Measurement and Analysis of Traffic in a Hybrid Satellite-Terrestrial Network

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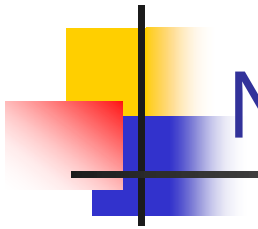
Road map

- Introduction and motivation
- Traffic:
  - collection
  - analysis
  - prediction
- Conclusions
- References



Network traffic measurements

- Focus of networking research during:
  - mid to late 1980's
  - early 1990's
- Motivation for traffic measurements:
  - understand traffic characteristics in deployed networks
  - develop traffic models
  - evaluate performance of protocols and applications
  - perform trace driven simulations

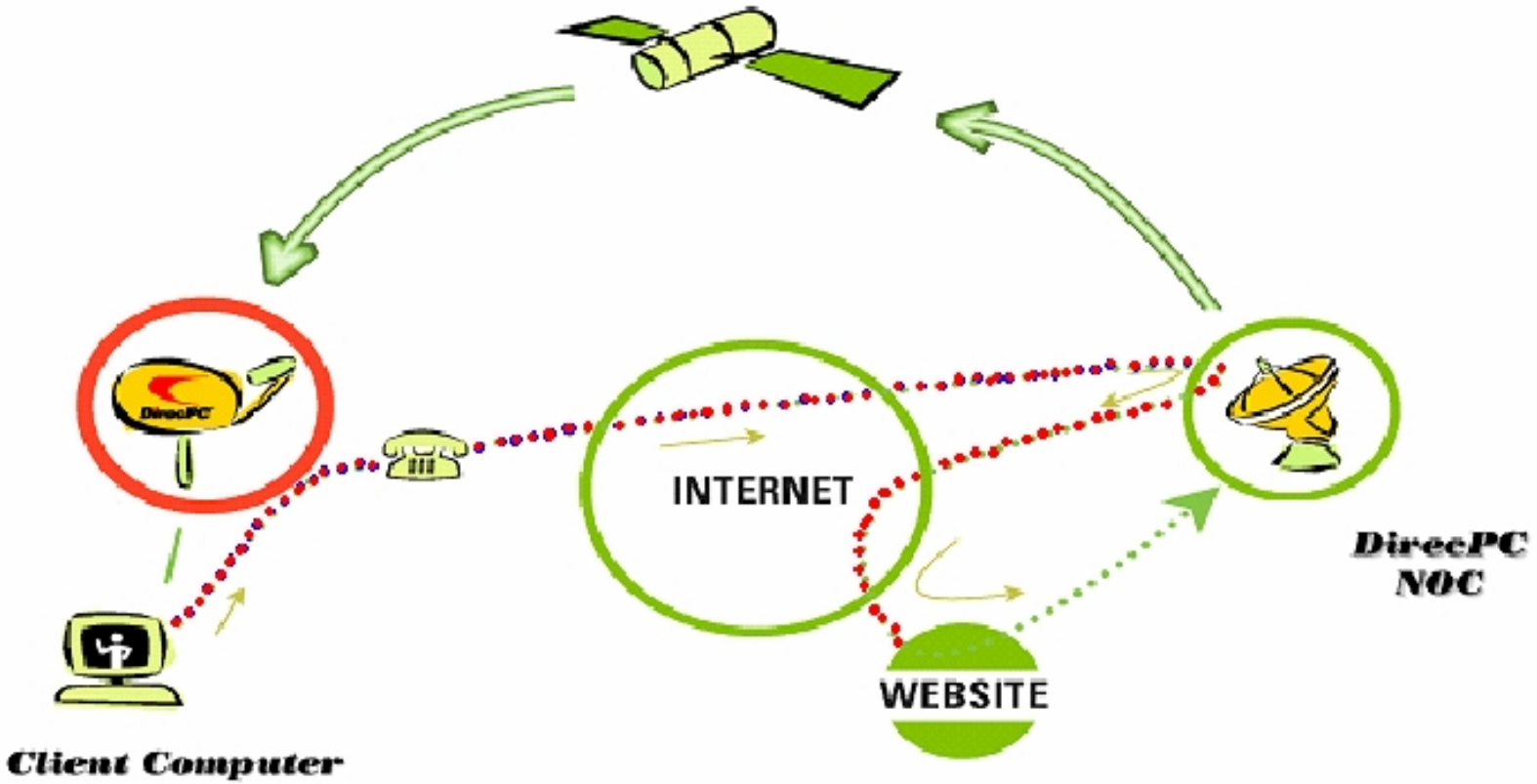
Traffic traces

- Most available traffic traces are from the wired networks within research communities:
  - Bellcore, LBNL, Auckland University
- Few traces were collected from wireless or satellite commercial networks
- Various factors affect Internet traffic patterns:
  - Web, Proxy, Napster, MP3, Web mail
- Used to evaluate the AutoRegressive Integrated Moving-Average (ARIMA) model for predicting uploaded and downloaded traffic

DirecPC system

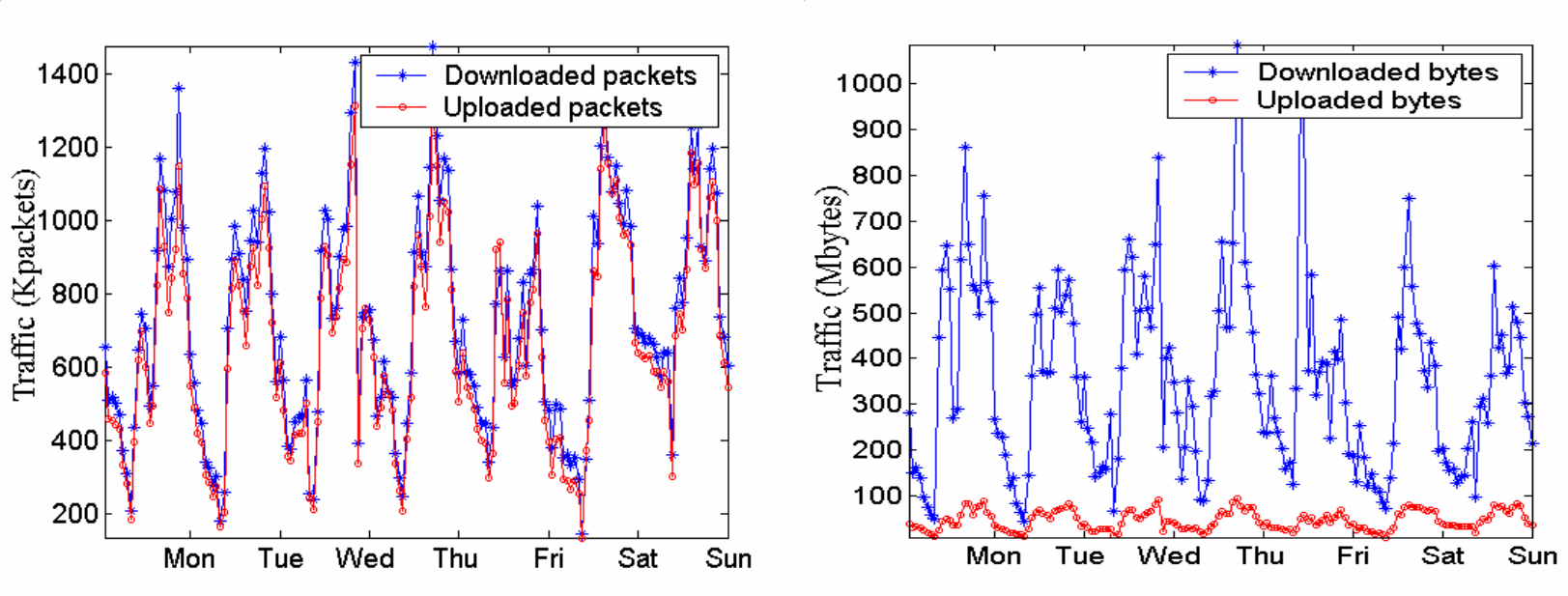
- Satellite one-way broadcast system manufactured by Hughes Network Systems
- DirecPC systems are deployed worldwide
- ChinaSat uses DirecPC system to provide Internet access to over 200 Internet cafés across provinces
- DirecPC utilizes two special techniques to improve network performance:
  - IP spoofing
  - TCP splitting

Traffic collection



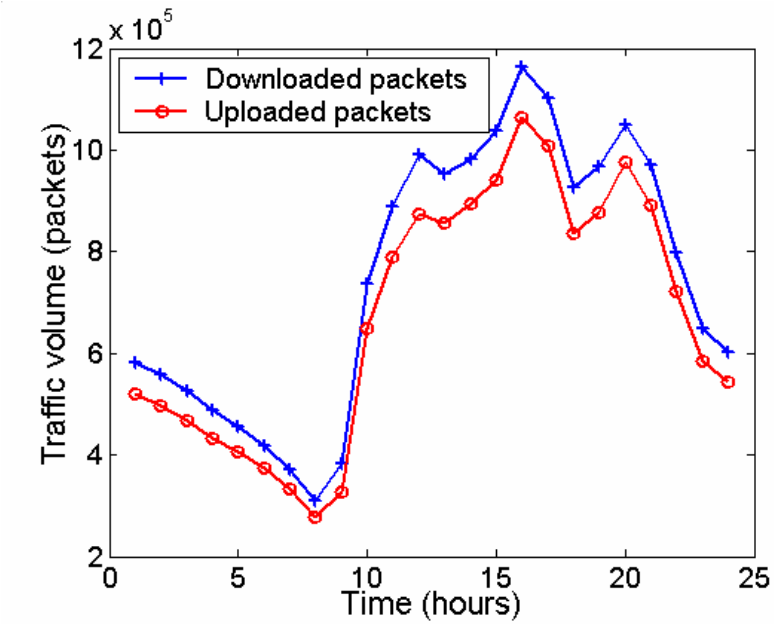
Red: uploaded traffic
Green: downloaded traffic

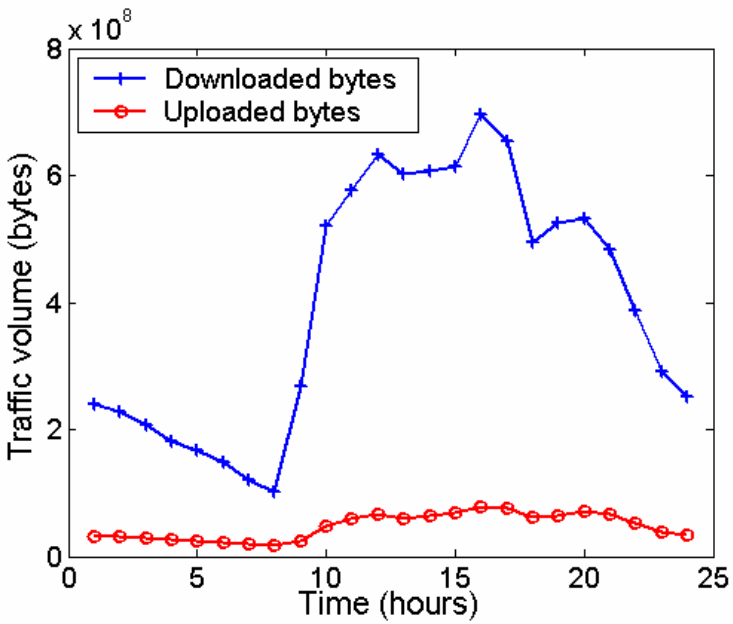
Analysis of weekly billing records



- Weekly traffic volume measured in packets (left) and bytes (right)
- Traffic data was collected from 09-12-2002 to 15-12-2002

Analysis of daily billing records





- Average traffic volume over a single day measured in packets (left) and bytes (right)
- Traffic data was collected from 9-12-2002 to 15-12-2002

Protocols and applications

Protocol	Packets	Packets (%)	Bytes	Bytes (%)
TCP	36,737,165	84.32	11,231,147,530	94.49
UDP	6,202,673	14.24	601,157,016	5.06
ICMP	630,528	1.45	53,128,377	0.45
Total	**43,570,366**	**~100**	**11,885,432,923**	**~100**

Applications	Connections	Connections (%)	Bytes	Bytes (%)
WWW	304,243	90.06	10,203,267,005	75.79
FTP-data	636	0.19	1,440,393,008	10.7
IRC	2,324	0.69	945,965	0.008
SMTP	562	0.17	2,326,373	0.01
POP-3	115	0.03	2,326,373	0.02
Telnet	70	0.02	280,286	0.002
Other	651	8.84	238,099,412	13.47
Total	**308,601**	**100**	**11,885,432,923**	**100**

- Traffic data was collected from 21-12-2002 22:08 to 23-12-2002 3:28

TCP connection level: Web traffic

- Zipf-like distribution: $fr \sim 1/r^{\beta}$

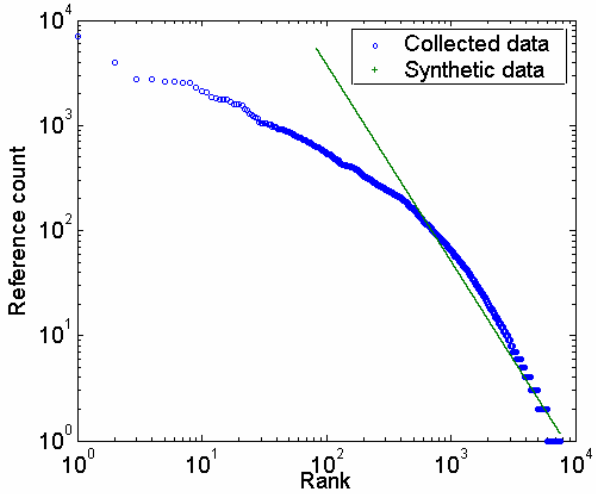
  the number of requests (frequency) is inversely proportional to its rank among the requests

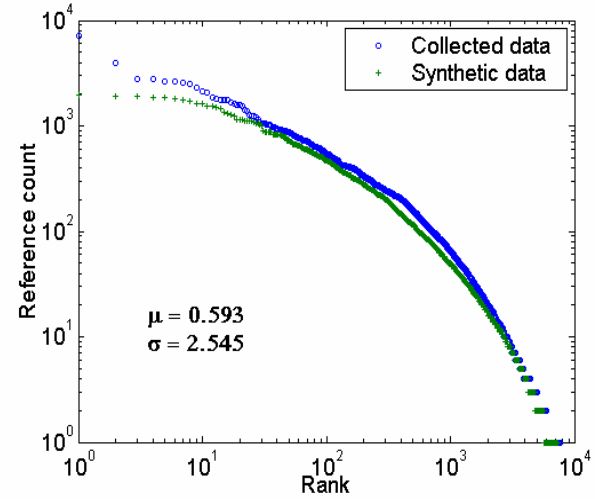
- DGX (discrete lognormal):

$$p(x=k)=\frac{A(\mu,\sigma)}{k}\exp[-\frac{(\ln k-\mu)^2}{2\sigma^2}]$$

$$A(\mu,\sigma)=\{\sum_{k=1}^{\infty}\frac{1}{k}[-\frac{(\ln k-\mu)^2}{2\sigma^2}]\}^{-1}$$

- DGX distribution fits better than the Zipf-like distribution

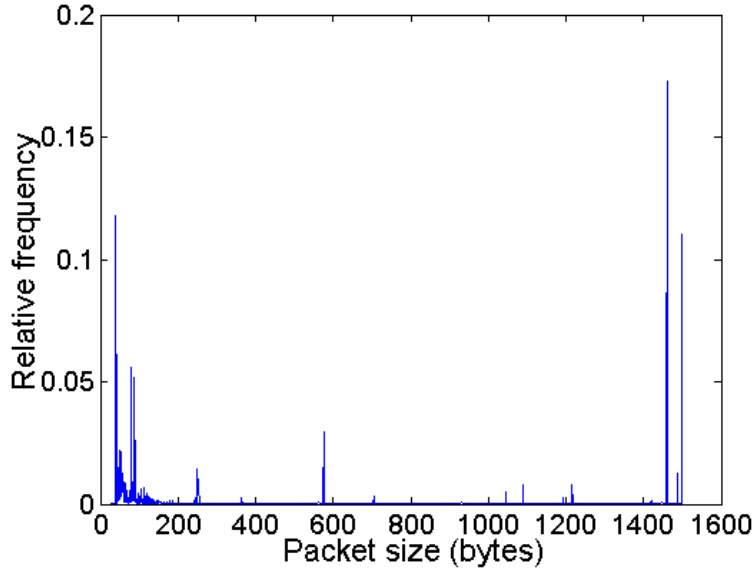


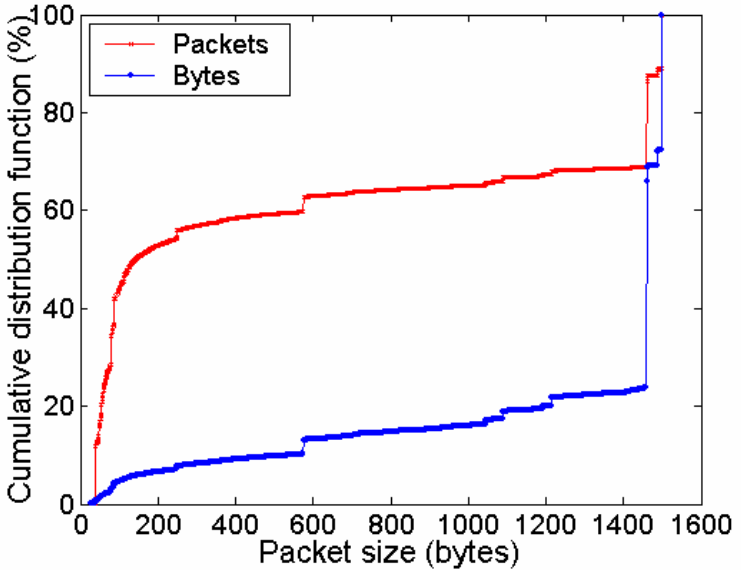


TCP connection level: Web traffic

- Traffic is non-uniformly distributed among the Internet hosts
- Ten busiest websites account for 60.23 % of the entire traffic load:
  - all registered under the Asia Pacific Network Information Centre
  - the most popular site: a Chinese search engine website
- Language, geographical, and commercial factors (popular sites) greatly affect the traffic distribution
- Important for designing content delivery networks and caching proxies

TCP packet size

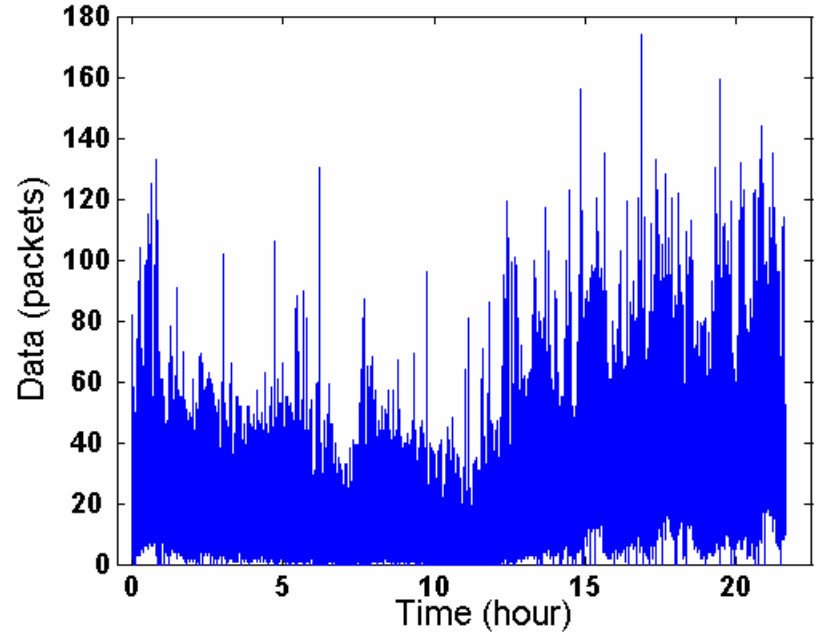


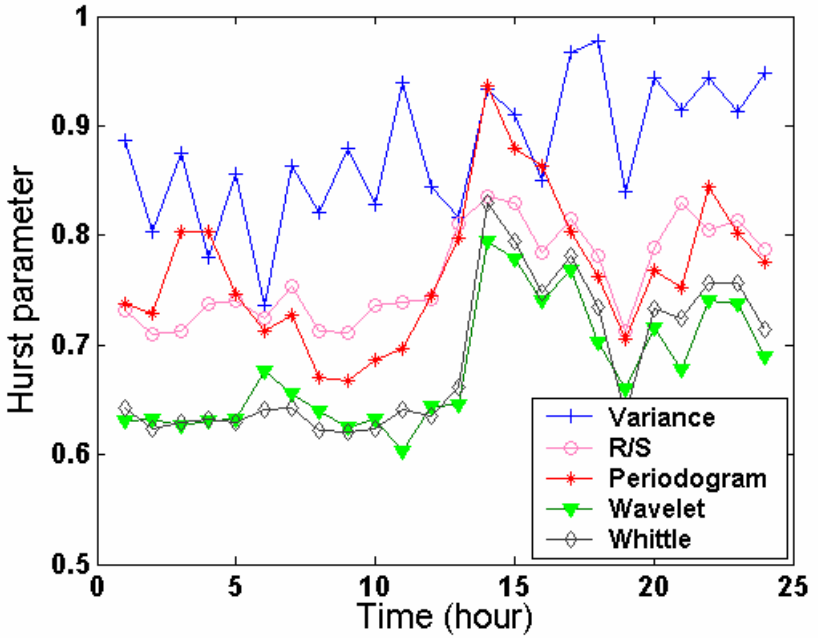


Traffic data was collected from 21-12-2002 22:08 to 23-12-2002 3:28

- Packet size distribution is bimodal:
  - 50 % of packets are less than 200 bytes
  - 30 % of packets are greater than 1,400 bytes
- Most bytes are transferred in large packets

Estimation of self-similarity





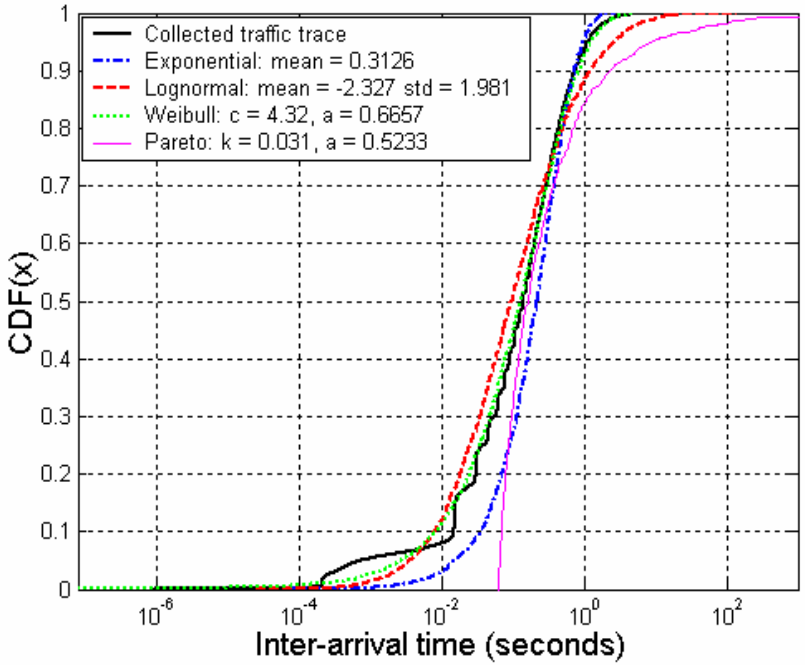
Traffic data was collected on 09-12-2002

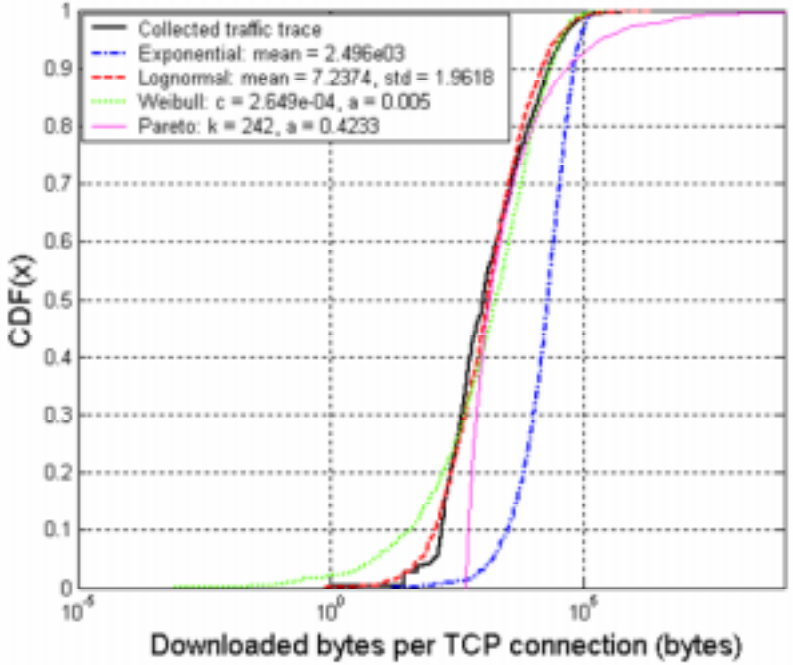
TCP connection model

- We consider two parameters of a TCP connection:
  - connection inter-arrival times
  - number of downloaded bytes per connection
- Four probability distributions:

Distribution	Probability density	Cumulative probability
Exponential	$f(x)=\frac{1}{\rho}e^{-x/\rho}$	$F(x)=1-e^{-x/\rho}$
Weibull	$f(x)=\frac{1}{a}\left(-\frac{x}{a}\right)^{c-1}\mathrm{e}^{-(x/a)^c}$	$F(x)=1-e^{-(x/a)^c}$
Pareto $(k>0, a>0; x\geq k)$	$f(x)=\frac{ak^a}{(x)^{k+1}}$	$F(x)=1-\left(\frac{k}{x}\right)^a$
Lognormal	$f(x)=\frac{1}{x\sqrt{2\pi\sigma}}e^{-[\log(x)-\xi]^2/2\sigma^2}$	No closed form

TCP connection model





- Best fit:
  - Lognormal: downloaded bytes per TCP connection
  - Weibull: inter-arrival times of TCP connections

Traffic prediction

- "Time series analysis - forecasting and control"
  - G. E. P. Box and G. M. Jenkins (1976)
- AutoRegressive Integrated Moving-Average (ARIMA):

$$X(t) = \phi_1 X(t-1) + \cdots + \phi_p X(t-p) + e(t) + \theta_1 e(t-1) \cdots + \theta_q e(t-q)$$

$$(p, d, q) \times (P, D, Q)_s$$

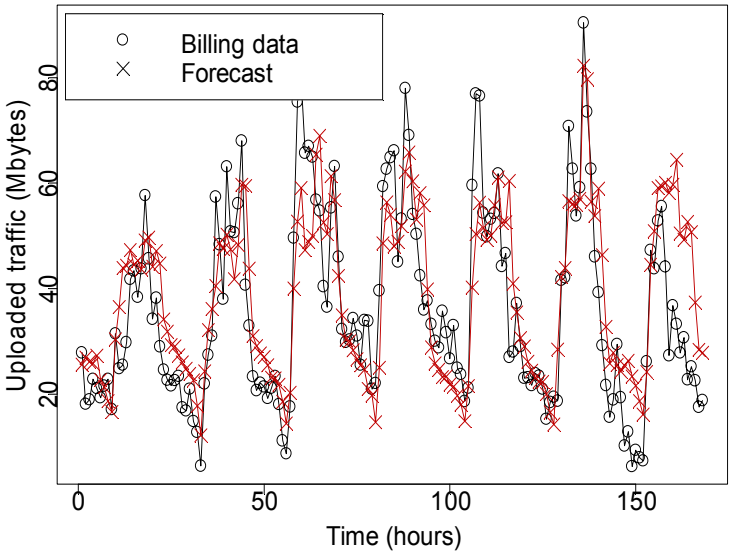
  - past values
    - AutoRegressive (AR) structure
  - past random fluctuant effect
    - Moving Average (MA) process

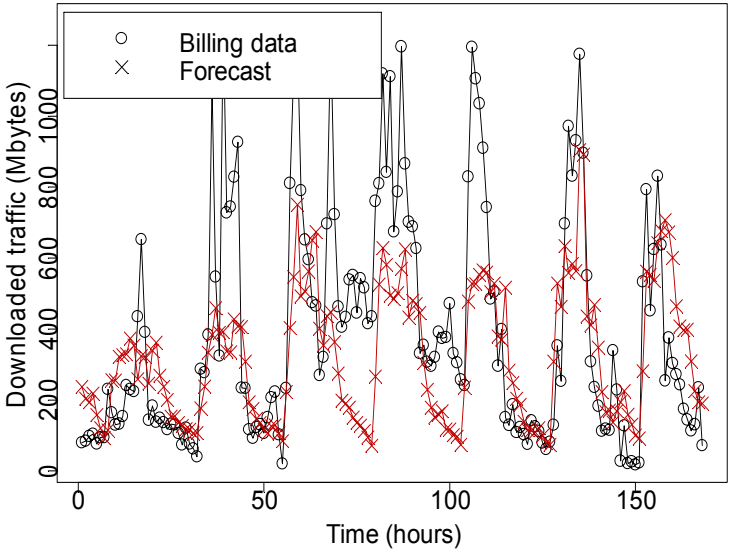
One week ahead prediction

- We applied Box-Jenkins method to six weeks of billing records
- Derived parameters:
  - d=0, D=1, s=168, p=1, q=0, P=0, Q=1
  - collected records fit the model $(1,0,0)\times(0,1,1)_{168}$
- Normalized mean squared error (nmse) is used to measure the performance of the predictor:

$$nmse = \frac{1}{\sigma^2 N}\sum_{k=1}^{N}(x(k)-\bar{x}(k))^2$$

Predictability evaluation



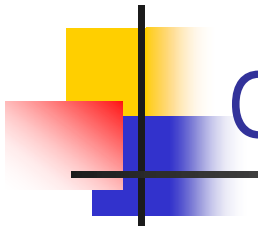


- Predicting downloaded traffic is more difficult than predicting uploaded traffic

Traffic type	Uploaded traffic	Downloaded traffic
nmse	0.3653	0.5988

Conclusions

- Analysis of collected traffic data:
  - Web applications and TCP protocol dominate the collected traffic
  - packet size distribution is bimodal: most bytes are transferred in big packets
  - few Web servers account for majority of data traffic
  - the frequency-rank relation of client connections matches the discrete lognormal distribution
  - various estimators of the Hurst parameter produced inconsistent results
  - more accurate estimation was achieved with the wavelet estimator



Conclusions

- TCP modeling:
  - Weibull: inter-arrival times of TCP connections
  - Lognormal: downloaded bytes per TCP connection
- Traffic prediction using the ARIMA model:
  - performs better for predicting the uploaded traffic
  - not suitable for predicting downloaded traffic

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tcpdump trace format

- timestamp src > dst: flags data-seqno ack window urgent options
  - 19:12:45.660701 61.159.59.162.12800 > 192.168.1.169.62246: udp 52
  - 19:12:45.672959 192.168.1.242.40849 > 210.51.17.67.9065: P 6541284:6541321(37) ack 1479344110 win 8192 (DF)
  - 19:12:45.674709 192.168.2.30.39042 > 202.101.165.124.4220: . ack 807850998 win 8192
  - 19:12:45.676255 61.152.249.71.55901 > 192.168.1.242.40770: P 2627573783:2627573791(8) ack 5795719 win 63343 (DF)
  - 19:12:45.676256 61.152.249.71.55901 > 192.168.1.242.40846: P 2775973525:2775973533(8) ack 11622145 win 64102 (DF)
  - 19:12:45.688514 192.168.1.242.40770 > 61.152.249.71.55901: . ack 8 win 8192
  - 19:12:45.688843 192.168.1.242.40846 > 61.152.249.71.55901: . ack 8 win 8192
  - 19:12:45.689095 192.168.1.169.63644 > 202.103.69.103.3010: P 1969195:1969259(64) ack 2995916216 win 8192 (DF)
  - 19:12:45.692475 202.101.165.134.80 > 192.168.2.3.45585: . ack 3153903 win 6432
  - 19:12:45.699193 207.46.104.20.80 > 192.168.1.239.4912: R 2405276149:2405276149(0) win 0

- Red: uploaded traffic
- Green: downloaded traffic

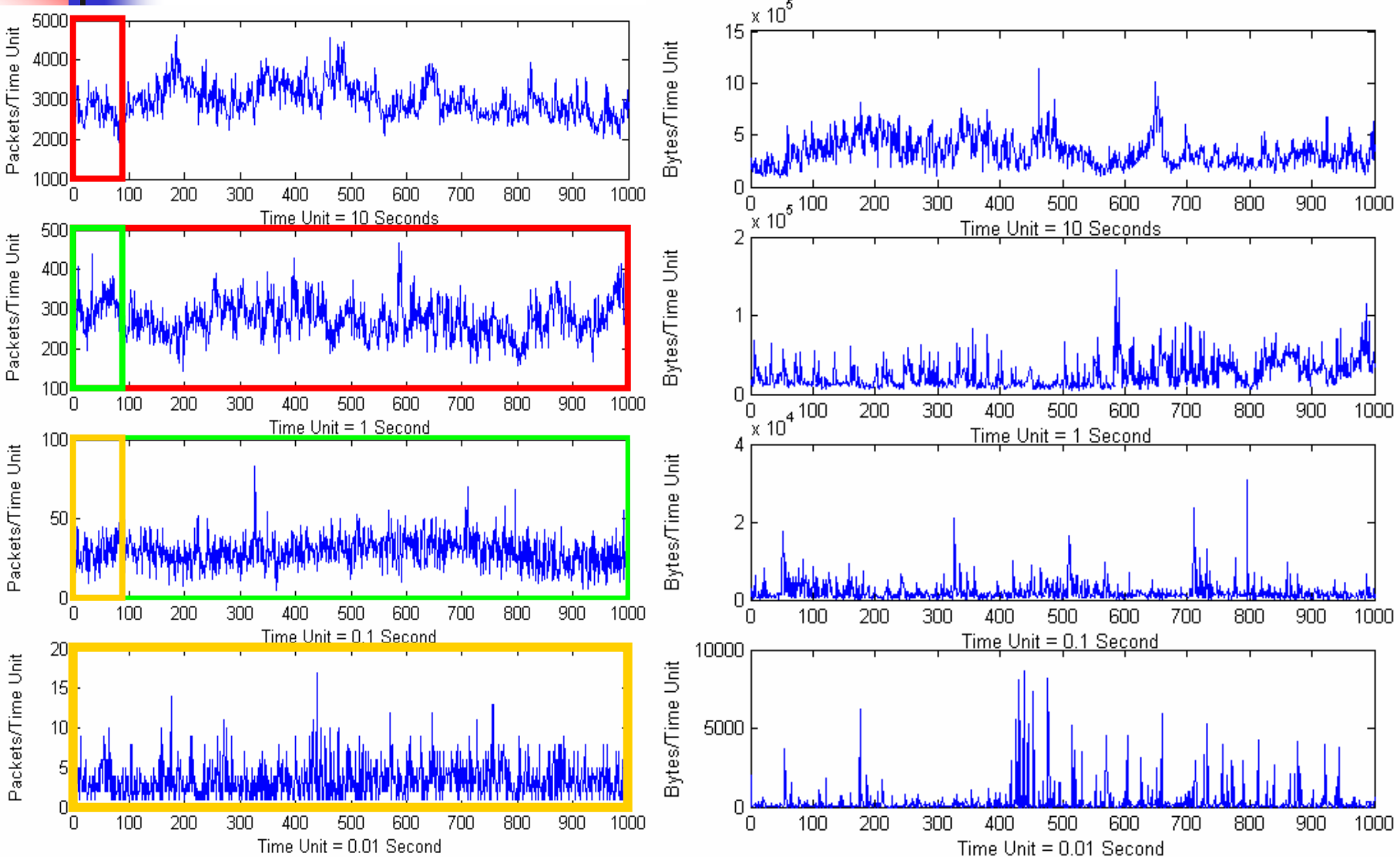
DirecPC system

- IP spoofing:
  - customer's requests are not directly sent to the website
  - they are rerouted to the satellite network operation center (NOC)
  - NOC resends the request to the website
  - website sends to the NOC data to be downloaded
- TCP splitting:
  - terrestrial links use standard TCP
  - to improve throughput, space links with long delay use modified TCP versions with enlarged TCP window size

Self-similarity

- Self-similarity implies a "fractal-like" behavior: data on various time scales have similar patterns
- A wide-sense stationary process X(n) is called (exactly second order) self-similar if:
  - $r^{(m)}(k) = r(k)$, $k \geqslant 0$, $m = 1, 2, \ldots, n$
- Implications:
  - no natural length of bursts
  - bursts exist across many time scales
  - traffic does not become "smoother" when aggregated (unlike Poisson traffic)

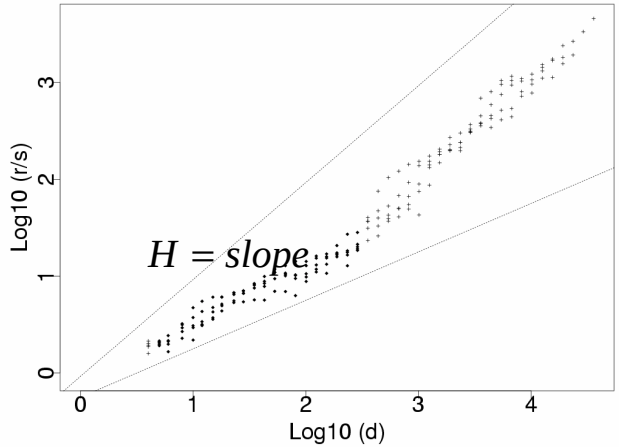
Estimation of self-similarity



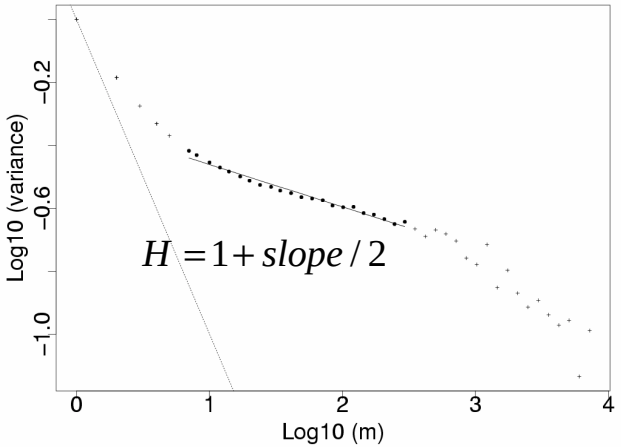
Self-similar processes

- Properties:
  - slow decaying variance
  - long-range dependence
  - Hurst parameter
- Processes with only short-range dependence (Poisson): H = 0.5
- Self-similar processes: 0.5 < H < 1.0
- As the traffic volume increases, the traffic becomes more bursty, more self-similar, and the Hurst parameter increases

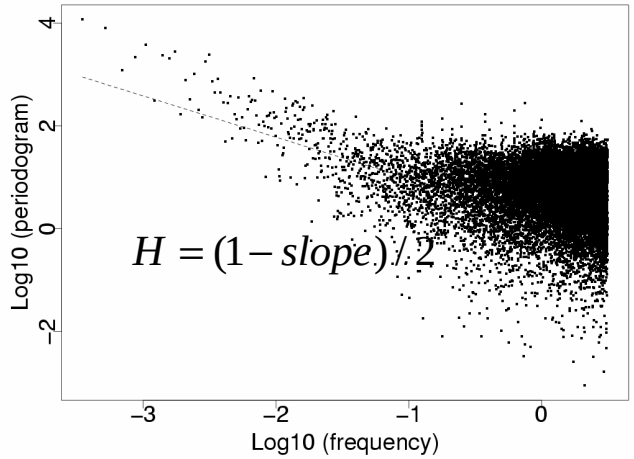
Estimation of self-similarity



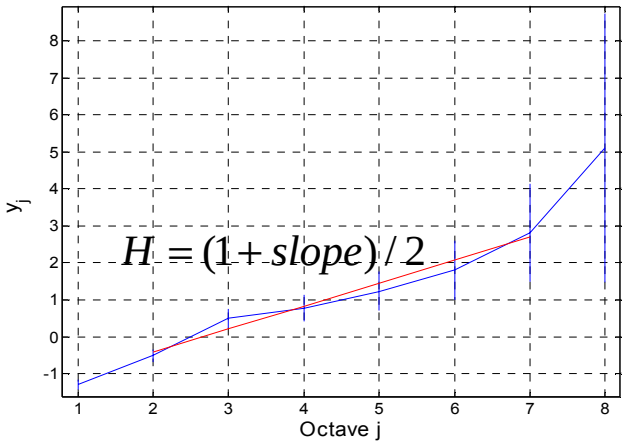
(a) R/S plot



(b) Variance-time plot



(c) Periodogram plot



(d) Wavelet plot

Modeling self-similar processes

- Self-similar process can be generated by aggregating multiple ON/OFF sources
- The ON/OFF periods are heavy-tailed distributed with infinite variance
- Web and ftp file sizes are heavy-tailed
- A probability distribution X is heavy-tailed if:

$$P[X > x] \sim cx^{-\alpha}, 0 < a < 2, x \to \infty$$

Reference: Mark E. Crovella and Azer Bestavros, “Self-similarity in world wide web traffic: evidence and possible causes,” in *IEEE/ACM Transactions on Networking*, vol. 5, no. 6, pp. 835 - 846, December 1997.