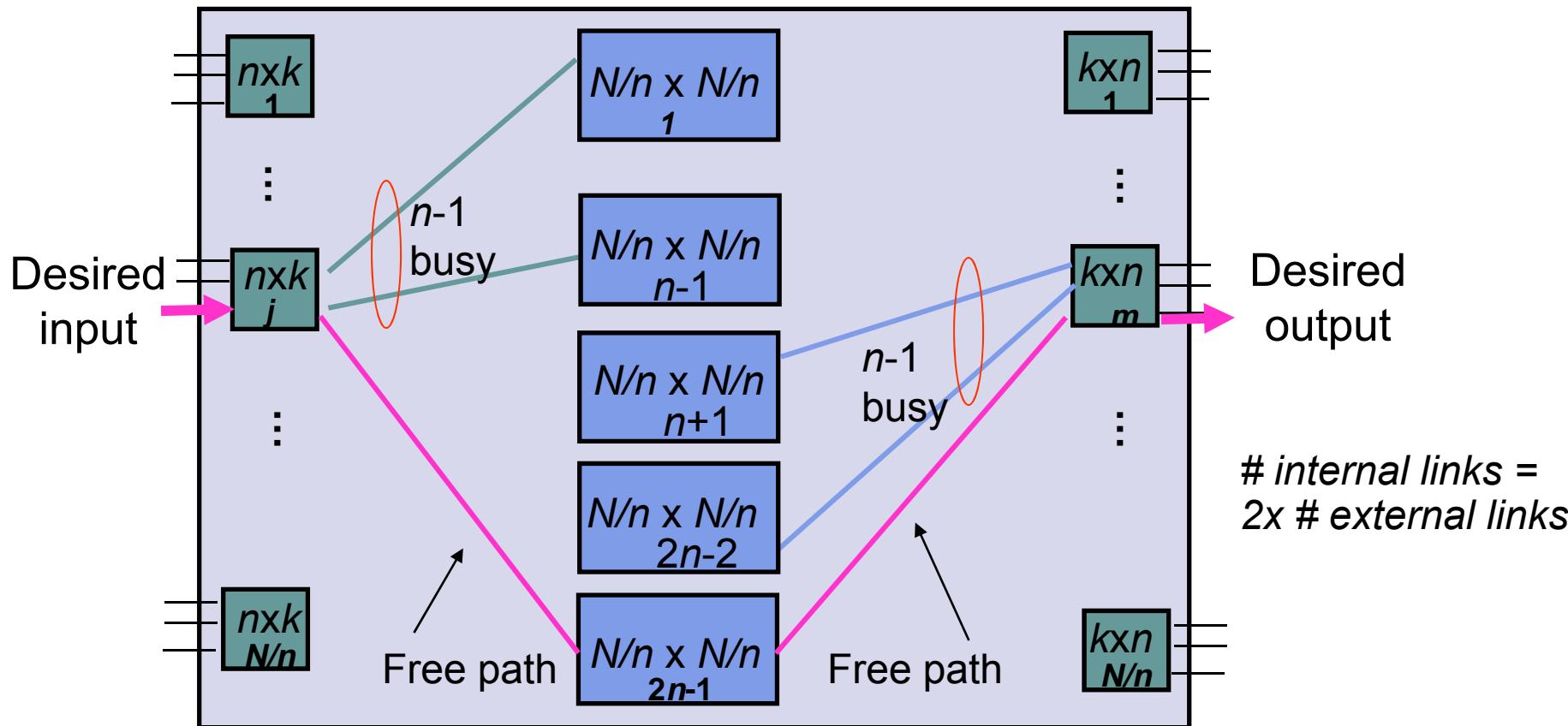
- Large switch built from multiple stages of small switches
- The  $n$  inputs to a first-stage switch share  $k$  paths through intermediate crossbar switches
- Larger  $k$  (more intermediate switches) means more paths to output
- In 1950s, Clos asked, “How many intermediate switches required to make switch nonblocking?”



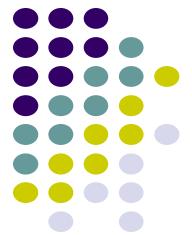
# Clos Non-Blocking Condition: $k=2n-1$



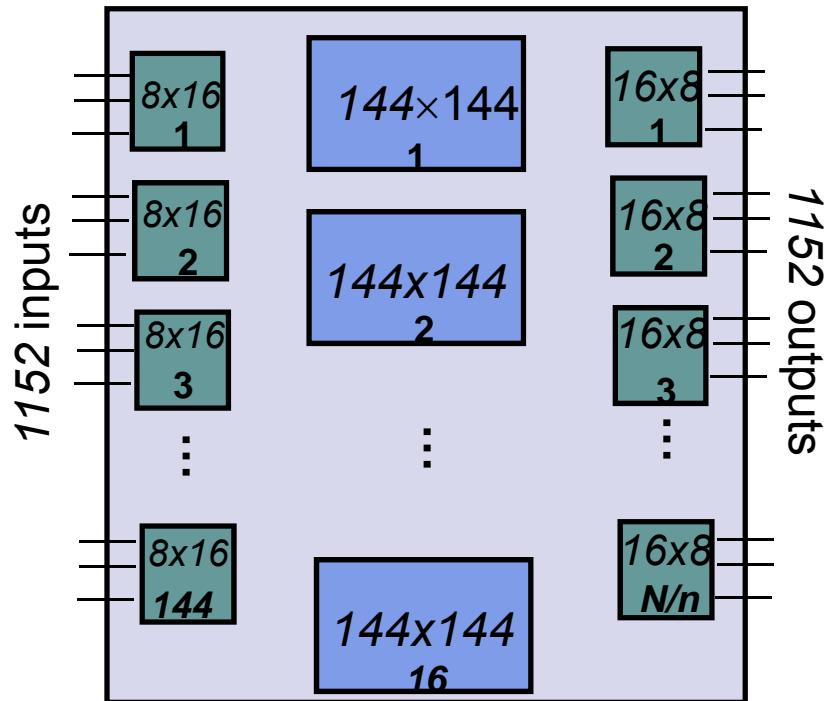
- Request connection from last input to input switch  $j$  to last output in output switch  $m$
- Worst Case: All other inputs have seized top  $n-1$  middle switches AND all other outputs have seized next  $n-1$  middle switches
- If  $k=2n-1$ , there is another path left to connect desired input to desired output



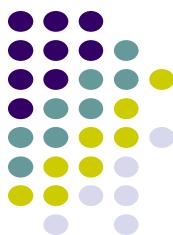
# Example: Clos Switch Design



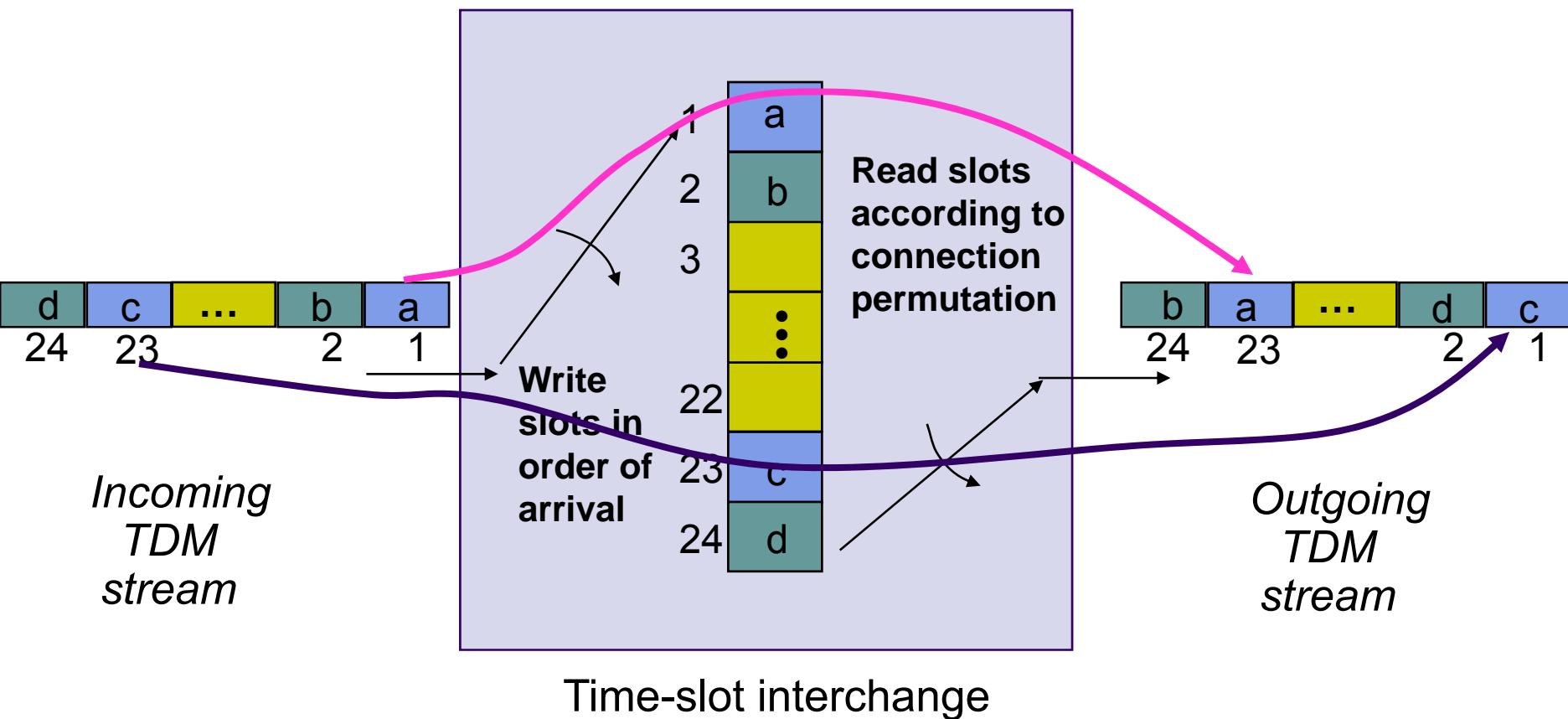
- Circa 2002, Mindspeed offered a Crossbar chip with the following specs:
  - 144 inputs x 144 outputs, 3.125 Gbps/line
  - Aggregate Crossbar chip throughput: 450 Gbps
- Clos Nonblocking Design for 1152x1152 switch
  - $N=1152$ ,  $n=8$ ,  $k=16$
  - $N/n=144$  8x16 switches in first stage
  - 16 144x144 in centre stage
  - 144 16x8 in third stage
  - Aggregate Throughput: 3.6 Tbps!
  - Note: the 144x144 crossbar can be partitioned into multiple smaller switches

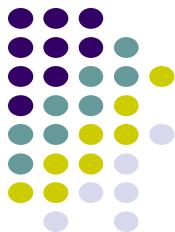


# Time-Slot Interchange (TSI) Switching



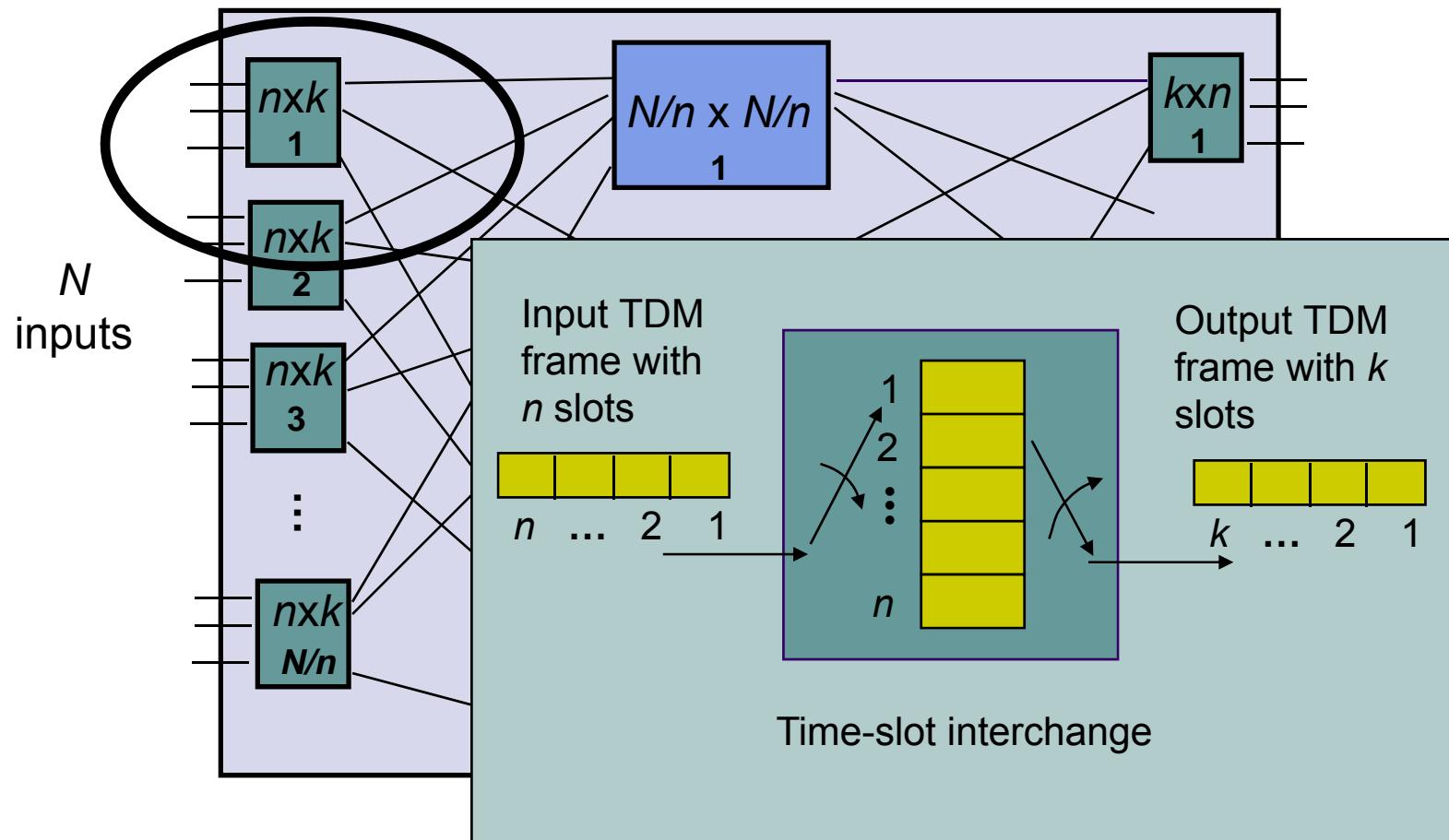
- Write bytes from arriving TDM stream into memory
- Read bytes in permuted order into outgoing TDM stream
- Max # slots =  $125 \mu\text{sec} / (2 \times \text{memory cycle time})$





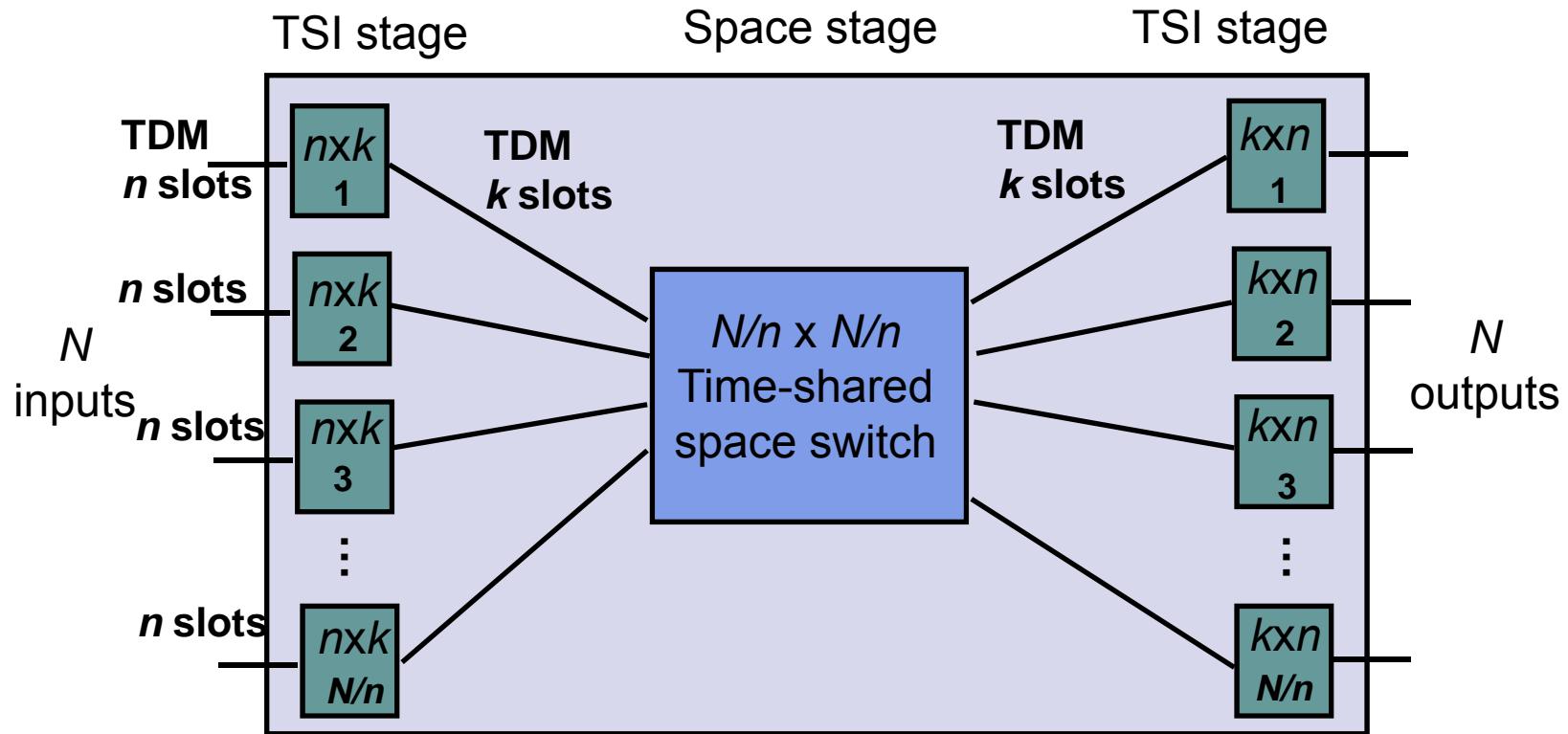
# Time-Space-Time Hybrid Switch

- Use TSI in first & third stage; Use crossbar in middle
- Replace  $n$  input  $\times k$  output space switch by TSI switch that takes  $n$ -slot input frame and switches it to  $k$ -slot output frame



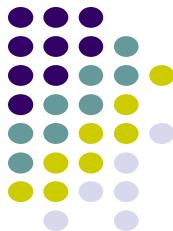


# Time-Share the Crossbar Switch



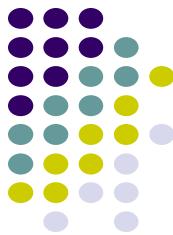
- Interconnection pattern of space switch is reconfigured every time slot
- Very compact design: fewer lines because of TDM & less space because of time-shared crossbar

# Available TSI Chips circa 2002



- OC-192 SONET Framer Chips
  - Decompose 192 STS1s and perform (restricted) TSI
- Single-chip TST
  - 64 inputs x 64 outputs
  - Each line @ STS-12 (622 Mbps)
  - Equivalent to 768x768 STS-1 switch

# Pure Optical Switching

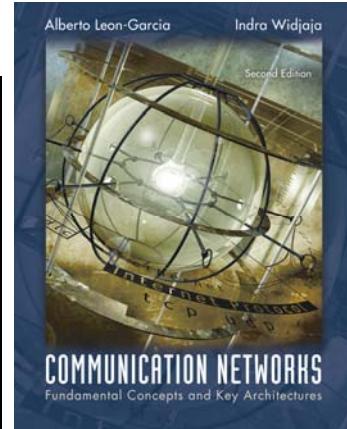


- Pure Optical switching: light-in, light-out, without optical-to-electronic conversion
- Space switching theory can be used to design optical switches
  - Multistage designs using small optical switches
  - Typically 2x2 or 4x4
  - MEMs and Electro-optic switching devices
- Wavelength switches
  - Very interesting designs when space switching is combined with wavelength conversion devices

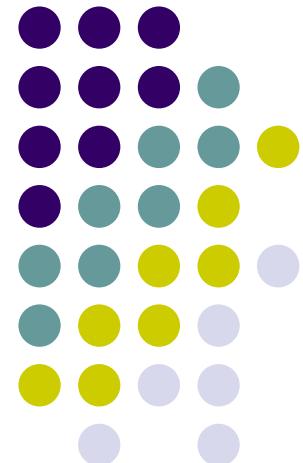
# Chapter 4

# Circuit-Switching

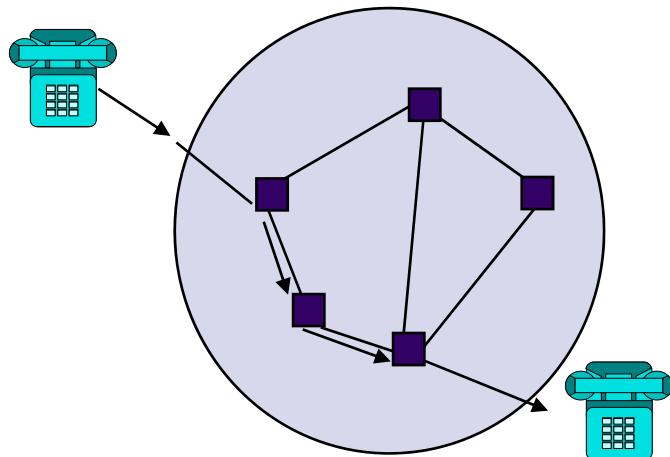
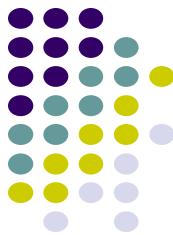
# Networks



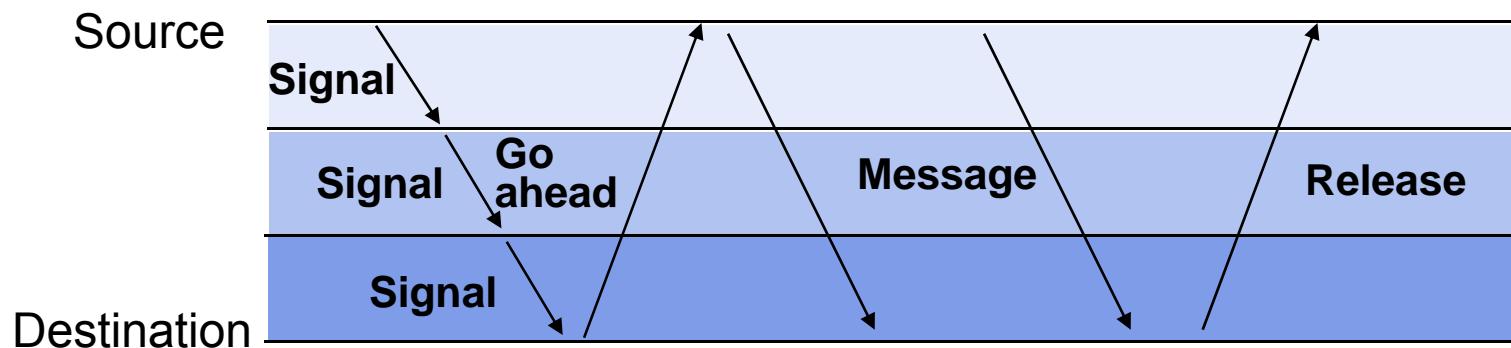
## *The Telephone Network*



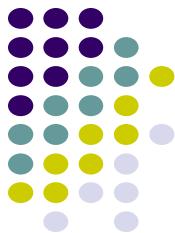
# Telephone Call



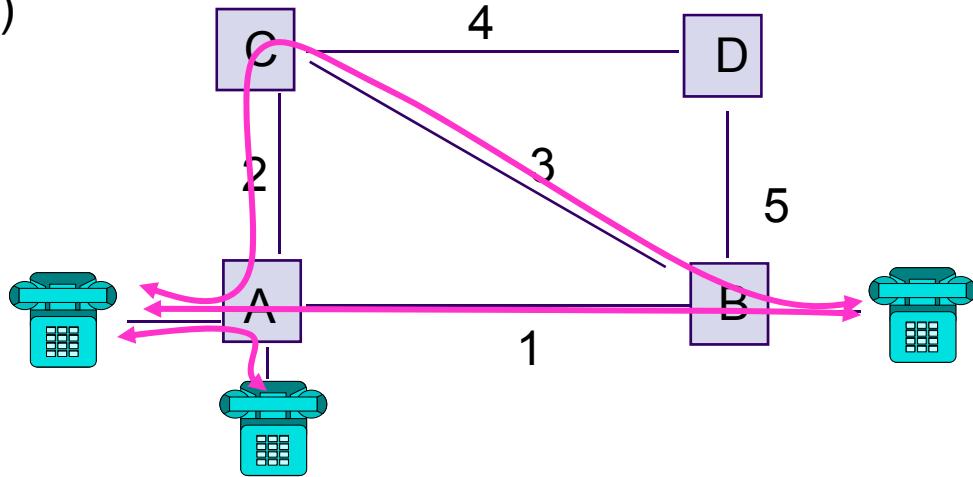
- User requests connection
- Network signaling establishes connection
- Speakers converse
- User(s) hang up
- Network releases connection resources



# Call Routing



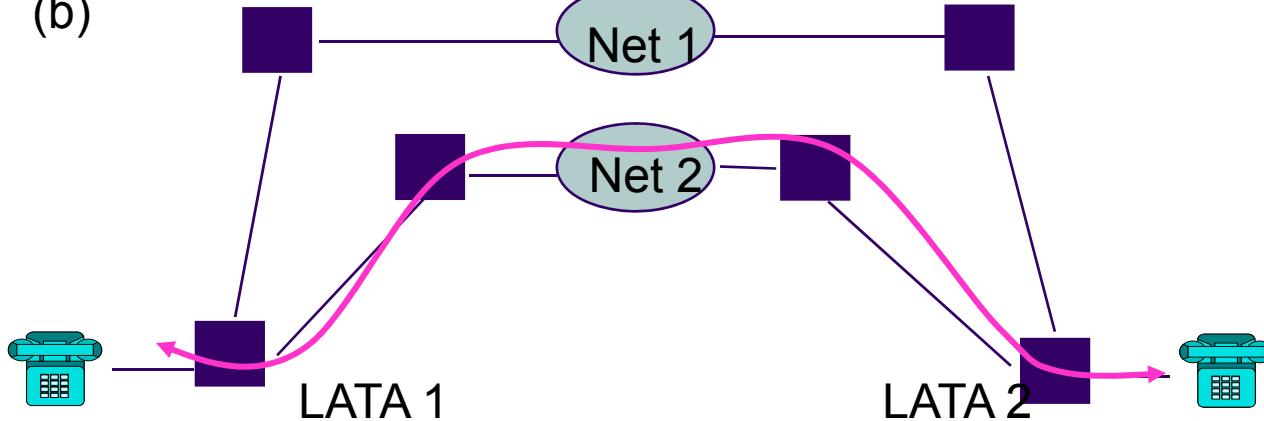
(a)



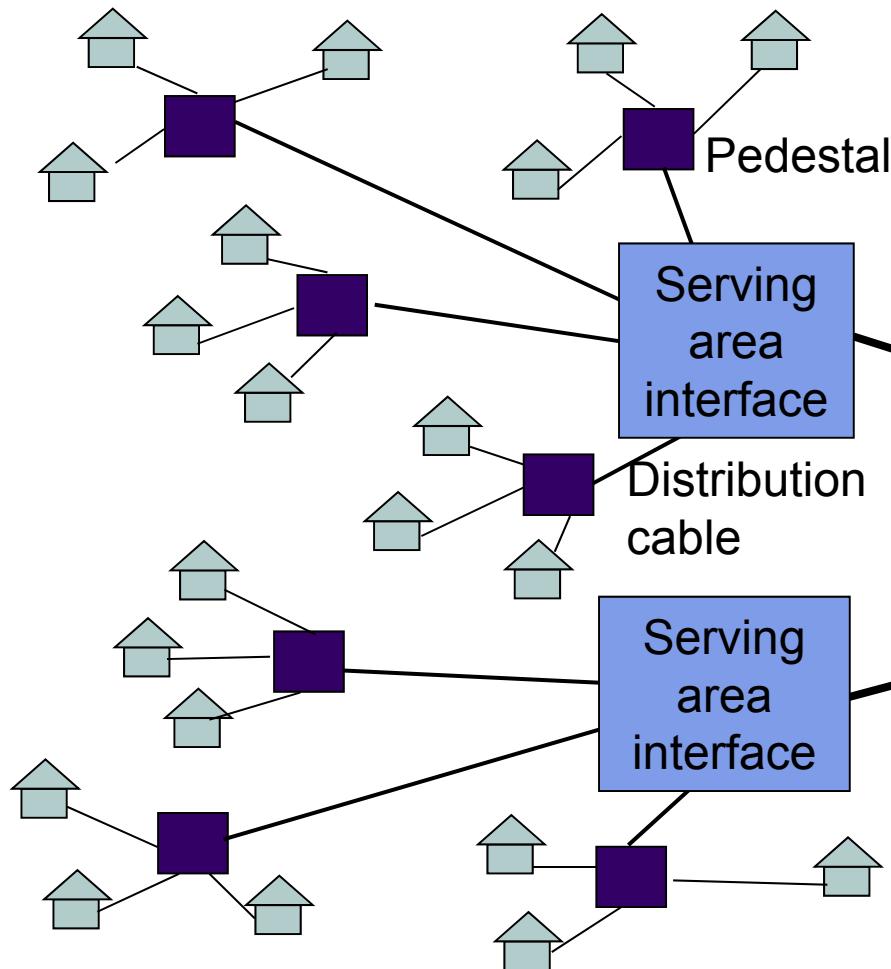
- Local calls routed through local network (In U.S. Local Access & Transport Area)

- Long distance calls routed to long distance service provider

(b)



# Telephone Local Loop



## Local Loop: “Last Mile”

- Copper pair from telephone to CO
- Pedestal to SAI to Main Distribution Frame (MDF)
- 2700 cable pairs in a feeder cable
- MDF connects
  - voice signal to telephone switch
  - DSL signal to routers

For interesting pictures of switches & MDF, see  
[web.mit.edu/is/is/delivery/5ess/photos.html](http://web.mit.edu/is/is/delivery/5ess/photos.html)  
[www.museumofcommunications.org/coe.html](http://www.museumofcommunications.org/coe.html)

# Fiber-to-the-Home or Fiber-to-the-Curve?

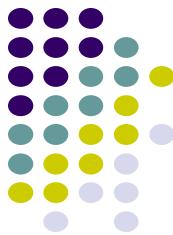
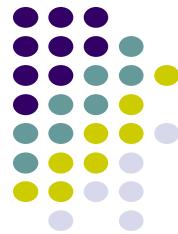


Table 3.5 Data rates of 24-gauge twisted pair

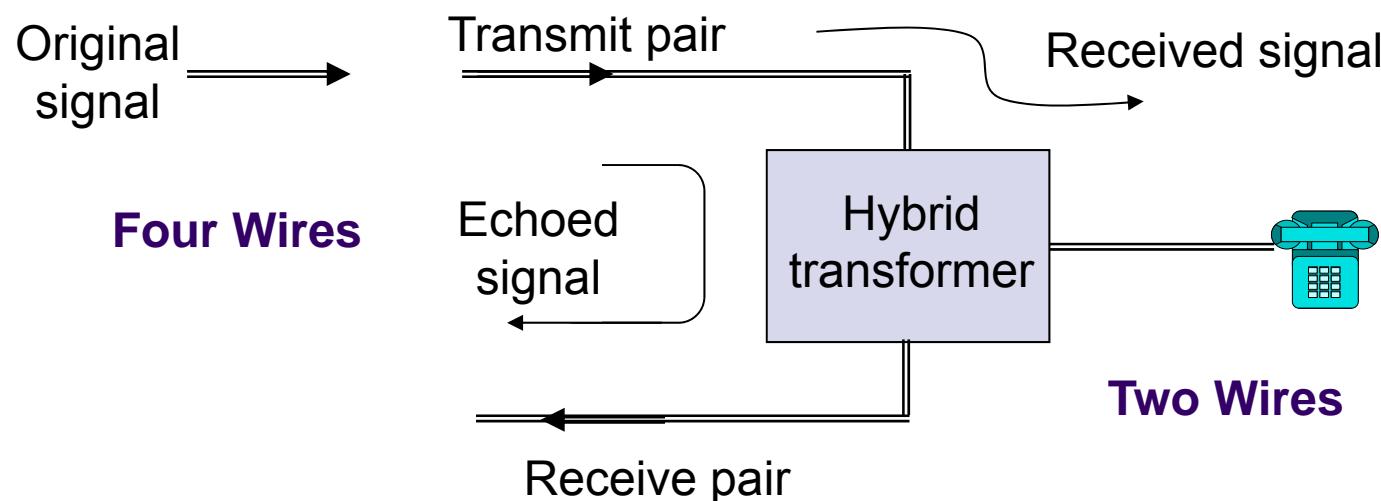
Standard	Data Rate	Distance
T-1	1.544 Mbps	18,000 feet, 5.5 km
DS2	6.312 Mbps	12,000 feet, 3.7 km
1/4 STS-1	12.960 Mbps	4500 feet, 1.4 km
1/2 STS-1	25.920 Mbps	3000 feet, 0.9 km
STS-1	51.840 Mbps	1000 feet, 300 m

- Fiber connection to the home provides huge amount of bandwidth, but cost of optical modems still high
- Fiber to the curve (pedestal) with shorter distance from pedestal to home can provide high speeds over copper pairs

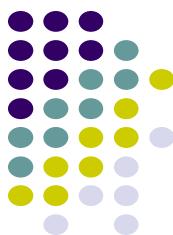
# Two- & Four-wire connections



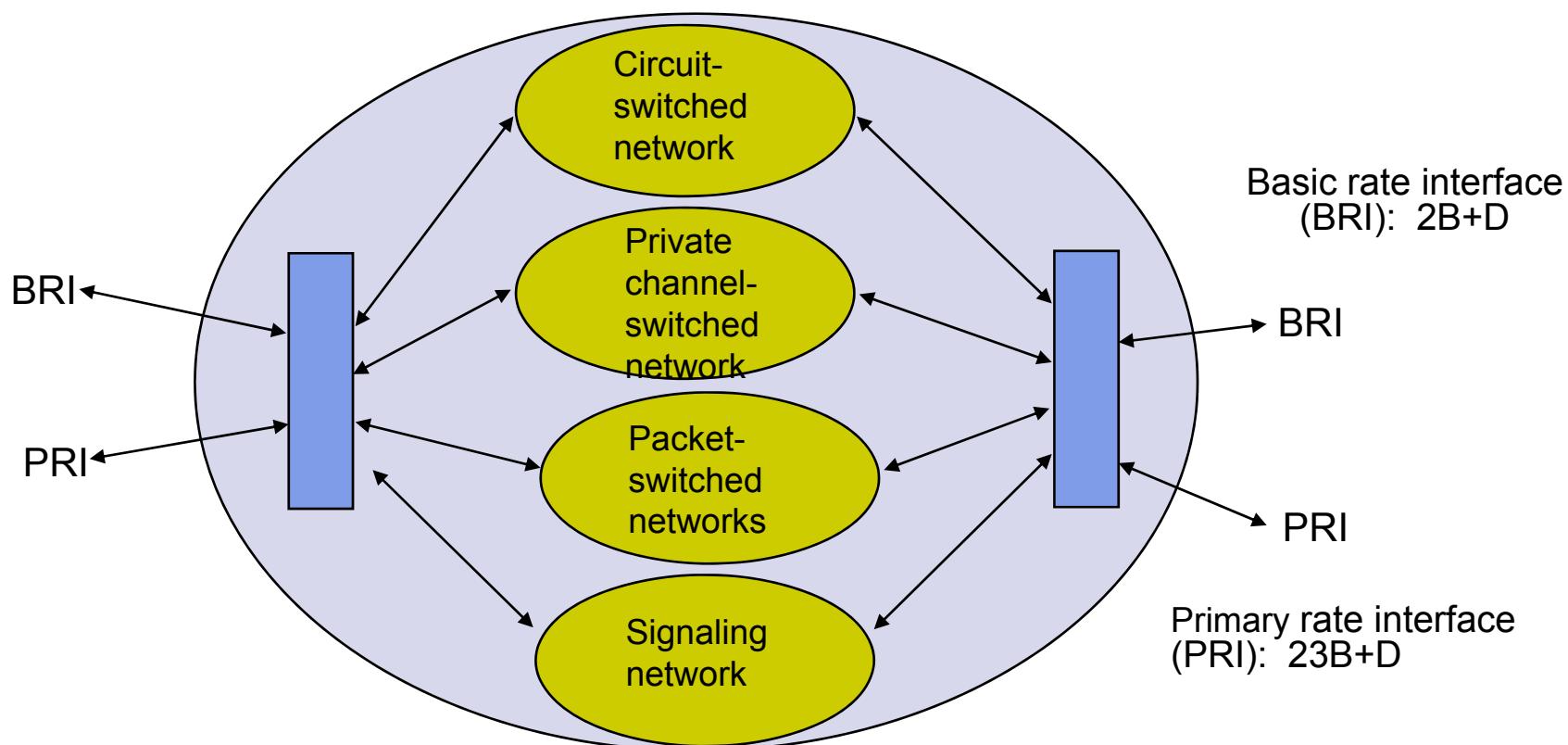
- From telephone to CO, two wires carry signals in both directions
- Inside network, 1 wire pair per direction
- Conversion from 2-wire to 4-wire occurs at hybrid transformer in the CO
- Signal reflections can occur causing speech echo
- Echo cancellers used to subtract the echo from the voice signals



# Integrated Services Digital Network (ISDN)

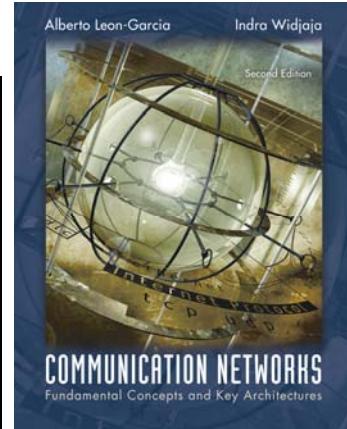


- First effort to provide end-to-end digital connections
- B channel = 64 kbps, D channel = 16 kbps
- ISDN defined interface to network
- Network consisted of separate networks for voice, data, signaling

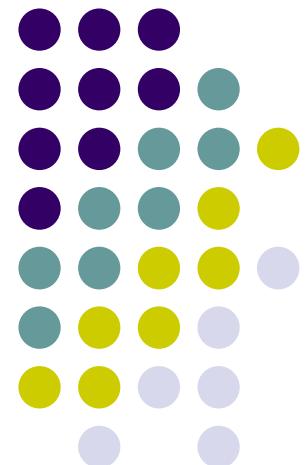


# Chapter 4

# Circuit-Switching Networks



*Signaling*



# Setting Up Connections



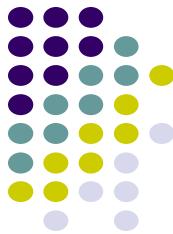
## Manually

- Human Intervention
- Telephone
  - Voice commands & switchboard operators
- Transport Networks
  - Order forms & dispatching of craftpersons

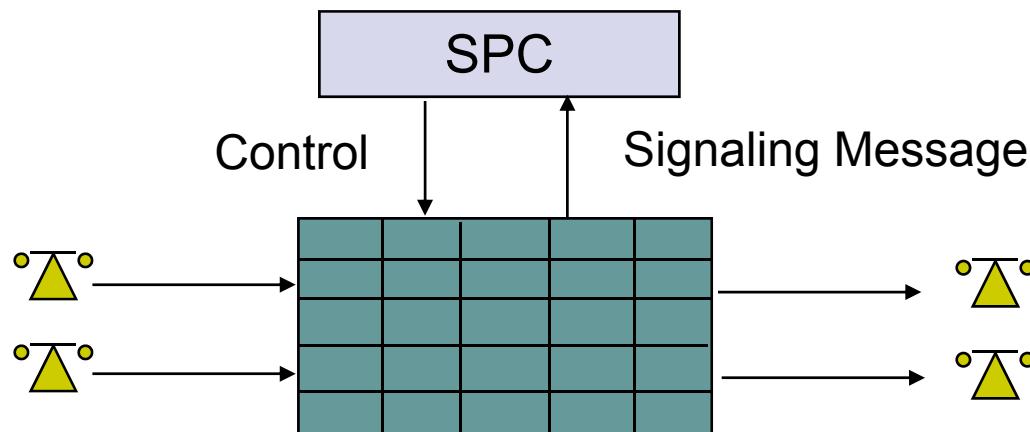
## Automatically

- Management Interface
  - Operator at console sets up connections at various switches
- Automatic signaling
  - Request for connection generates signaling messages that control connection setup in switches

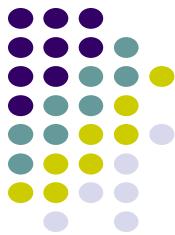
# Stored-Program Control Switches



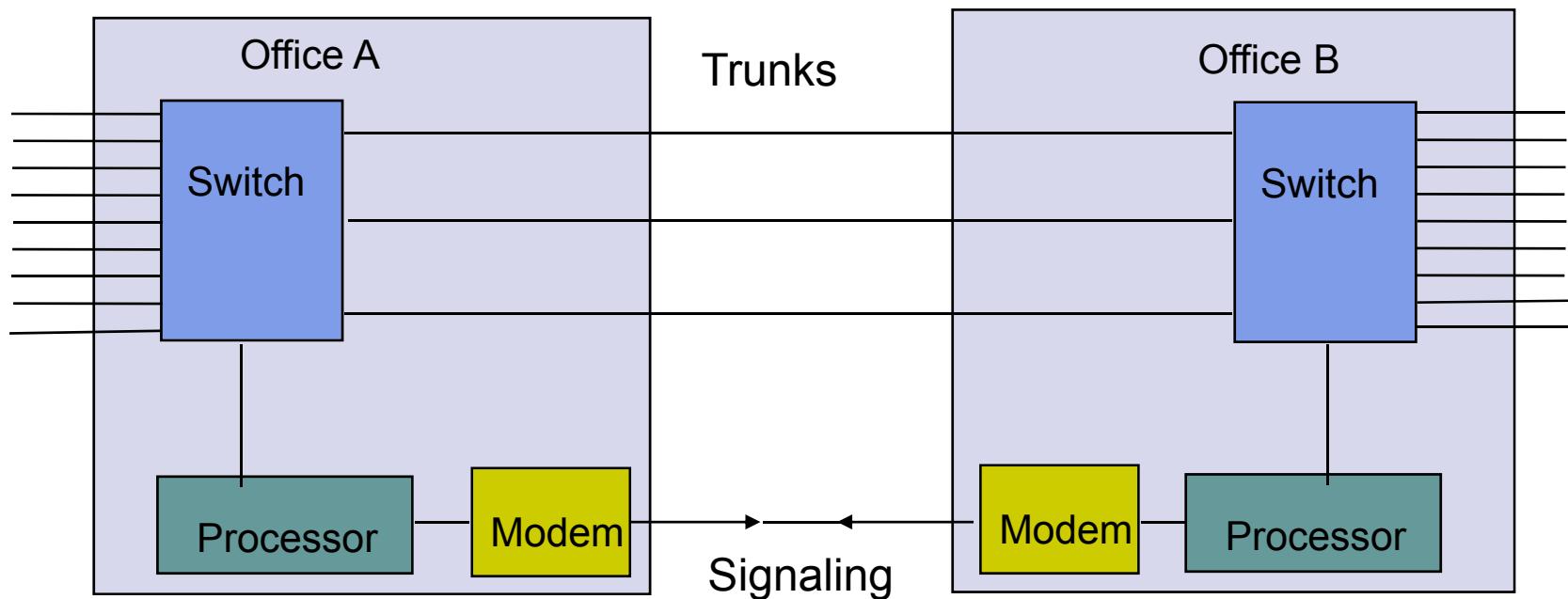
- SPC switches (1960s)
  - Crossbar switches with crossbars built from relays that open/close mechanically through electrical control
  - Computer program controls set up opening/closing of crosspoints to establish connections between switch inputs and outputs
- Signaling required to coordinate path set up across network



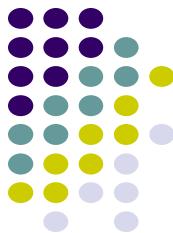
# Message Signaling



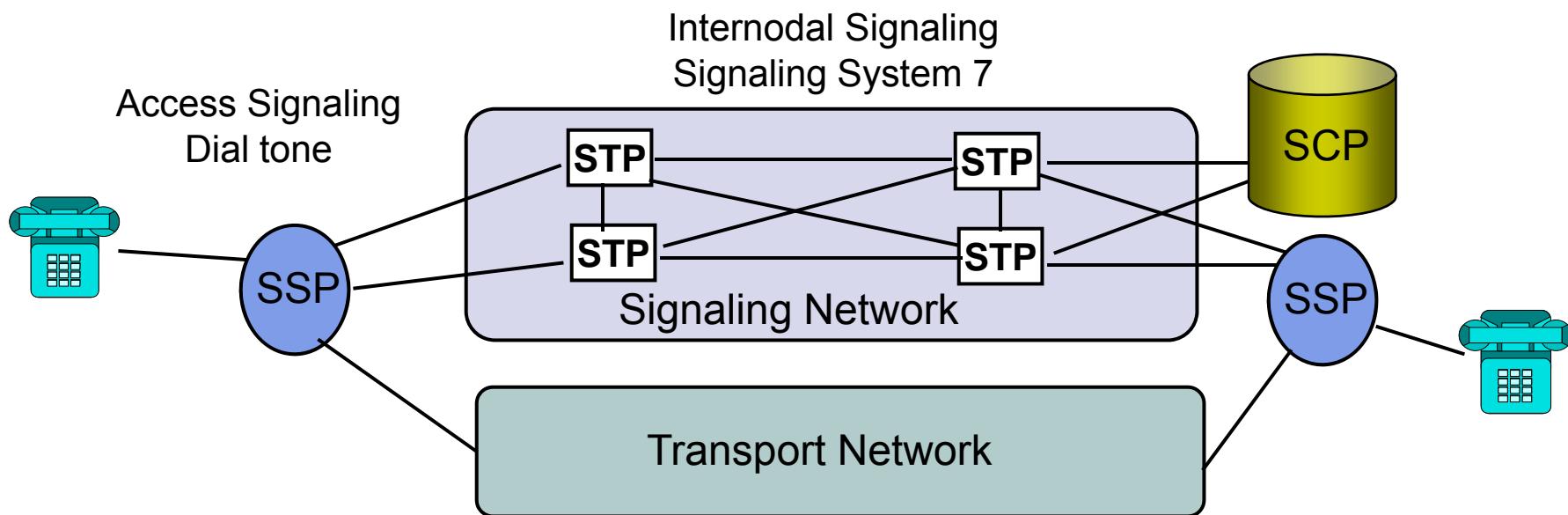
- Processors that control switches exchange signaling messages
- Protocols defining messages & actions defined
- Modems developed to communicate digitally over converted voice trunks



# Signaling Network

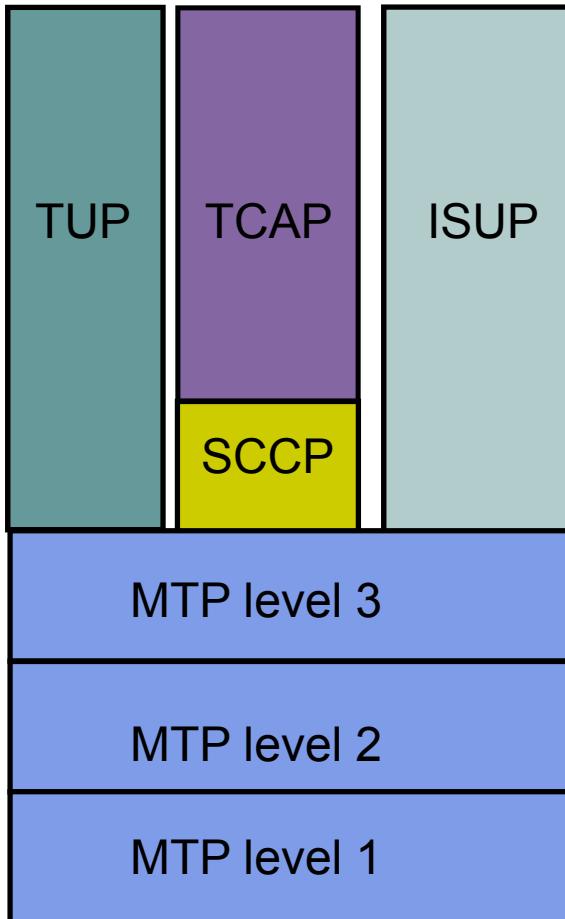
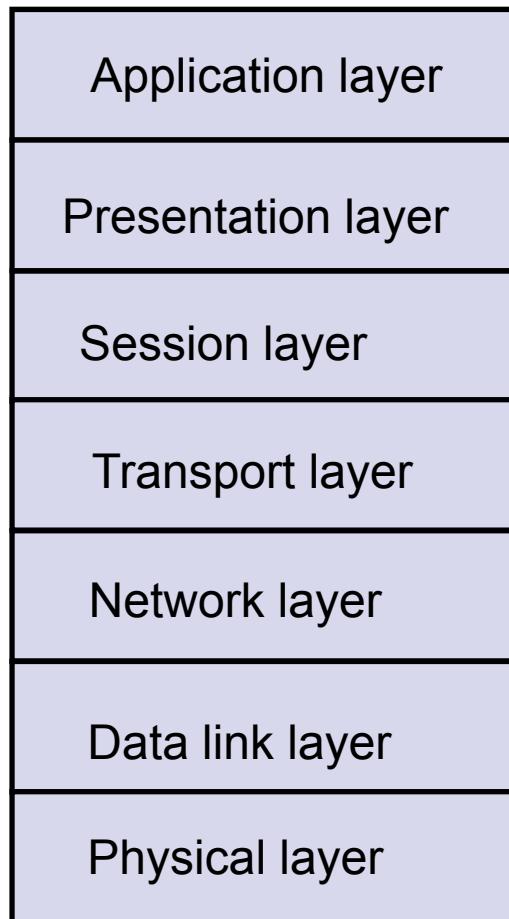
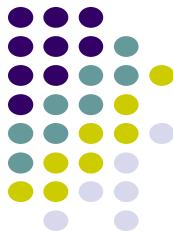


- Common Channel Signaling (CCS) #7 deployed in 1970s to control call setup
- Protocol stack developed to support signaling
- Signaling network based on highly reliable packet switching network
- Processors & databases attached to signaling network enabled many new services: caller id, call forwarding, call waiting, user mobility



SSP = service switching point (signal to message)  
STP = signal transfer point (packet switch)  
SCP = service control point (processing)

# Signaling System Protocol Stack



- Lower 3 layers ensure delivery of messages to signaling nodes
- SCCP allows messages to be directed to applications
- TCAP defines messages & protocols between applications
- ISUP performs basic call setup & release
- TUP instead of ISUP in some countries

ISUP = ISDN user part

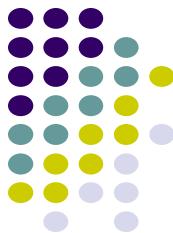
SSCP = signaling connection control part

TUP = telephone user part

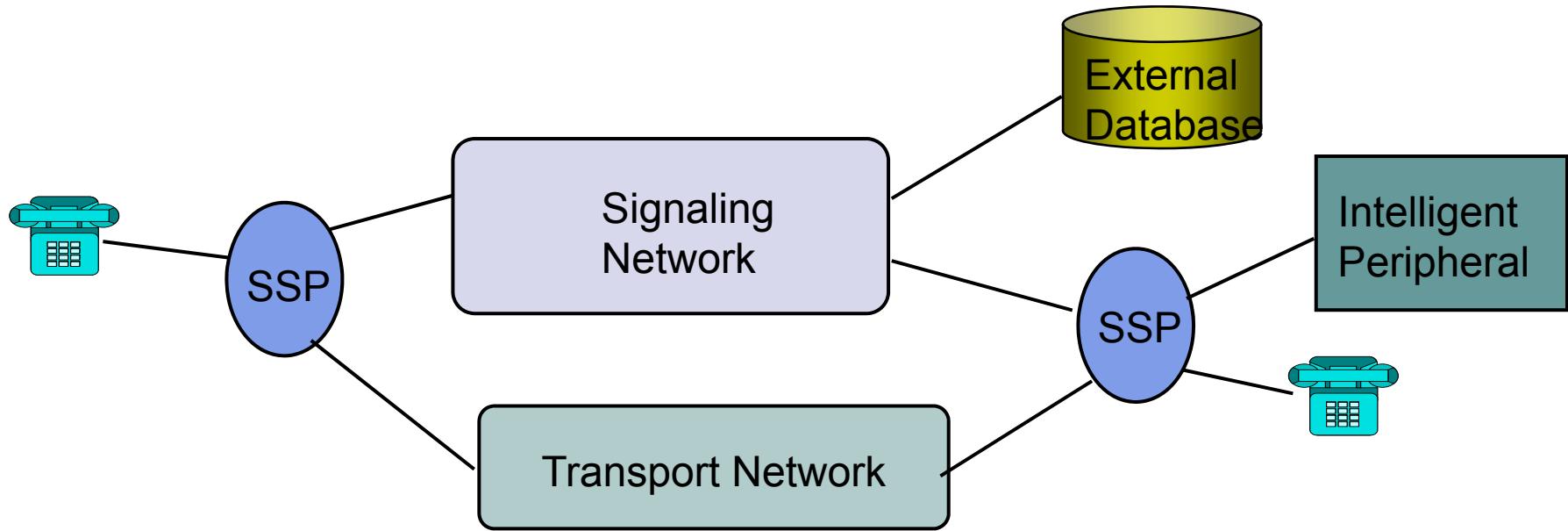
MTP = message transfer part

TCAP = transaction capabilities part

# Network Intelligence

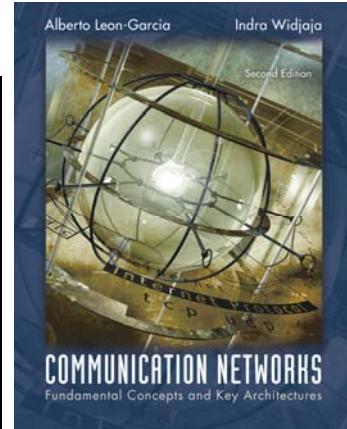


- Intelligent Peripherals provide additional service capabilities
- Voice Recognition & Voice Synthesis systems allow users to access applications via speech commands
- “Voice browsers” currently under development (See: [www.voicexml.org](http://www.voicexml.org))
- Long-term trend is for IP network to replace signaling system and provide equivalent services
- Services can then be provided by telephone companies as well as new types of service companies

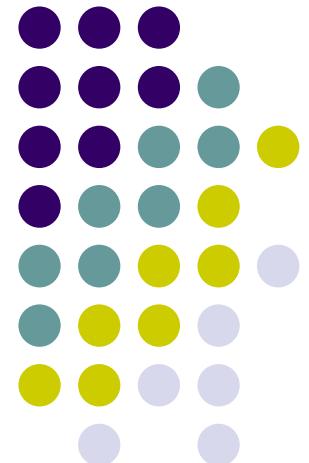


# Chapter 4

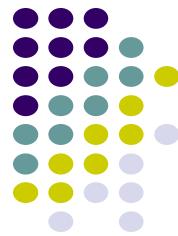
# Circuit-Switching Networks



*Traffic and Overload Control in  
Telephone Networks*

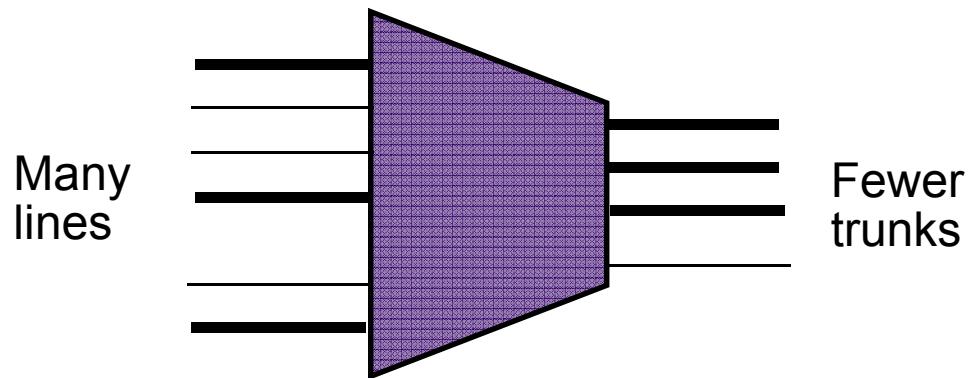


# Traffic Management & Overload Control

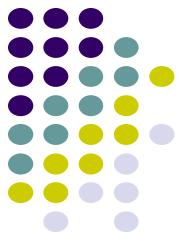


- Telephone calls come and go
- People activity follow patterns
  - Mid-morning & mid-afternoon at office
  - Evening at home
  - Summer vacation
- Outlier Days are extra busy
  - Mother's Day, Christmas, ...
- Disasters & other events cause surges in traffic
- Need traffic management & overload control

# Traffic concentration

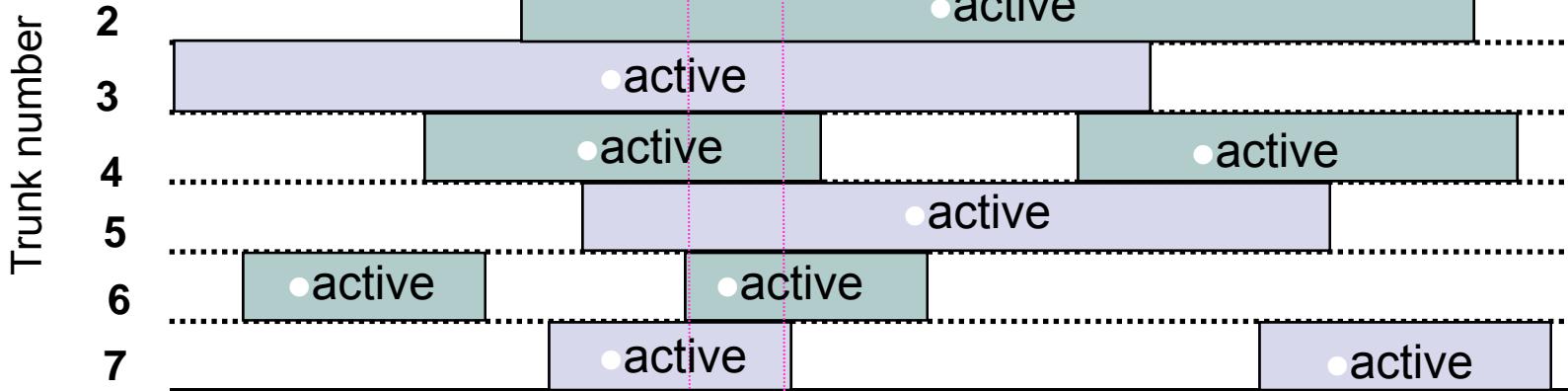
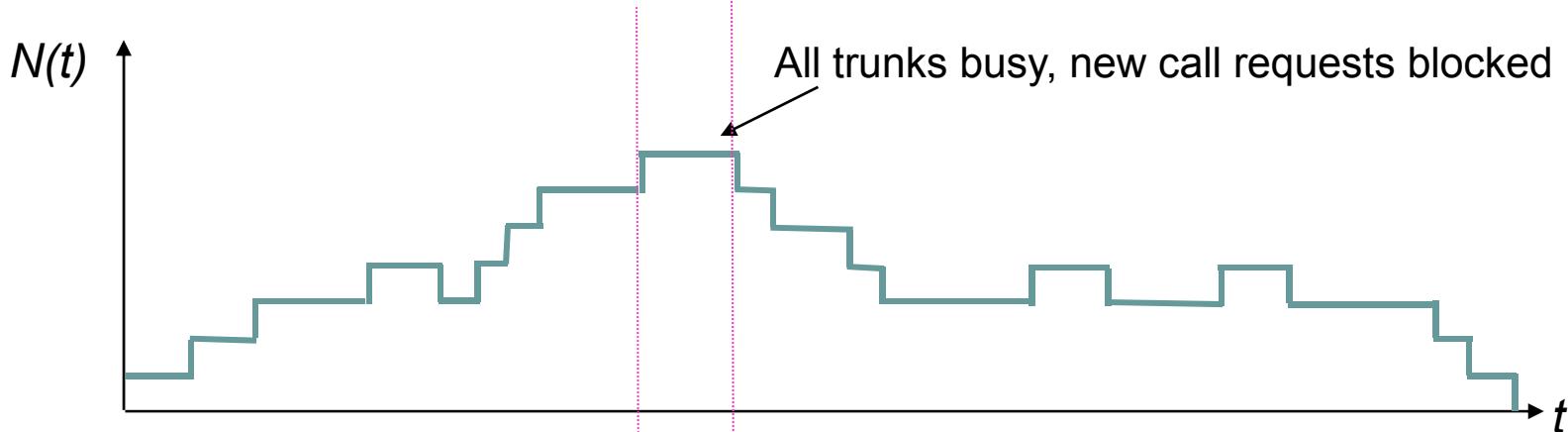


- Traffic fluctuates as calls initiated & terminated
  - Driven by human activity
- Providing resources so
  - Call requests *always* met is too expensive
  - Call requests met *most of the time* cost-effective
- Switches concentrate traffic onto *shared* trunks
  - Blocking of requests will occur from time to time
- Traffic engineering provisions resources to meet blocking performance targets



# Fluctuation in Trunk Occupancy

Number of busy trunks





# Modeling Traffic Processes

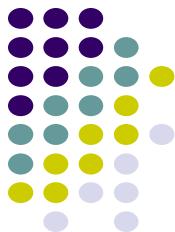
- Find the statistics of  $N(t)$  the number of calls in the system

## Model

- Call request ***arrival rate***:  $\lambda$  requests per second
- In a very small time interval  $\Delta$ ,
  - Prob[ new request ] =  $\lambda\Delta$
  - Prob[no new request] =  $1 - \lambda\Delta$
- The resulting random process is a Poisson arrival process:

$$\text{Prob}(k \text{ arrivals in time } T) = \frac{(\lambda T)^k e^{-\lambda T}}{k!}$$

- Holding time***: Time a user maintains a connection
  - $X$  a random variable with mean  $E(X)$
- Offered load***: rate at which work is offered by users:
  - $a = \lambda \text{ calls/sec} * E(X) \text{ seconds/call}$  (Erlangs)



# Blocking Probability & Utilization

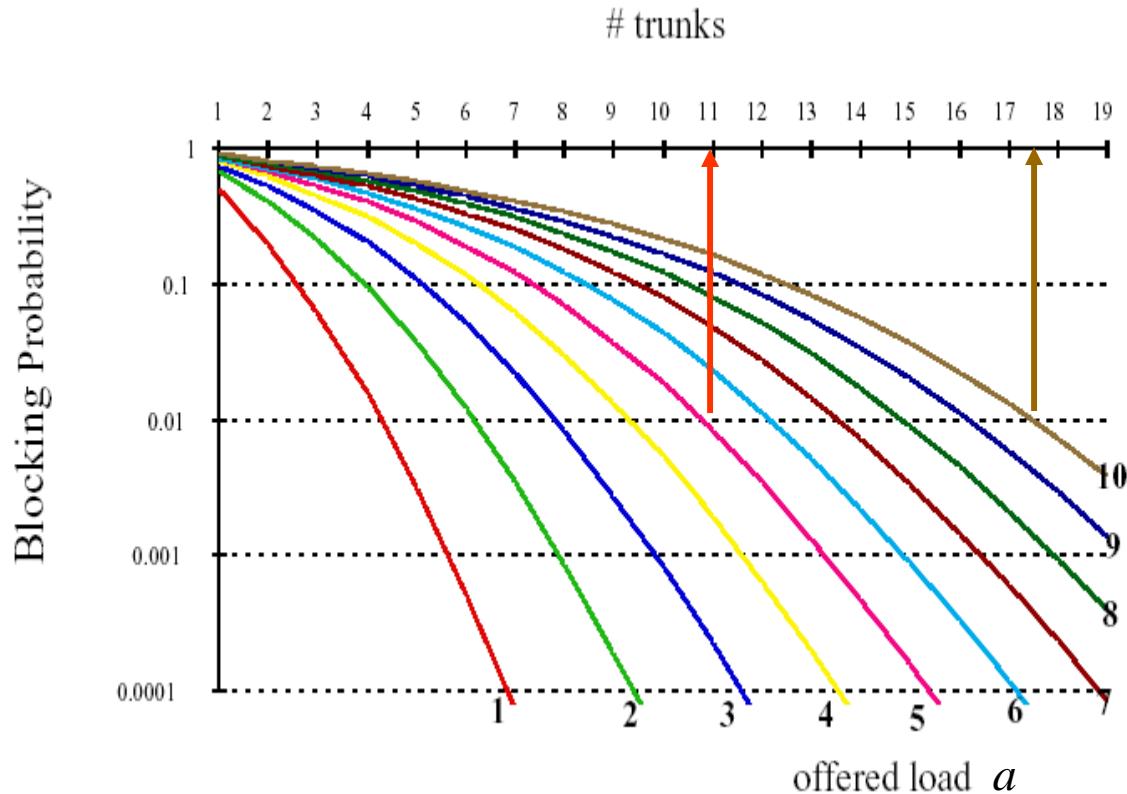
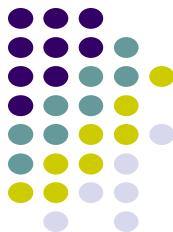
- $c$  = Number of Trunks
- Blocking occurs if all trunks are busy, i.e.  $N(t)=c$
- If call requests are Poisson, then blocking probability  $P_b$  is given by **Erlang B Formula**

$$P_b = \frac{a^c / c!}{\sum_{k=0}^c a^k / k!}$$

- The *utilization* is the average # of trunks in use

$$\text{Utilization} = \lambda(1 - P_b) E[X]/c = (1 - P_b) a/c$$

# Blocking Performance



To achieve 1% blocking probability:

$a = 5$  Erlangs requires 11 trunks

$a = 10$  Erlangs requires 18 trunks

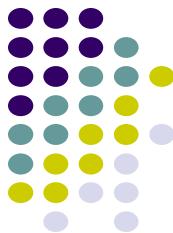


# Multiplexing Gain

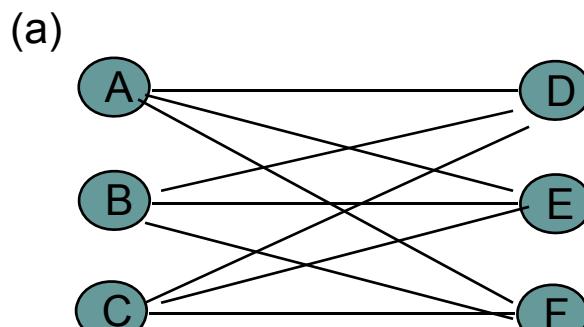
Load	Trunks@1%	Utilization
1	5	0.20
2	7	0.29
3	8	0.38
4	10	0.40
5	11	0.45
6	13	0.46
7	14	0.50
8	15	0.53
9	17	0.53
10	18	0.56
30	42	0.71
50	64	0.78
60	75	0.80
90	106	0.85
100	117	0.85

- At a given  $P_b$ , the system becomes more efficient in utilizing trunks with increasing system size
- Aggregating traffic flows to share centrally allocated resources is more efficient
- This effect is called *Multiplexing Gain*

# Routing Control

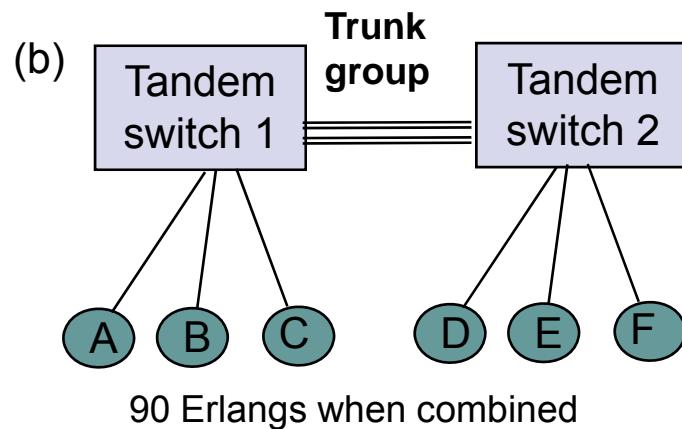


- Routing control: selection of connection paths
- Large traffic flows should follow direct route because they are efficient in use of resources
- Useful to combine smaller flows to share resources
- Example: 3 close CO's & 3 other close COs
- 10 Erlangs between each pair of COs



10 Erlangs between each pair

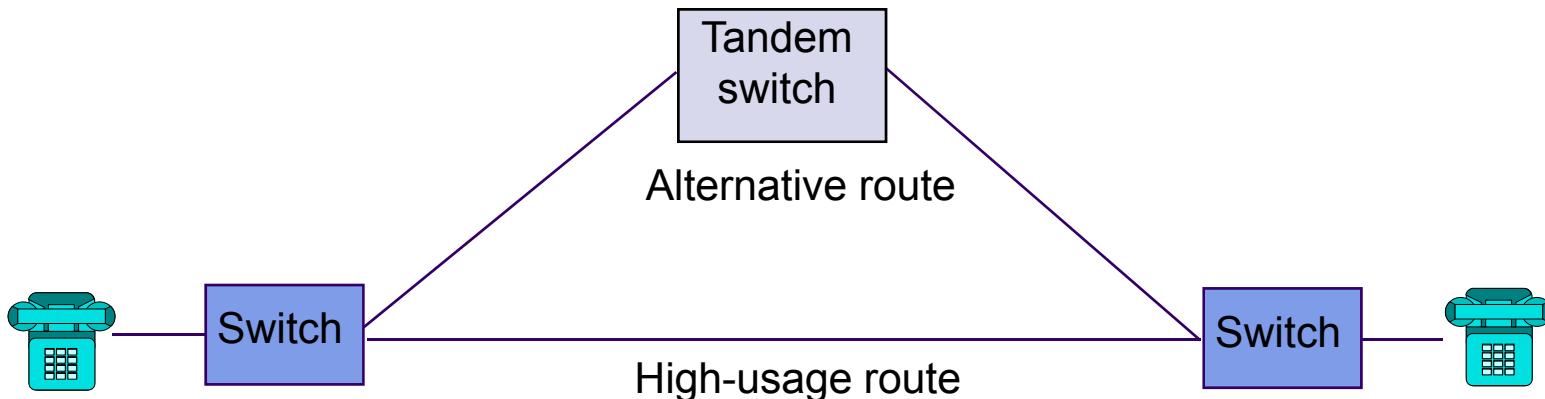
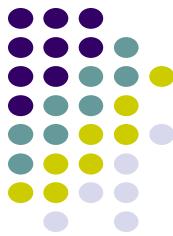
17 trunks for 10 Erlangs  
 $9 \times 17 = 153$  trunks  
Efficiency =  $90/153 = 53\%$



90 Erlangs when combined

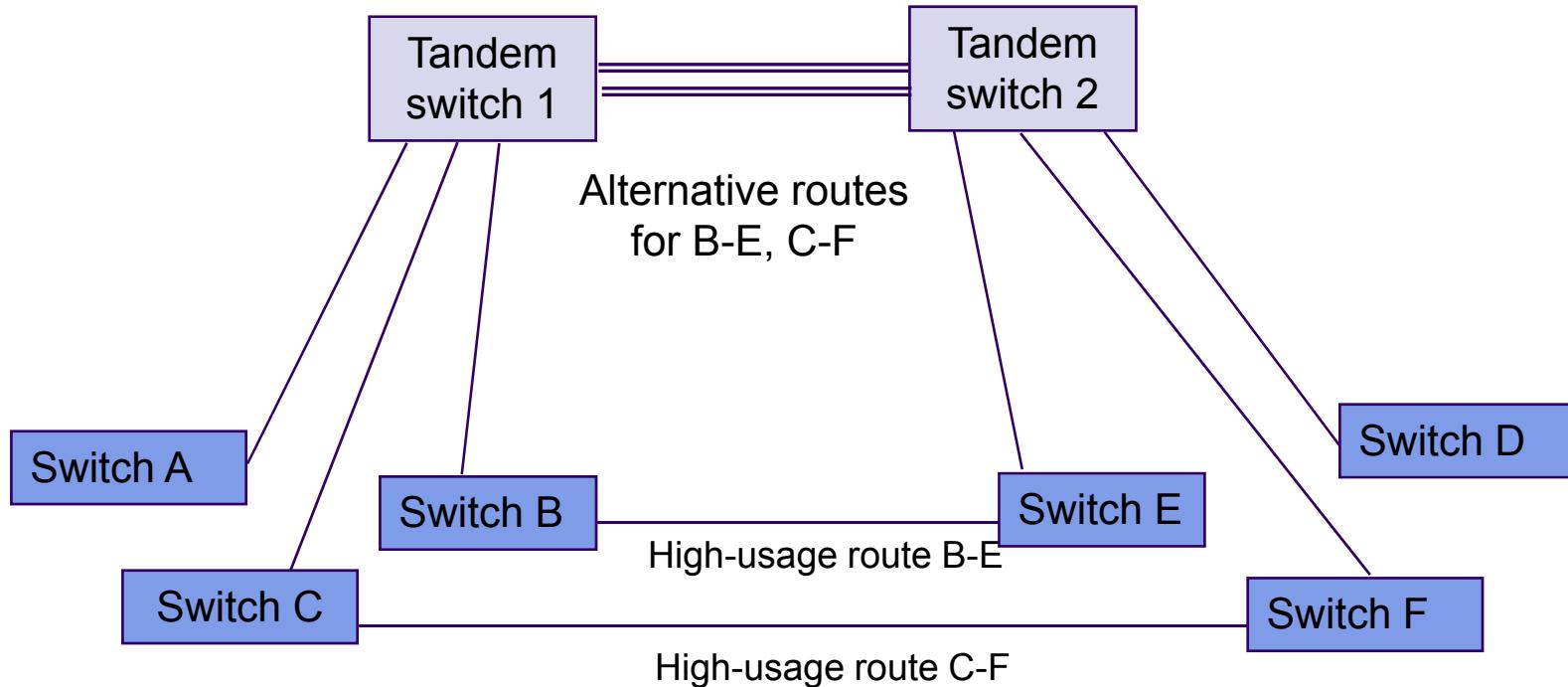
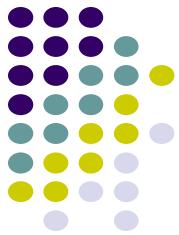
106 trunks for 90 Erlangs  
Efficiency = 85%

# Alternative Routing

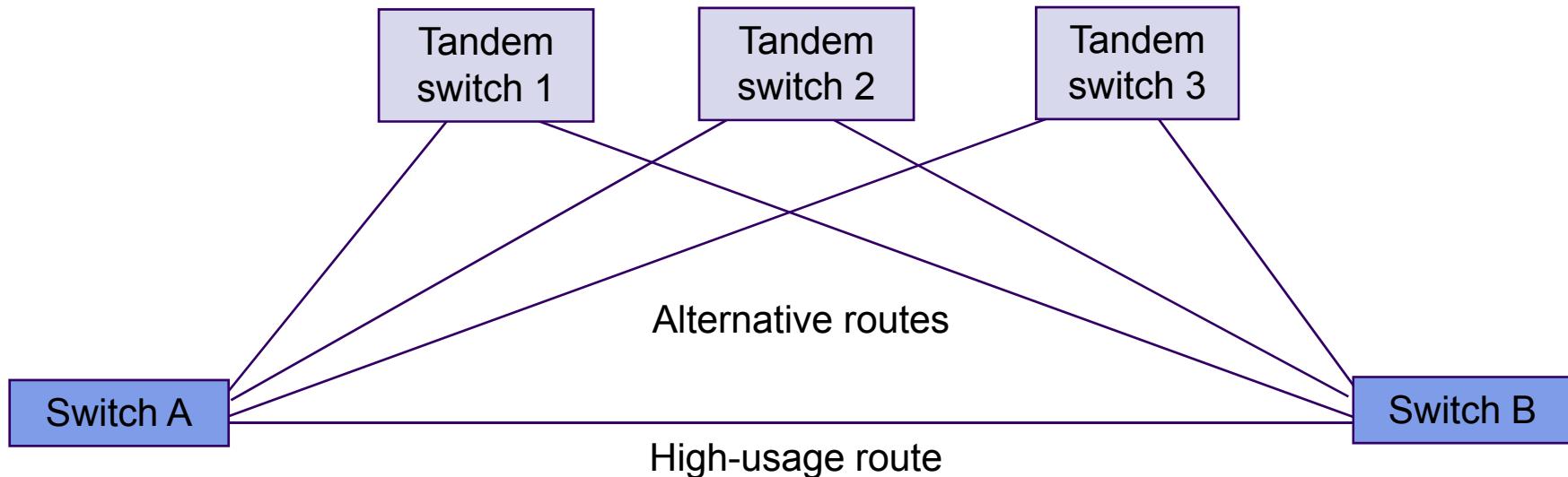
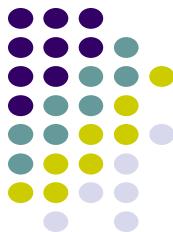


- Deploy trunks between switches with significant traffic volume
- Allocate trunks with high blocking, say 10%, so utilization is high
- Meet 1% end-to-end blocking requirement by overflowing to longer paths over tandem switch
- Tandem switch handles overflow traffic from other switches so it can operate efficiently
- Typical scenario shown in next slide

# Typical Routing Scenario

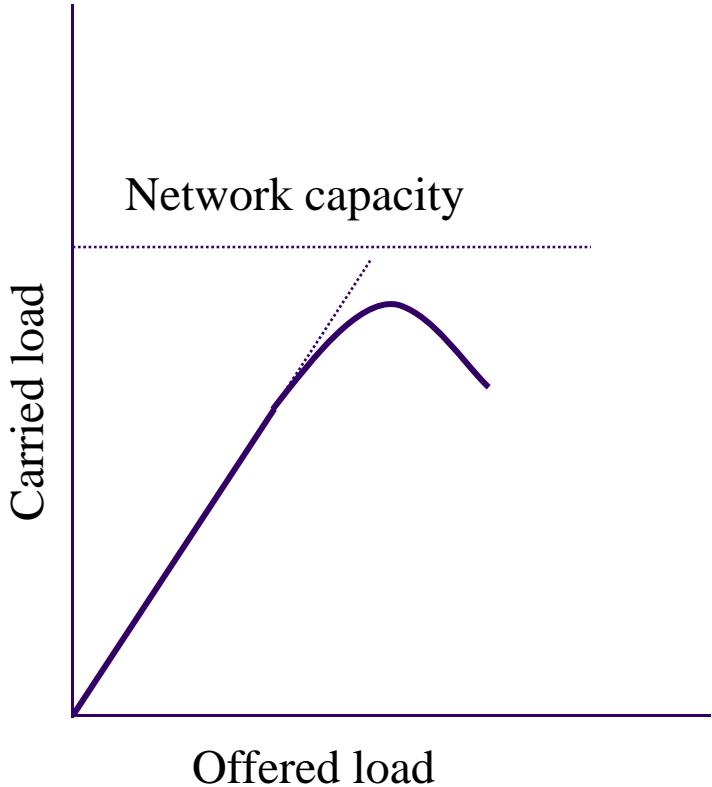
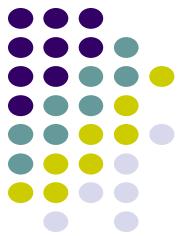


# Dynamic Routing



- Traffic varies according to time of day, day of week
  - East coast of North America busy while West coast idle
- Network can use idle resources by adapting route selection dynamically
  - Route some intra-East-coast calls through West-coast switches
- Try high-usage route and overflow to alternative routes

# Overload Control



## Overload Situations

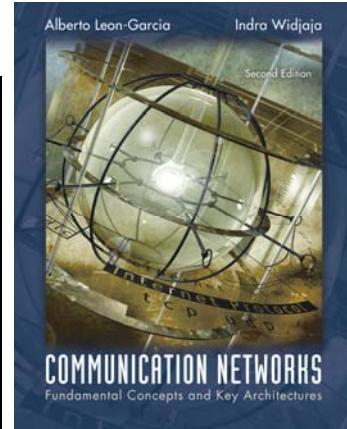
- Mother's Day, Xmas
- Catastrophes
- Network Faults

## Strategies

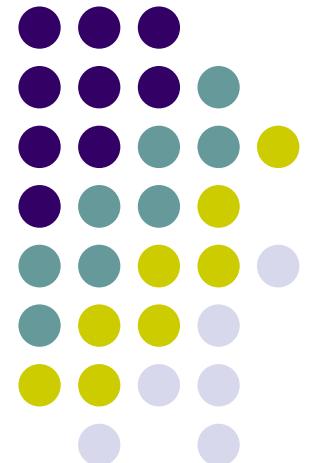
- Direct routes first
- Outbound first
- Code blocking
- Call request pacing

# Chapter 4

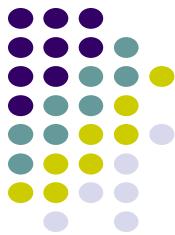
# Circuit-Switching Networks



## *Cellular Telephone Networks*

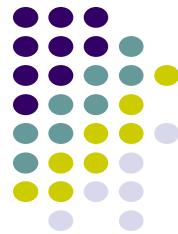


# Radio Communications



- 1900s: Radio telephony demonstrated
- 1920s: Commercial radio broadcast service
- 1930s: Spectrum regulation introduced to deal with interference
- 1940s: Mobile Telephone Service
  - Police & ambulance radio service
  - Single antenna covers transmission to mobile users in city
  - Less powerful car antennas transmit to network of antennas around a city
  - Very limited number of users can be supported

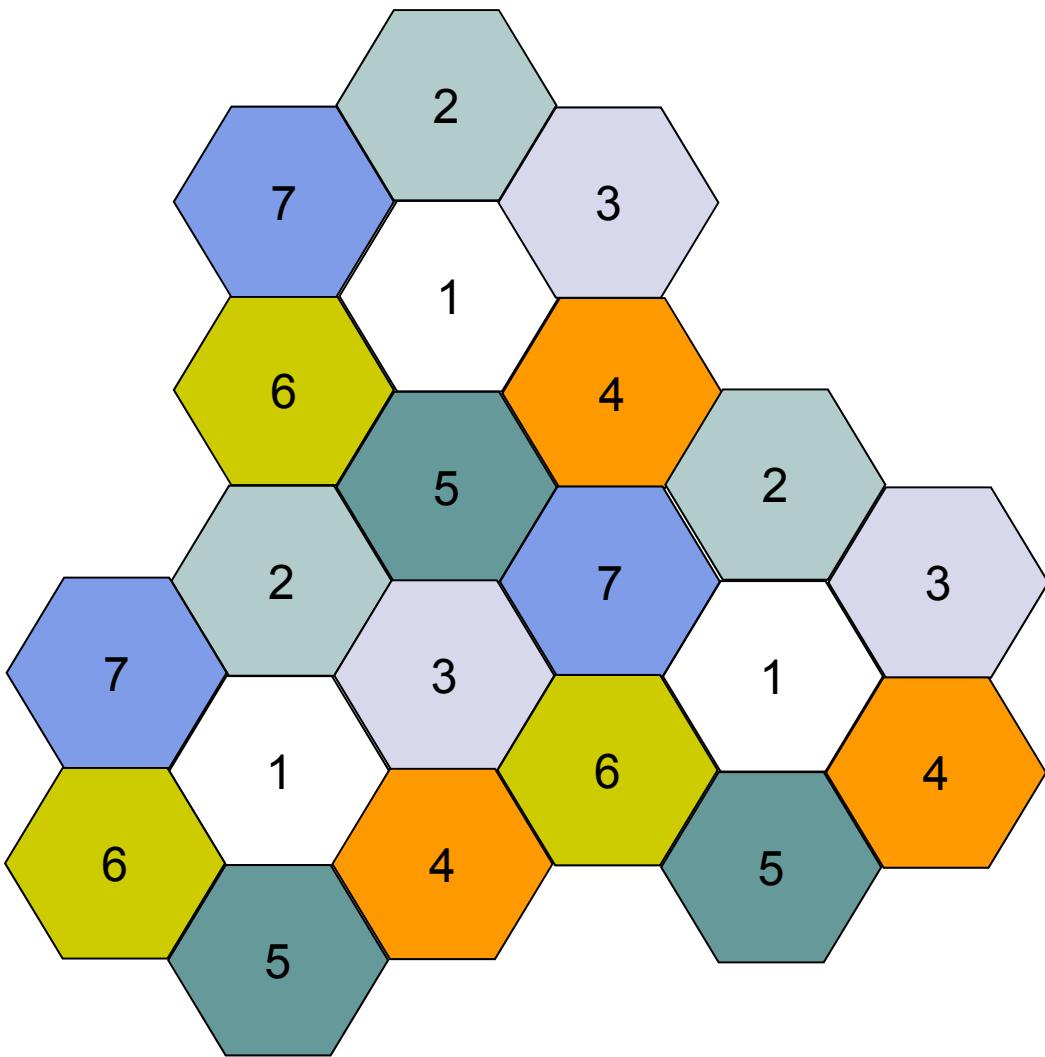
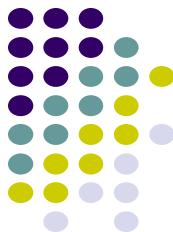
# Cellular Communications



Two basic concepts:

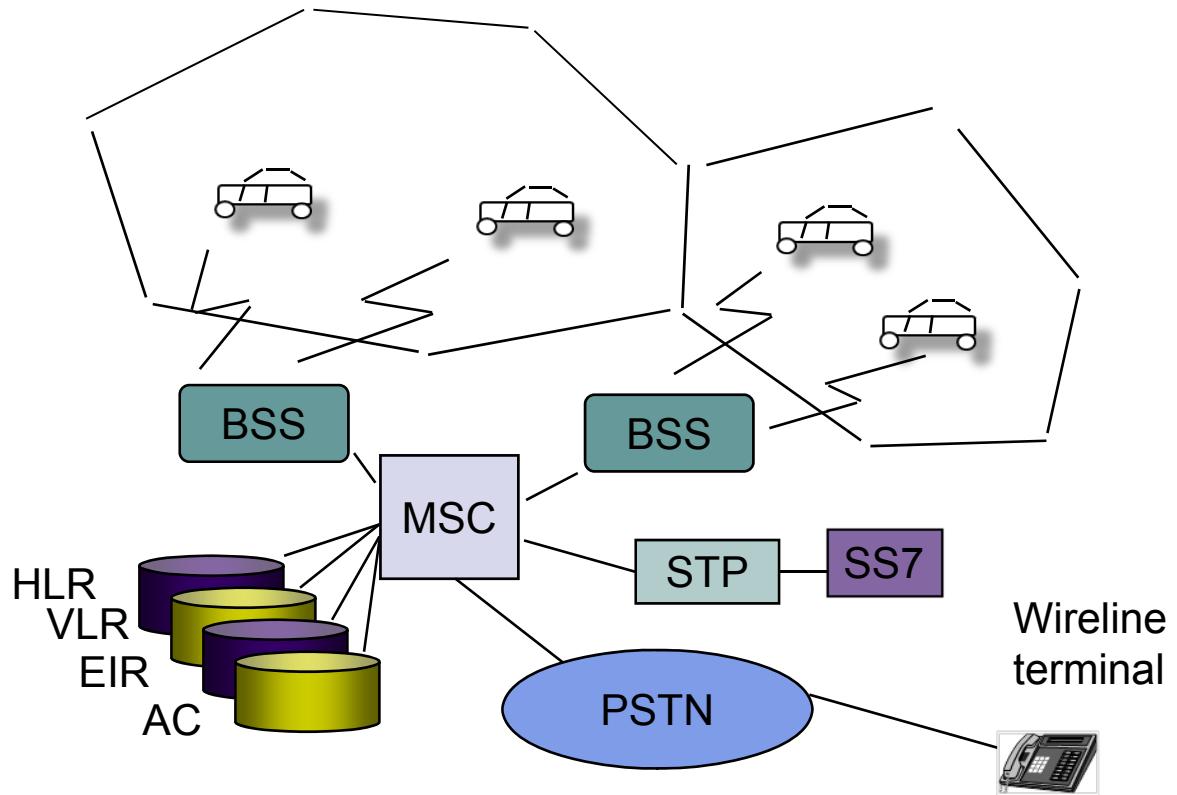
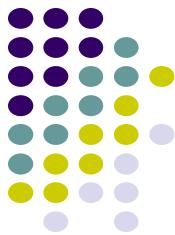
- Frequency Reuse
  - A region is partitioned into *cells*
  - Each cell is covered by *base station*
  - Power transmission levels controlled to minimize inter-cell interference
  - Spectrum can be reused in other cells
- Handoff
  - Procedures to ensure continuity of call as user moves from cell to another
  - Involves setting up call in new cell and tearing down old one

# Frequency Reuse



- Adjacent cells may not use same band of frequencies
- Frequency Reuse Pattern specifies how frequencies are reused
- Figure shows 7-cell reuse: frequencies divided into 7 groups & reused as shown
- Also 4-cell & 12-cell reuse possible
- Note: CDMA allows adjacent cells to use same frequencies (Chapter 6)

# Cellular Network



AC = authentication center

BSS = base station subsystem

EIR = equipment identity register

HLR = home location register

MSC = mobile switching center

PSTN = public switched telephone network

STP = signal transfer point

VLR = visitor location register

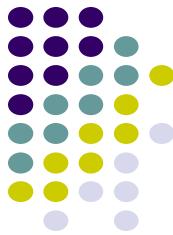
## Base station

- Transmits to users on *forward channels*
- Receives from users on *reverse channels*

## Mobile Switching Center

- Controls connection setup within cells & to telephone network

# Signaling & Connection Control

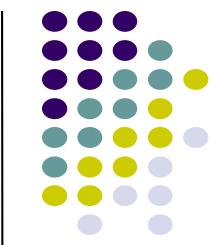


- *Setup channels* set aside for call setup & handoff
  - Mobile unit selects setup channel with strongest signal & monitors this channel
- **Incoming call to mobile unit**
  - MSC sends call request to all BSSs
  - BSSs broadcast request on all setup channels
  - Mobile unit replies on reverse setup channel
  - BSS forwards reply to MSC
  - BSS assigns forward & reverse voice channels
  - BSS informs mobile to use these
  - Mobile phone rings



# Mobile Originated Call

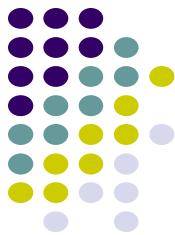
- Mobile sends request in reverse setup channel
- Message from mobile includes serial # and possibly authentication information
- BSS forwards message to MSC
- MSC consults Home Location Register for information about the subscriber
- MSC may consult Authentication center
- MSC establishes call to PSTN
- BSS assigns forward & reverse channel



# Handoff

- Base station monitors signal levels from its mobiles
- If signal level drops below threshold, MSC notified & mobile instructed to transmit on setup channel
- Base stations in vicinity of mobile instructed to monitor signal from mobile on setup channel
- Results forward to MSC, which selects new cell
- Current BSS & mobile instructed to prepare for handoff
- MSC releases connection to first BSS and sets up connection to new BSS
- Mobile changes to new channels in new cell
- Brief interruption in connection (except for CDMA)

# Roaming



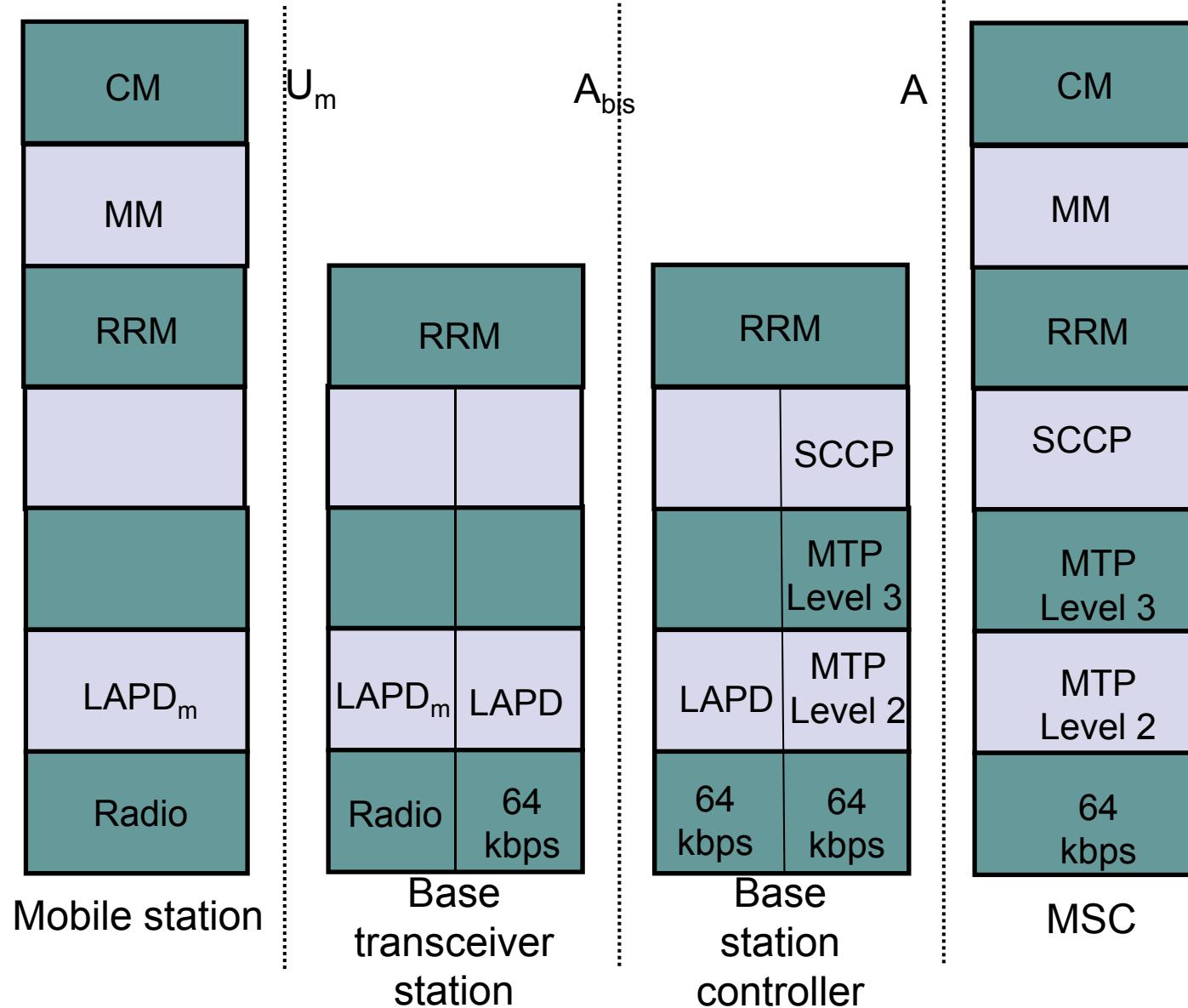
- Users subscribe to roaming service to use service outside their home region
- Signaling network used for message exchange between home & visited network
- Roamer uses setup channels to register in new area
- MSC in visited areas requests authorization from user's Home Location Register
- Visitor Location Register informed of new user
- User can now receive & place calls

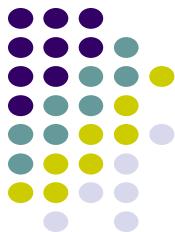
# GSM Signaling Standard



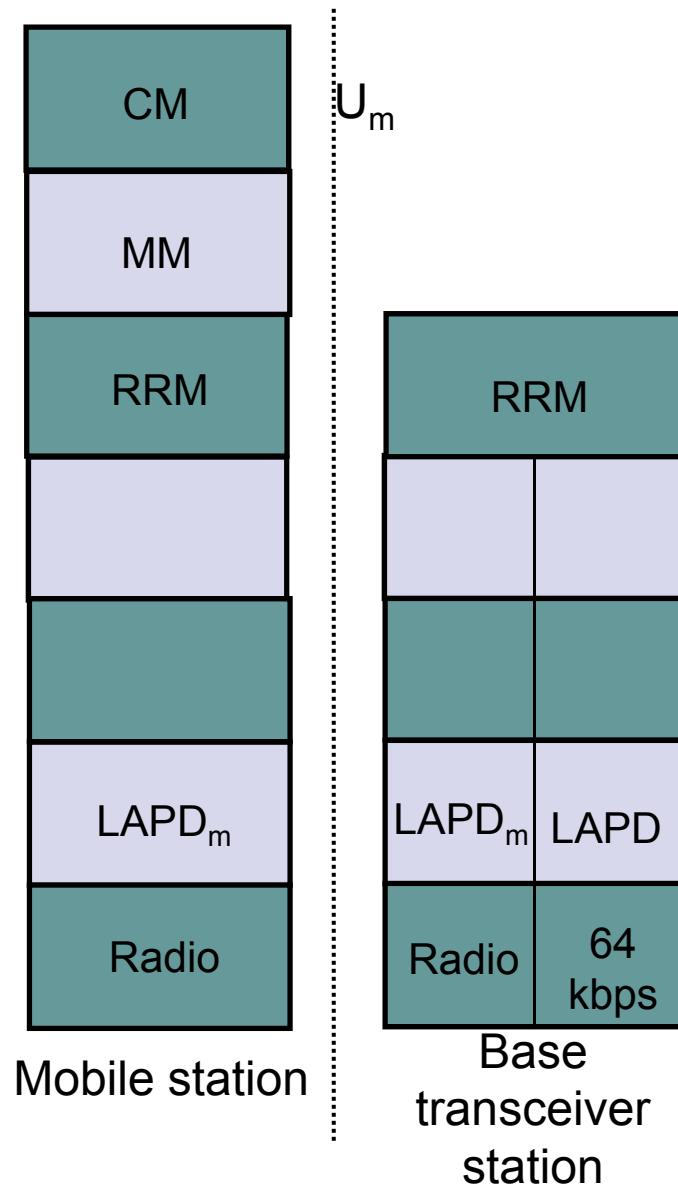
- Base station
  - Base Transceiver Station (BTS)
    - Antenna + Transceiver to mobile
    - Monitoring signal strength
  - Base Station Controller
    - Manages radio resources or 1 or more BTSs
    - Set up of channels & handoff
    - Interposed between BTS & MSC
- Mobile & MSC Applications
  - Call Management (CM)
  - Mobility Management (MM)
- Radio Resources Management (RRM) concerns mobile, BTS, BSC, and MSC

# Cellular Network Protocol Stack





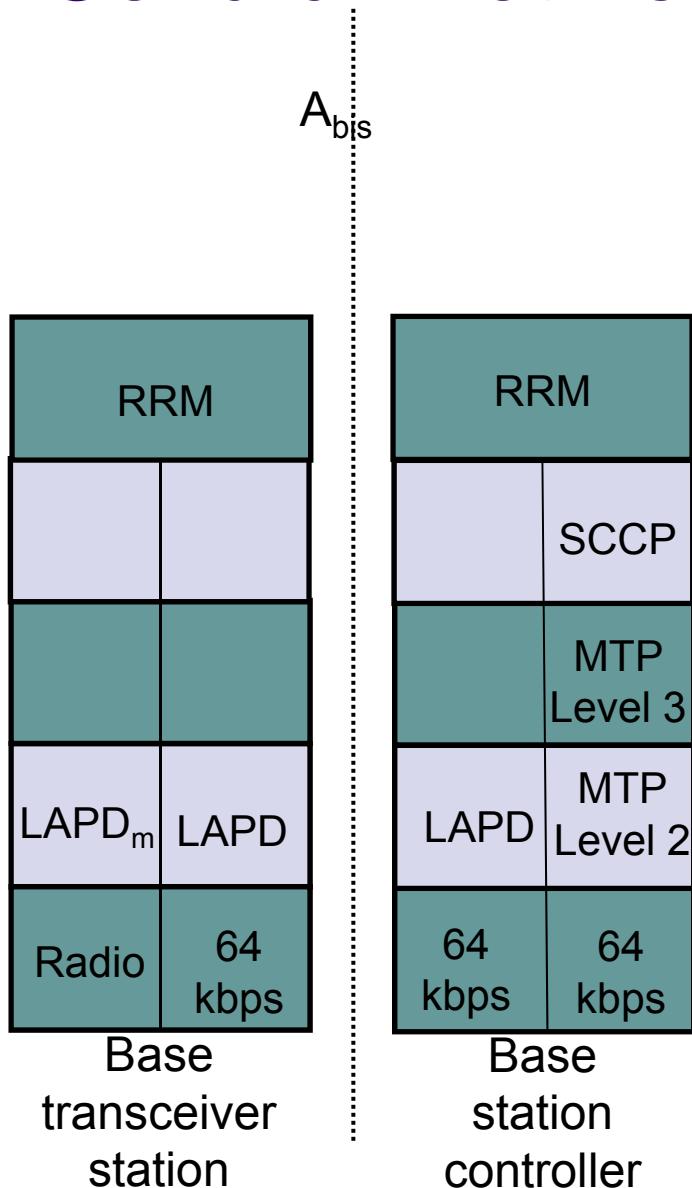
# Cellular Network Protocol Stack



## Radio Air Interface ( $U_m$ )

- LAPD<sub>m</sub> is data link control adapted to mobile
- RRM deals with setting up of radio channels & handover

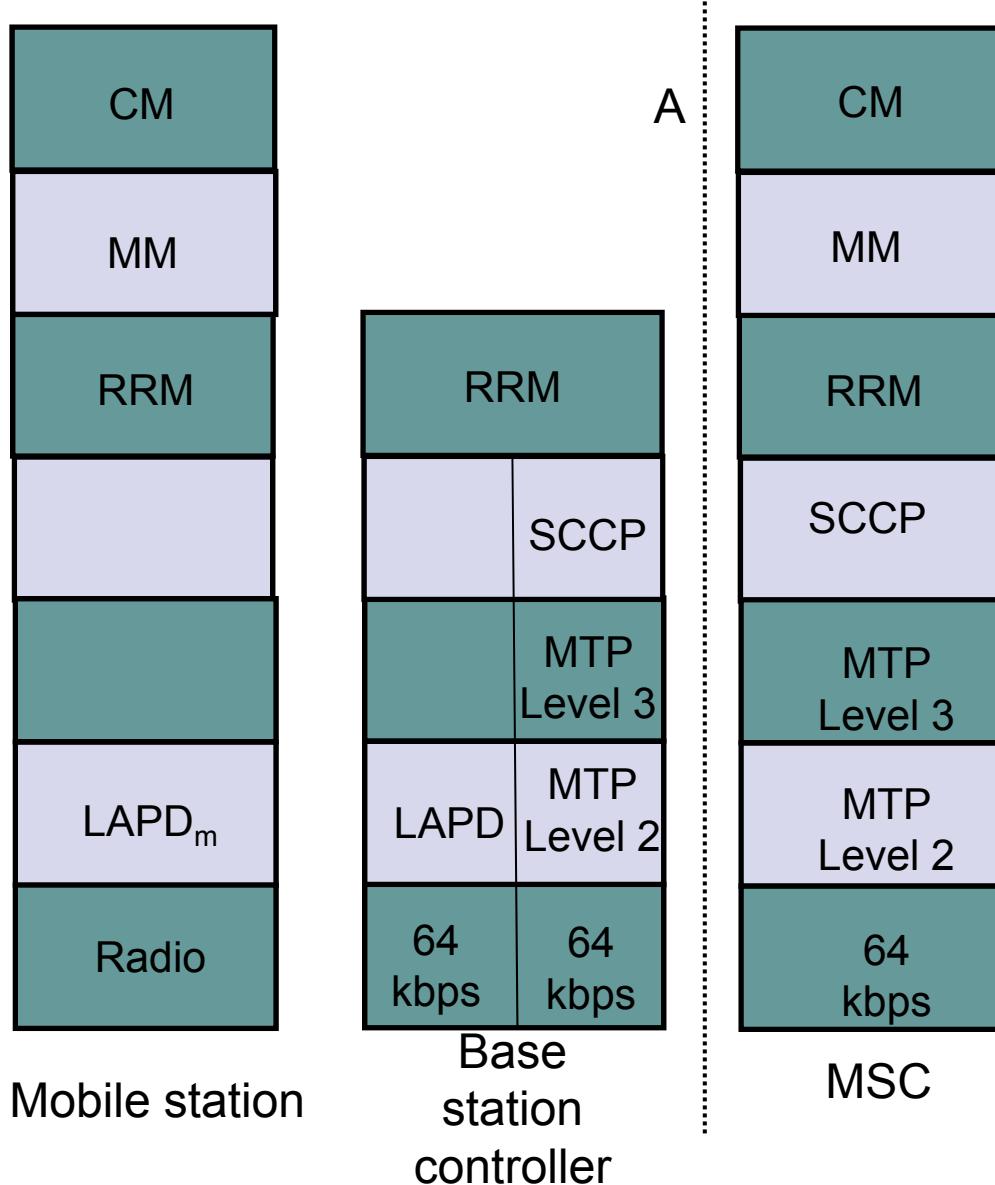
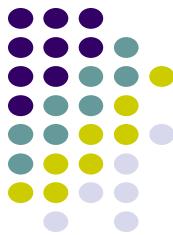
# Cellular Network Protocol Stack



## $A_{bis}$ Interface

- 64 kbps link physical layer
- $LAPD_m$
- BSC RRM can handle handover for cells within its control

# Cellular Network Protocol Stack



## Signaling Network (A) Interface

- RRM deals handover involving cells with different BSCs
- MM deals with mobile user location, authentication
- CM deals with call setup & release using modified ISUP